Red Hat Enterprise Linux 9 Managing, monitoring, and updating the kernel

A guide to managing the Linux kernel on Red Hat Enterprise Linux 9
Abstract

This document provides the users and administrators with necessary information about configuring their workstations on the Linux kernel level. Such adjustments bring performance enhancements, easier troubleshooting or optimized system.
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MAKING OPEN SOURCE MORE INCLUSIVE

Red Hat is committed to replacing problematic language in our code, documentation, and web properties. We are beginning with these four terms: master, slave, blacklist, and whitelist. Because of the enormity of this endeavor, these changes will be implemented gradually over several upcoming releases. For more details, see our CTO Chris Wright’s message.
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We appreciate your feedback on our documentation. Let us know how we can improve it.

Submitting comments on specific passages

1. View the documentation in the Multi-page HTML format and ensure that you see the Feedback button in the upper right corner after the page fully loads.
2. Use your cursor to highlight the part of the text that you want to comment on.
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1. Log in to the Bugzilla website.
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3. Enter a descriptive title in the Summary field.
4. Enter your suggestion for improvement in the Description field. Include links to the relevant parts of the documentation.
5. Click Submit Bug.
CHAPTER 1. THE LINUX KERNEL

The following sections bring information about the Linux kernel and the Linux kernel RPM package provided and maintained by Red Hat (Red Hat kernel). The integral part of this chapter is instruction on how to keep the Red Hat kernel updated, which ensures the operating system has all the latest bug fixes, performance enhancements, and patches, and is compatible with new hardware.

1.1. WHAT THE KERNEL IS

The kernel is a core part of a Linux operating system that manages the system resources and provides interface between hardware and software applications. The Red Hat kernel is a custom-built kernel based on the upstream Linux mainline kernel that Red Hat engineers further develop and harden with a focus on stability and compatibility with the latest technologies and hardware.

Before Red Hat releases a new kernel version, the kernel needs to pass a set of rigorous quality assurance tests.

The Red Hat kernels are packaged in the RPM format so that they are easily upgraded and verified by the `dnf` package manager.

WARNING

Kernels that have not been compiled by Red Hat are not supported by Red Hat.

1.2. WHAT AN RPM IS

An RPM package is a file containing other files and their metadata (information about the files that are needed by the system). Specifically, an RPM package consists of the `cpio` archive.

The `cpio` archive contains:

- Files
- RPM header (package metadata)
  The `rpm` package manager uses this metadata to determine dependencies, where to install files, and other information.

Types of RPM packages

There are two types of RPM packages. Both types share the file format and tooling, but have different contents and serve different purposes:

- Source RPM (SRPM)
  An SRPM contains source code and a SPEC file, which describes how to build the source code into a binary RPM. Optionally, the patches to source code are included as well.

- Binary RPM
  A binary RPM contains the binaries built from the sources and patches.

1.3. THE LINUX KERNEL RPM PACKAGE OVERVIEW
The kernel RPM is a meta package that does not contain any files, but rather ensures that the following required sub-packages are properly installed:

- **kernel-core** - contains the binary image of the kernel, all initramfs-related objects to bootstrap the system, and a minimal number of kernel modules to ensure core functionality. This sub-package alone could be used in virtualized and cloud environments to provide a Red Hat Enterprise Linux 9 kernel with a quick boot time and a small disk size footprint.

- **kernel-modules** - contains the remaining kernel modules that are not present in kernel-core.

The small set of kernel sub-packages above aims to provide a reduced maintenance surface to system administrators especially in virtualized and cloud environments.

Optional kernel packages are for example:

- **kernel-modules-extra** - contains kernel modules for rare hardware and modules which loading is disabled by default.

- **kernel-debug** – contains a kernel with numerous debugging options enabled for kernel diagnosis, at the expense of reduced performance.

- **kernel-tools** – contains tools for manipulating the Linux kernel and supporting documentation.

- **kernel-devel** – contains the kernel headers and makefiles sufficient to build modules against the kernel package.

- **kernel-abi-stablelists** – contains information pertaining to the RHEL kernel ABI, including a list of kernel symbols that are needed by external Linux kernel modules and a dnf plug-in to aid enforcement.

- **kernel-headers** – includes the C header files that specify the interface between the Linux kernel and user-space libraries and programs. The header files define structures and constants that are needed for building most standard programs.

Additional resources

- *What are the kernel-core, kernel-modules, and kernel-modules-extras packages?*

### 1.4. DISPLAING CONTENTS OF THE KERNEL PACKAGE

The following procedure describes how to view the contents of the kernel package and its sub-packages without installing them using the rpm command.

**Prerequisites**

- Obtained kernel, kernel-core, kernel-modules, kernel-modules-extra RPM packages for your CPU architecture

**Procedure**

- List modules for kernel:

  ```
  $ rpm -qlp <kernel_rpm>
  (contains no files)
  ...
  ```
List modules for **kernel-core**:

```bash
$ rpm -qlp <kernel-core_rpm>
```

```
…
/lib/modules/5.14.0-1.el9.x86_64/kernel/fs/udf/udf.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/kernel/fs/xfs
/lib/modules/5.14.0-1.el9.x86_64/kernel/fs/xfs/xfs.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/kernel/kernel
/lib/modules/5.14.0-1.el9.x86_64/kernel/kernel/trace
/lib/modules/5.14.0-1.el9.x86_64/kernel/kernel/trace/ring_buffer_benchmark.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/kernel/lib
/lib/modules/5.14.0-1.el9.x86_64/kernel/lib/cordic.ko.xz
…
```

List modules for **kernel-modules**:

```bash
$ rpm -qlp <kernel-modules_rpm>
```

```
…
/lib/modules/5.14.0-1.el9.x86_64/kernel/drivers/infiniband/hw/mlx4/mlx4_ib.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/kernel/drivers/infiniband/hw/mlx5/mlx5_ib.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/kernel/drivers/infiniband/hw/qedr/qedr.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/kernel/drivers/infiniband/hw/usnic/usnic_verbs.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/kernel/drivers/infiniband/hw/vmw_pvrdma/vmw_pvrdma.ko.xz
…
```

List modules for **kernel-modules-extra**:

```bash
$ rpm -qlp <kernel-modules-extra_rpm>
```

```
…
/lib/modules/5.14.0-1.el9.x86_64/extra/net/sched/sch_cbq.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/extra/net/sched/sch_choke.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/extra/net/sched/sch_drr.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/extra/net/sched/sch_dsmark.ko.xz
/lib/modules/5.14.0-1.el9.x86_64/extra/net/sched/sch_gred.ko.xz
…
```

Additional resources

- The **rpm(8)** manual page
- RPM packages

### 1.5. UPDATING THE KERNEL

The following procedure describes how to update the kernel using the **dnf** package manager.

**Procedure**

1. To update the kernel, enter the following command:

   ```bash
   # dnf update kernel
   ```

   This command updates the kernel along with all dependencies to the latest available version.
2. Reboot your system for the changes to take effect.

Additional resources

- package manager
- The `dnf(8)` manual page

### 1.6. INSTALLING SPECIFIC KERNEL VERSIONS

The following procedure describes how to install new kernels using the `dnf` package manager.

**Procedure**

- To install a specific kernel version, enter the following command:

  ```
  # dnf install kernel-{version}
  ```

**Additional resources**

- [Red Hat Code Browser](#)
- [Red Hat Enterprise Linux Release Dates](#)
CHAPTER 2. MANAGING KERNEL MODULES

The following sections explain what kernel modules are, how to display their information, and how to perform basic administrative tasks with kernel modules.

2.1. INTRODUCTION TO KERNEL MODULES

The Red Hat Enterprise Linux kernel can be extended with optional, additional pieces of functionality, called kernel modules, without having to reboot the system. On Red Hat Enterprise Linux 9, kernel modules are extra kernel code which is built into compressed `<KERNEL_MODULE_NAME>.ko.xz` object files.

The most common functionality enabled by kernel modules are:

- Device driver which adds support for new hardware
- Support for a file system such as GFS2 or NFS
- System calls

On modern systems, kernel modules are automatically loaded when needed. However, in some cases it is necessary to load or unload modules manually.

Like the kernel itself, the modules can take parameters that customize their behavior if needed.

Tooling is provided to inspect which modules are currently running, which modules are available to load into the kernel and which parameters a module accepts. The tooling also provides a mechanism to load and unload kernel modules into the running kernel.

2.2. KERNEL MODULE DEPENDENCIES

Certain kernel modules sometimes depend on one or more other kernel modules. The `/lib/modules/<KERNEL_VERSION>/modules.dep` file contains a complete list of kernel module dependencies for the respective kernel version.

The dependency file is generated by the `depmod` program, which is a part of the `kmod` package. Many of the utilities provided by `kmod` take module dependencies into account when performing operations so that manual dependency-tracking is rarely necessary.

**WARNING**

The code of kernel modules is executed in kernel-space in the unrestricted mode. Because of this, you should be mindful of what modules you are loading.

Additional resources

- `modules.dep(5)` manual page
- `depmod(8)` manual page
2.3. LISTING INSTALLED KERNEL MODULES

The `grubby --info=ALL` command displays an indexed list of installed kernels on !BLS and BLS installs.

Procedure

- List the installed kernels using the following command:

  ```
  # grubby --info=ALL | grep title
  ```

  The list of all installed kernels is displayed as follows:

  ```
  title="Red Hat Enterprise Linux (5.14.0-1.el9.x86_64) 9.0 (Plow)"
  title="Red Hat Enterprise Linux (0-rescue-0d772916a9724907a5d1350bcd39ac92) 9.0 (Plow)"
  ```

  The above example displays the installed kernels list of grubby-8.40-17, from the GRUB menu.

2.4. LISTING CURRENTLY LOADED KERNEL MODULES

The following procedure describes how to view the currently loaded kernel modules.

Prerequisites

- The `kmod` package is installed.

Procedure

- To list all currently loaded kernel modules, execute:

  ```
  $ lsmod
  ```

  Module                  Size  Used by  
  fuse                  126976  3
  uinput                 20480  1
  xt_CHECKSUM            16384  1
  ipt_masquerade         16384  1
  xt_contrack            16384  1
  ipt_REJECT             16384  1
  nft_counter            16384  16
  nf_nat_tftp            16384  0
  nf_contrack_tftp       16384  1 nf_nat_tftp
  tun                    49152  1
  bridge                 192512  0
  stp                    16384  1 bridge
  llc                    16384  2 bridge,stp
  nf_tables_set          32768  5
  nf_fib_inet            16384  1
  ...

  In the example above:

  - The first column provides the names of currently loaded modules.
The second column displays the amount of memory per module in kilobytes.

The last column shows the number, and optionally the names of modules that are dependent on a particular module.

Additional resources

- `/usr/share/doc/kmod/README` file
- `lsmod(8)` manual page

### 2.5. SETTING A KERNEL AS DEFAULT

The following procedure describes how to set a specific kernel as default using the `grubby` command-line tool and GRUB.

#### Procedure

- **Setting the kernel as default, using the `grubby` tool**
  
  Execute the following command to set the kernel as default using the `grubby` tool:
  
  ```
  # grubby --set-default $kernel_path
  ```

  The command uses a machine ID without the `.conf` suffix as an argument.

  **NOTE**
  
  The machine ID is located in the `/boot/loader/entries/` directory.

- **Setting the kernel as default, using the `id` argument**
  
  List the boot entries using the `id` argument and then set an intended kernel as default:
  
  ```
  # grubby --info ALL | grep id
  # grubby --set-default /boot/vmlinuz-<version>.<architecture>
  ```

  **NOTE**
  
  To list the boot entries using the `title` argument, execute the `# grubby --info=ALL | grep title` command.

- **Setting the default kernel for only the next boot**
  
  Execute the following command to set the default kernel for only the next reboot using the `grub2-reboot` command:
  
  ```
  # grub2-reboot <index|title|id>
  ```
2.6. DISPLAYING INFORMATION ABOUT KERNEL MODULES

When working with a kernel module, you may want to see further information about that module. This procedure describes how to display extra information about kernel modules.

Prerequisites

- The `kmod` package is installed.

Procedure

- To display information about any kernel module, execute:

  ```
  $ modinfo <KERNEL_MODULE_NAME>
  ```

  For example:

  ```
  $ modinfo virtio_net
  ```

  ```
  filename:       /lib/modules/5.14.0-1.el9.x86_64/kernel/drivers/net/virtio_net.ko.xz
  license:        GPL
  description:    Virtio network driver
  rhelversion:    9.0
  srcversion:     8809CDDBE7202A1B00B9F1C
  alias:          virtio:d00000001v*
  depends:        net_failover
  retpoline:      Y
  intree:         Y
  name:           virtio_net
  vermagic:       5.14.0-1.el9.x86_64 SMP mod_unload modversions
  ...
  parm:           napi_weight:int
  parm:           csum:bool
  parm:           gso:bool
  parm:           napi_tx:bool
  ```

The `modinfo` command displays some detailed information about the specified kernel module. You can query information about all available modules, regardless of whether they are loaded or not. The `parm` entries show parameters the user is able to set for the module, and what type of value they expect.
NOTE

When entering the name of a kernel module, do not append the .ko.xz extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

Additional resources

- modinfo(8) manual page

2.7. LOADING KERNEL MODULES AT SYSTEM RUNTIME

The optimal way to expand the functionality of the Linux kernel is by loading kernel modules. The following procedure describes how to use the modprobe command to find and load a kernel module into the currently running kernel.

Prerequisites

- Root permissions
- The kmod package is installed.
- The respective kernel module is not loaded. To ensure this is the case, list the loaded kernel modules.

Procedure

1. Select a kernel module you want to load. The modules are located in the /lib/modules/$(uname -r)/kernel/<SUBSYSTEM>/ directory.

2. Load the relevant kernel module:

   ```
   # modprobe <MODULE_NAME>
   ```

   NOTE

   When entering the name of a kernel module, do not append the .ko.xz extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

3. Optionally, verify the relevant module was loaded:

   ```
   $ lsmod | grep <MODULE_NAME>
   ```

   If the module was loaded correctly, this command displays the relevant kernel module. For example:

   ```
   $ lsmod | grep serio_raw
   serio_raw       16384     0
   ```
The changes described in this procedure will not persist after rebooting the system. For information on how to load kernel modules to persist across system reboots, see Loading kernel modules automatically at system boot time.

Additional resources

- `modprobe(8)` manual page

## 2.8. UNLOADING KERNEL MODULES AT SYSTEM RUNTIME

At times, you find that you need to unload certain kernel modules from the running kernel. The following procedure describes how to use the `modprobe` command to find and unload a kernel module at system runtime from the currently loaded kernel.

### Prerequisites

- Root permissions
- The `kmod` package is installed.

### Procedure

1. Execute the `lsmod` command and select a kernel module you want to unload. If a kernel module has dependencies, unload those prior to unloading the kernel module. For details on identifying modules with dependencies, see Listing currently loaded kernel modules and Kernel module dependencies.

2. Unload the relevant kernel module:

   ```
   # modprobe -r `<MODULE_NAME>`
   ```

   When entering the name of a kernel module, do not append the `.ko.xz` extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

   **WARNING**

   Do not unload kernel modules when they are used by the running system. Doing so can lead to an unstable or non-operational system.

3. Optionally, verify the relevant module was unloaded:

   ```
   $ lsmod | grep `<MODULE_NAME>`
   ```

   If the module was unloaded successfully, this command does not display any output.
IMPORTANT

After finishing this procedure, the kernel modules that are defined to be automatically loaded on boot, **will not stay unloaded** after rebooting the system. For information on how to counter this outcome, see Preventing kernel modules from being automatically loaded at system boot time.

Additional resources

- modprobe(8) manual page

2.9. UNLOADING KERNEL MODULES AT EARLY STAGES OF THE BOOT PROCESS

In certain situations it is necessary to unload a kernel module very early in the booting process. For example, when the kernel module contains a code, which causes the system to become unresponsive, and the user is not able to reach the stage to permanently disable the rogue kernel module. In that case it is possible to temporarily block the loading of the kernel module using a bootloader.

IMPORTANT

The changes described in this procedure **will not persist** after the next reboot. For information on how to add a kernel module to a denylist so that it will not be automatically loaded during the boot process, see Preventing kernel modules from being automatically loaded at system boot time.

Prerequisites

- You have a loadable kernel module, which you want to prevent from loading for some reason.

Procedure

- Edit the relevant bootloader entry to unload the desired kernel module before the booting sequence continues.
  - Use the cursor keys to highlight the relevant bootloader entry.
  - Press **e** key to edit the entry.
Use the cursor keys to navigate to the line that starts with `linux`.

- Append `modprobe.blacklist=module_name` to the end of the line.

The `serio_raw` kernel module illustrates a rogue module to be unloaded early in the boot process.

- Press `CTRL+x` keys to boot using the modified configuration.

**Verification**

- Once the system fully boots, verify that the relevant kernel module is not loaded.

```
# lsmod | grep serio_raw
```

**Additional resources**

- Managing kernel modules
The following procedure describes how to configure a kernel module so that it is loaded automatically during the boot process.

**Prerequisites**

- Root permissions
- The `kmod` package is installed.

**Procedure**

1. Select a kernel module you want to load during the boot process. The modules are located in the `/lib/modules/$(uname -r)/kernel/<SUBSYSTEM>/` directory.

2. Create a configuration file for the module:

   ```
   # echo <MODULE_NAME> > /etc/modules-load.d/<MODULE_NAME>.conf
   ```

   **NOTE**

   When entering the name of a kernel module, do not append the `.ko.xz` extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

3. Optionally, after reboot, verify the relevant module was loaded:

   ```
   $ lsmod | grep <MODULE_NAME>
   ```

   The example command above should succeed and display the relevant kernel module.

   **IMPORTANT**

   The changes described in this procedure will persist after rebooting the system.

**Additional resources**

- `modules-load.d(5)` manual page

---

2.11. PREVENTING KERNEL MODULES FROM BEING AUTOMATICALLY LOADED AT SYSTEM BOOT TIME

The following procedure describes how to add a kernel module to a denylist so that it will not be automatically loaded during the boot process.

**Prerequisites**

- Root permissions
- The `kmod` package is installed.
- Ensure that a kernel module in a denylist is not vital for your current system configuration.
Procedure

1. Select a kernel module that you want to put in a denylist:

```
$ lsmod
Module                  Size  Used by
fuse                  126976  3
xt_CHECKSUM            16384  1
ipt_MASQUERADE         16384  1
uinput                 20480  1
xt_conntrack           16384  1
...
```

The `lsmod` command displays a list of modules loaded to the currently running kernel.

- Alternatively, identify an unloaded kernel module you want to prevent from potentially loading.
  All kernel modules are located in the `/lib/modules/<KERNEL_VERSION>/kernel/<SUBSYSTEM>/` directory.

2. Create a configuration file for a denylist:

```
# vim /etc/modprobe.d/blacklist.conf

# Blacklists <KERNEL_MODULE_1>
blacklist <MODULE_NAME_1>
install <MODULE_NAME_1> /bin/false

# Blacklists <KERNEL_MODULE_2>
blacklist <MODULE_NAME_2>
install <MODULE_NAME_2> /bin/false

# Blacklists <KERNEL_MODULE_n>
blacklist <MODULE_NAME_n>
install <MODULE_NAME_n> /bin/false
...
```

The example shows the contents of the `blacklist.conf` file, edited by the `vim` editor. The `blacklist` line ensures that the relevant kernel module will not be automatically loaded during the boot process. The `blacklist` command, however, does not prevent the module from being loaded as a dependency for another kernel module that is not in a denylist. Therefore the `install` line causes the `/bin/false` to run instead of installing a module.

The lines starting with a hash sign are comments to make the file more readable.

**NOTE**

When entering the name of a kernel module, do not append the `.ko.xz` extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

3. Create a backup copy of the current initial ramdisk image before rebuilding:
The command above creates a backup initramfs image in case the new version has an unexpected problem.

- Alternatively, create a backup copy of other initial ramdisk image which corresponds to the kernel version for which you want to put kernel modules in a denylist:

```
# cp /boot/initramfs-<SOME_VERSION>.img /boot/initramfs-<SOME_VERSION>.img.bak.$(date +%m-%d-%H%M%S)
```

4. Generate a new initial ramdisk image to reflect the changes:

```
# dracut -f -v
```

- If you are building an initial ramdisk image for a different kernel version than you are currently booted into, specify both target initramfs and kernel version:

```
# dracut -f -v /boot/initramfs-<TARGET_VERSION>.img <CORRESPONDING_TARGET_KERNEL_VERSION>
```

5. Reboot the system:

```
$ reboot
```

**IMPORTANT**

The changes described in this procedure will take effect and persist after rebooting the system. If you improperly put a key kernel module in a denylist, you can face an unstable or non-operational system.

Additional resources

- How do I prevent a kernel module from loading automatically?
- dracut(8) manual page

### 2.12. COMPILING CUSTOM KERNEL MODULES

You can build a sampling kernel module as requested by various configurations at hardware and software level.

**Prerequisites**

- You installed the kernel-devel, gcc, and elfutils-libelf-devel packages.
- You have root permissions.
- You created the /root/testmodule/ directory where you compile the custom kernel module.

**Procedure**
1. Create the `/root/testmodule/test.c` file with the following content.

```c
#include <linux/module.h>
#include <linux/kernel.h>

int init_module(void)
{ printk("Hello World\n This is a test\n"); return 0; }

void cleanup_module(void)
{ printk("Good Bye World"); }

MODULE_LICENSE("GPL");
```

The `test.c` file is a source file that provides the main functionality to the kernel module. The file has been created in a dedicated `/root/testmodule/` directory for organizational purposes. After the module compilation, the `/root/testmodule/` directory will contain multiple files.

The `test.c` file includes from the system libraries:

- The `linux/kernel.h` header file is necessary for the `printk()` function in the example code.
- The `linux/module.h` file contains function declarations and macro definitions to be shared between several source files written in C programming language.

Next follow the `init_module()` and `cleanup_module()` functions to start and end the kernel logging function `printk()`, which prints text.

2. Create the `/root/testmodule/Makefile` file with the following content.

```make
obj-m := test.o
```

The Makefile contains instructions that the compiler has to produce an object file specifically named `test.o`. The `obj-m` directive specifies that the resulting `test.ko` file is going to be compiled as a loadable kernel module. Alternatively, the `obj-y` directive would instruct to build `test.ko` as a built-in kernel module.

3. Compile the kernel module.

```bash
# make -C /lib/modules/$(uname -r)/build M=/root/testmodule modules
make: Entering directory '/usr/src/kernels/5.14.0-70.17.1.el9_0.x86_64'
CC [M] /root/testmodule/test.o
MODPOST /root/testmodule/Module.symvers
CC [M] /root/testmodule/test.mod.o
LD [M] /root/testmodule/test.ko
BTF [M] /root/testmodule/test.ko
Skipping BTF generation for /root/testmodule/test.ko due to unavailability of vmlinux
make: Leaving directory '/usr/src/kernels/5.14.0-70.17.1.el9_0.x86_64'
```

The compiler creates an object file (`test.o`) for each source file (`test.c`) as an intermediate step before linking them together into the final kernel module (`test.ko`).

After a successful compilation, `/root/testmodule/` contains additional files that relate to the compiled custom kernel module. The compiled module itself is represented by the `test.ko` file.

**Verification**
1. Optional: check the contents of the `/root/testmodule/` directory:

```
# ls -l /root/testmodule/
total 152
-rw-r--r--. 1 root root    16 Jul 26 08:19 Makefile
-rw-r--r--. 1 root root    25 Jul 26 08:20 modules.order
-rw-r--r--. 1 root root    0 Jul 26 08:20 Module.symvers
-rw-r--r--. 1 root root   224 Jul 26 08:18 test.c
-rw-r--r--. 1 root root  62176 Jul 26 08:20 test.ko
-rw-r--r--. 1 root root    25 Jul 26 08:20 test.mod
-rw-r--r--. 1 root root    849 Jul 26 08:20 test.mod.c
-rw-r--r--. 1 root root  50936 Jul 26 08:20 test.mod.o
-rw-r--r--. 1 root root  12912 Jul 26 08:20 test.o
```

2. Copy the kernel module to the `/lib/modules/$(uname -r)/` directory:

```
# cp /root/testmodule/test.ko /lib/modules/$(uname -r)/
```

3. Update the modular dependency list:

```
# depmod -a
```

4. Load the kernel module:

```
# modprobe -v test
insmod /lib/modules/5.14.0-1.el9.x86_64/test.ko
```

5. Verify that the kernel module was successfully loaded:

```
# lsmod | grep test
    test       16384  0
```

6. Read the latest messages from the kernel ring buffer:

```
# dmesg
[74422.545004] Hello World
    This is a test
```

Additional resources

- Managing kernel modules
You can enhance the security of your system by using signed kernel modules. The following sections describe how to self-sign privately built kernel modules for use with RHEL 9 on UEFI-based build systems where Secure Boot is enabled. These sections also provide an overview of available options for importing your public key into a target system where you want to deploy your kernel modules.

If Secure Boot is enabled, the UEFI operating system boot loaders, the Red Hat Enterprise Linux kernel, and all kernel modules have to be signed with a private key and authenticated with the corresponding public key. If they are not signed and authenticated, the system will not be allowed to finish the booting process.

The RHEL 9 distribution includes:

- Signed boot loaders
- Signed kernels
- Signed kernel modules

In addition, the signed first-stage boot loader and the signed kernel include embedded Red Hat public keys. These signed executable binaries and embedded keys enable RHEL 9 to install, boot, and run with the Microsoft UEFI Secure Boot Certification Authority keys that are provided by the UEFI firmware on systems that support UEFI Secure Boot. Note that not all UEFI-based systems include support for Secure Boot.

**Prerequisites**

To be able to sign externally built kernel modules, install the utilities listed in the following table on the build system.

<table>
<thead>
<tr>
<th>Utility</th>
<th>Provided by package</th>
<th>Used on</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td>openssl</td>
<td>openssl</td>
<td>Build system</td>
<td>Generates public and private X.509 key pair</td>
</tr>
<tr>
<td>sign-file</td>
<td>kernel-devel</td>
<td>Build system</td>
<td>Executable file used to sign a kernel module with the private key</td>
</tr>
<tr>
<td>mokutil</td>
<td>mokutil</td>
<td>Target system</td>
<td>Optional utility used to manually enroll the public key</td>
</tr>
<tr>
<td>keyctl</td>
<td>keyutils</td>
<td>Target system</td>
<td>Optional utility used to display public keys in the system keyring</td>
</tr>
</tbody>
</table>

**NOTE**

The build system, where you build and sign your kernel module, does not need to have UEFI Secure Boot enabled and does not even need to be a UEFI-based system.
3.1. WHAT IS UEFI SECURE BOOT

With the Unified Extensible Firmware Interface (UEFI) Secure Boot technology you can ensure that the system boot loader is signed with a cryptographic key. The database of public keys which is contained in the firmware, authorizes the signing key. You can verify a signature in the next-stage boot loader and the kernel. As a result, you can prevent the execution of the kernel space code which has not been signed by a trusted key.

UEFI Secure Boot establishes a chain of trust from the firmware to the signed drivers and kernel modules as follows:

- A UEFI private key signs, and a public key authenticates the *shim* first-stage boot loader. A certificate authority (CA) in turn signs the public key. The CA is stored in the firmware database.

- The *shim.efi* contains the Red Hat public key *Red Hat Secure Boot (CA key 1)* to authenticate the GRUB boot loader, and the kernel.

- The kernel in turn contains public keys to authenticate drivers and modules.

Secure Boot is the boot path validation component of the UEFI specification. The specification defines:

- Programming interface for cryptographically protected UEFI variables in non-volatile storage
- Storing the trusted X.509 root certificates in UEFI variables
- Validation of UEFI applications like boot loaders and drivers
- Procedures to revoke known-bad certificates and application hashes

UEFI Secure Boot does not prevent installation or removal of second-stage boot loaders. Also, it does not require explicit user confirmation of such changes. Signatures are verified during booting, not when the boot loader is installed or updated. Therefore, UEFI Secure Boot does not stop boot path manipulations. It helps in the detection of unauthorized changes.

**NOTE**

The boot loader or the kernel work as long as a system trusted key signs them.

3.2. UEFI SECURE BOOT SUPPORT

RHEL 9 includes support for the Unified Extensible Firmware Interface (UEFI) Secure Boot feature. It means that you can install and run RHEL 9 on machines with enabled UEFI Secure Boot. On these systems, you must sign the all loaded drivers with a trusted key. Otherwise the system will not accept them. Red Hat provides drivers which are signed and authenticated by the relevant Red Hat keys.

If you want to load externally built drivers you must sign them as well.

**Restrictions Imposed by UEFI Secure Boot**

- The system only runs the kernel mode code after its signature has been properly authenticated.

- GRUB module loading is disabled because there is no infrastructure for signing and verification of GRUB modules. Allowing them to be loaded constitutes execution of untrusted code inside the security perimeter that Secure Boot defines.

- Red Hat provides a signed GRUB binary that contains all the supported modules on RHEL 9.
Additional resources

- Restrictions Imposed by UEFI Secure Boot

### 3.3. REQUIREMENTS FOR AUTHENTICATING KERNEL MODULES WITH X.509 KEYS

In RHEL 9, when a kernel module is loaded, the kernel checks the signature of the module against the public X.509 keys from the kernel system keyring (.builtin_trusted_keys) and the kernel platform keyring (.platform). The .platform keyring contains keys from third-party platform providers and custom public keys. The keys from the kernel system .blacklist keyring are excluded from verification. The following sections provide an overview of sources of keys, keyrings and examples of loaded keys from different sources in the system. Also, you can see how to authenticate a kernel module.

You need to meet certain conditions to load kernel modules on systems with enabled UEFI Secure Boot functionality.

If UEFI Secure Boot is enabled or if the module.sig_enforce kernel parameter has been specified:

- You can only load those signed kernel modules whose signatures were authenticated against keys from the system keyring (.builtin_trusted_keys) and the platform keyring (.platform).
- The public key must not be on the system revoked keys keyring (.blacklist).

If UEFI Secure Boot is disabled and the module.sig_enforce kernel parameter has not been specified:

- You can load unsigned kernel modules and signed kernel modules without a public key.

If the system is not UEFI-based or if UEFI Secure Boot is disabled:

- Only the keys embedded in the kernel are loaded onto .builtin_trusted_keys and .platform.
- You have no ability to augment that set of keys without rebuilding the kernel.

<table>
<thead>
<tr>
<th>Module signed</th>
<th>Public key found and signature valid</th>
<th>UEFI Secure Boot state</th>
<th>sig_enforce</th>
<th>Module load</th>
<th>Kernel tainted</th>
</tr>
</thead>
<tbody>
<tr>
<td>Unsigned</td>
<td>-</td>
<td>Not enabled</td>
<td>Not enabled</td>
<td>Succeeds</td>
<td>Yes</td>
</tr>
<tr>
<td></td>
<td>Not enabled</td>
<td>Enabled</td>
<td>Fails</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Enabled</td>
<td>-</td>
<td>Fails</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>Signed</td>
<td>No</td>
<td>Not enabled</td>
<td>Not enabled</td>
<td>Succeeds</td>
<td>Yes</td>
</tr>
<tr>
<td></td>
<td>Not enabled</td>
<td>Enabled</td>
<td>Fails</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Enabled</td>
<td>-</td>
<td>Fails</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>Signed</td>
<td>Yes</td>
<td>Not enabled</td>
<td>Not enabled</td>
<td>Succeeds</td>
<td>No</td>
</tr>
</tbody>
</table>

Table 3.2. Kernel module authentication requirements for loading
### Table 3.3. Sources for system keyrings

<table>
<thead>
<tr>
<th>Source of X.509 keys</th>
<th>User can add keys</th>
<th>UEFI Secure Boot state</th>
<th>Keys loaded during boot</th>
</tr>
</thead>
<tbody>
<tr>
<td>Embedded in kernel</td>
<td>No</td>
<td>-</td>
<td>.builtin_trusted_keys</td>
</tr>
<tr>
<td>UEFI Secure Boot &quot;db&quot;</td>
<td>Limited</td>
<td>Not enabled</td>
<td>No</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Enabled</td>
<td>.platform</td>
</tr>
<tr>
<td>Embedded in shim.efi boot loader</td>
<td>No</td>
<td>Not enabled</td>
<td>No</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Enabled</td>
<td>.builtin_trusted_keys</td>
</tr>
<tr>
<td>Machine Owner Key (MOK) list</td>
<td>Yes</td>
<td>Not enabled</td>
<td>No</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Enabled</td>
<td>.platform</td>
</tr>
</tbody>
</table>

**.builtin_trusted_keys:**

- a keyring that is built on boot
- contains trusted public keys
- root privileges are needed to view the keys

**.platform:**

- a keyring that is built on boot
- contains keys from third-party platform providers and custom public keys
- **root** privileges are needed to view the keys

**blacklist**

- a keyring with X.509 keys which have been revoked
- a module signed by a key from **blacklist** will fail authentication even if your public key is in **builtin_trusted_keys**

UEFI Secure Boot db:

- a signature database
- stores keys (hashes) of UEFI applications, UEFI drivers, and bootloaders
- the keys can be loaded on the machine

UEFI Secure Boot dbx:

- a revoked signature database
- prevents keys from being loaded
- the revoked keys from this database are added to the **blacklist** keyring

### 3.5. GENERATING A PUBLIC AND PRIVATE KEY PAIR

You need to generate a public and private X.509 key pair to succeed in your efforts of using kernel modules on a Secure Boot-enabled system. You will later use the private key to sign the kernel module. You will also have to add the corresponding public key to the Machine Owner Key (MOK) for Secure Boot to validate the signed module.

Some of the parameters for this key pair generation are best specified with a configuration file.

**Procedure**

1. Create a configuration file with parameters for the key pair generation:

```bash
# cat << EOF > configuration_file.config
[ req ]
default_bits = 4096
distinguished_name = req_distinguished_name
prompt = no
string_mask = utf8only
x509_extensions = myexts

[ req_distinguished_name ]
O = Organization
CN = Organization signing key
e-mailAddress = E-mail address

[ myexts ]
basicConstraints=critical,CA:FALSE
keyUsage=digitalSignature
EOF
```
2. Create an X.509 public and private key pair as shown in the following example:

```bash
# openssl req -x509 -new -nodes -utf8 -sha256 -days 36500 \
   -batch -config configuration_file.config -outform DER \
   -out my_signing_key_pub.der \
   -keyout my_signing_key.priv
```

The public key will be written to the `my_signing_key_pub.der` file and the private key will be written to the `my_signing_key.priv` file.

**IMPORTANT**

In RHEL 9, the validity dates of the key pair do not matter. The key does not expire, but it is recommended for good security practices that the kernel module be signed within the validity period of its signing key. However, the `sign-file` utility will not warn you and the key will be usable regardless of the validity dates.

3. Optionally, you can review the validity dates of your public keys like in the example below:

```bash
# openssl x509 -inform der -text -noout -in my_signing_key_pub.der
```

**Validity**

- Not Before: Feb 14 16:34:37 2019 GMT
- Not After : Feb 11 16:34:37 2029 GMT

4. Enroll your public key on all systems where you want to authenticate and load your kernel module.

**WARNING**

Apply strong security measures and access policies to guard the contents of your private key. In the wrong hands, the key could be used to compromise any system which is authenticated by the corresponding public key.

**Additional resources**

- `openssl(1)` manual page
- `RHEL Security Guide`
- Enrolling public key on target system by adding the public key to the MOK list

**3.6. EXAMPLE OUTPUT OF SYSTEM KEYRINGS**
You can display information about the keys on the system keyrings using the `keyctl` utility from the `keyutils` package.

The following is a shortened example output of `.builtin_trusted_keys`, `.platform`, and `.blacklist` keyrings from a RHEL 9 system where UEFI Secure Boot is enabled.

Prerequisites

- You have root permissions.
- You installed the `keyctl` utility from the `keyutils` package.

```bash
# keyctl list %:.builtin_trusted_keys
6 keys in keyring:
...asymmetric: Red Hat Enterprise Linux Driver Update Program (key 3): bf57f3e87...
...asymmetric: Red Hat Secure Boot (CA key 1): 4016841644ce3a810408050766e8f8a29...
...asymmetric: Microsoft Corporation UEFI CA 2011: 13adbf4309bd82709c8cd54f316ed...
...asymmetric: Microsoft Windows Production PCA 2011: a92902398e16c49778cd90f99e...
...asymmetric: Red Hat Enterprise Linux kernel signing key: 4249689eefc77e95880b...
...asymmetric: Red Hat Enterprise Linux kpatch signing key: 4d38fd864ebe18c5f0b7...
```

```bash
# keyctl list %:.platform
4 keys in keyring:
...asymmetric: VMware, Inc.: 4ad8da0472073...
...asymmetric: Red Hat Secure Boot CA 5: cc6fafe72...
...asymmetric: Microsoft Windows Production PCA 2011: a929f298e1...
...asymmetric: Microsoft Corporation UEFI CA 2011: 13adbf4e0bd82...
```

```bash
# keyctl list %:.blacklist
4 keys in keyring:
...blacklist: bin:f5ff83a...
...blacklist: bin:0dfdbec...
...blacklist: bin:38f1d22...
...blacklist: bin:51f831f...
```

The `.builtin_trusted_keys` keyring above shows the addition of two keys from the UEFI Secure Boot "db" keys as well as the Red Hat Secure Boot (CA key 1), which is embedded in the `shim.efi` boot loader.

The following example shows the kernel console output. The messages identify the keys with a UEFI Secure Boot related source. These include UEFI Secure Boot db, embedded shim, and MOK list.

```bash
# dmesg | egrep 'integrity.*cert'
[1.512966] integrity: Loading X.509 certificate: UEFI:db
[1.513027] integrity: Loaded X.509 cert 'Microsoft Windows Production PCA 2011: a92902398e16c49778cd90f99e...
[1.513028] integrity: Loading X.509 certificate: UEFI:db
[1.513057] integrity: Loaded X.509 cert 'Microsoft Corporation UEFI CA 2011: 13adbf4309bd82709c8cd54f316ed...
[1.513298] integrity: Loading X.509 certificate: UEFI:MokListRT (MOKvar table)
[1.513549] integrity: Loaded X.509 cert 'Red Hat Secure Boot CA 5: cc6fa5e72868ba494e93...
```

Additional resources

- `keyctl(1), dmesg(1)` manual pages
3.7. ENROLLING PUBLIC KEY ON TARGET SYSTEM BY ADDING THE PUBLIC KEY TO THE MOK LIST

When RHEL 9 boots on a UEFI-based system with Secure Boot enabled, the kernel loads onto the system keyring (.builtin_trusted_keys) all public keys that are in the Secure Boot db key database. At the same time the kernel excludes the keys in the dbx database of revoked keys. The sections below describe different ways of importing a public key on a target system so that the system keyring (.builtin_trusted_keys) is able to use the public key to authenticate a kernel module.

The Machine Owner Key (MOK) facility feature can be used to expand the UEFI Secure Boot key database. When RHEL 9 boots on a UEFI-enabled system with Secure Boot enabled, the keys on the MOK list are also added to the system keyring (.builtin_trusted_keys) in addition to the keys from the key database. The MOK list keys are also stored persistently and securely in the same fashion as the Secure Boot database keys, but these are two separate facilities. The MOK facility is supported by shim.efi, MokManager.efi, grubx64.efi, and the mokutil utility.

Enrolling a MOK key requires manual interaction by a user at the UEFI system console on each target system. Nevertheless, the MOK facility provides a convenient method for testing newly generated key pairs and testing kernel modules signed with them.

Procedure

1. Request the addition of your public key to the MOK list:

```
# mokutil --import my_signing_key_pub.der
```

You will be asked to enter and confirm a password for this MOK enrollment request.

2. Reboot the machine.

The pending MOK key enrollment request will be noticed by shim.efi and it will launch MokManager.efi to allow you to complete the enrollment from the UEFI console.

3. Choose "Enroll MOK" and enter the password you previously associated with this request when prompted and confirm the enrollment.

Your public key is added to the MOK list, which is persistent.

Once a key is on the MOK list, it will be automatically propagated to the system keyring on this and subsequent boots when UEFI Secure Boot is enabled.

NOTE

To facilitate authentication of your kernel module on your systems, consider requesting your system vendor to incorporate your public key into the UEFI Secure Boot key database in their factory firmware image.

3.8. SIGNING KERNEL MODULES WITH THE PRIVATE KEY

Users are able to obtain enhanced security benefits on their systems by loading signed kernel modules if the UEFI Secure Boot mechanism is enabled. The following sections describe how to sign kernel modules with the private key.

Prerequisites
You generated a public and private key pair and know the validity dates of your public keys. For details, see Generating a public and private key pair.

You enrolled your public key on the target system. For details, see Enrolling public key on target system by adding the public key to the MOK list.

You have a kernel module in ELF image format available for signing.

Procedure

- Execute the sign-file utility with parameters as shown in the example below:

```
# /usr/src/kernels/$(uname -r)/scripts/sign-file sha256 my_signing_key.priv my_signing_key_pub.der my_module.ko
```

sign-file computes and appends the signature directly to the ELF image in your kernel module file. The modinfo utility can be used to display information about the kernel module’s signature, if it is present.

NOTE

The appended signature is not contained in an ELF image section and is not a formal part of the ELF image. Therefore, utilities such as readelf will not be able to display the signature on your kernel module.

Your kernel module is now ready for loading. Note that your signed kernel module is also loadable on systems where UEFI Secure Boot is disabled or on a non-UEFI system. That means you do not need to provide both a signed and unsigned version of your kernel module.

IMPORTANT

In RHEL 9, the validity dates of the key pair matter. The key does not expire, but the kernel module must be signed within the validity period of its signing key. The sign-file utility will not warn you of this. For example, a key that is only valid in 2021 can be used to authenticate a kernel module signed in 2021 with that key. However, users cannot use that key to sign a kernel module in 2022.

Additional resources

- Displaying information about kernel modules

3.9. LOADING SIGNED KERNEL MODULES

Once your public key is enrolled in the system keyring (.builtin_trusted_keys) and the MOK list, and after you have signed the respective kernel module with your private key, you can finally load your signed kernel module with the modprobe command as described in the following section.

Prerequisites

- You have generated the public and private keypair. For details, see Generating a public and private key pair.

- You have enrolled the public key into the system keyring. For details, see Enrolling public key on target system by adding the public key to the MOK list.
You have signed a kernel module with the private key. For details, see Signing kernel modules with the private key.

**Procedure**

1. Verify that your public keys are on the system keyring:
   ```
   # keyctl list %:.platform
   ```

2. Install the **kernel-modules-extra** package which will create the `/lib/modules/$(uname -r)/extra/` directory:
   ```
   # dnf -y install kernel-modules-extra
   ```

3. Copy the kernel module into the `/extra/` directory of the kernel you want:
   ```
   # cp my_module.ko /lib/modules/$(uname -r)/extra/
   ```

4. Update the modular dependency list:
   ```
   # depmod -a
   ```

5. Load the kernel module and verify that it was successfully loaded:
   ```
   # modprobe -v my_module
   # lsmod | grep my_module
   ```
   a. Optionally, to load the module on boot, add it to the `/etc/modules-loaded.d/my_module.conf` file:
      ```
      # echo "my_module" > /etc/modules-load.d/my_module.conf
      ```

**Additional resources**

- Managing kernel modules
CHAPTER 4. CONFIGURING KERNEL COMMAND-LINE PARAMETERS

Kernel command-line parameters are a way to change the behavior of certain aspects of the Red Hat Enterprise Linux kernel at boot time. As a system administrator, you have full control over what options get set at boot. Certain kernel behaviors are only able to be set at boot time, so understanding how to make these changes is a key administration skill.

IMPORTANT

Opting to change the behavior of the system by modifying kernel command-line parameters may have negative effects on your system. You should therefore test changes prior to deploying them in production. For further guidance, contact Red Hat Support.

4.1. UNDERSTANDING KERNEL COMMAND-LINE PARAMETERS

Kernel command-line parameters are used for boot time configuration of:

- The Red Hat Enterprise Linux kernel
- The initial RAM disk
- The user space features

Kernel boot time parameters are often used to overwrite default values and for setting specific hardware settings.

By default, the kernel command-line parameters for systems using the GRUB bootloader are defined in the boot entry configuration file for each kernel boot entry.

Additional resources

- kernel-command-line(7), bootparam(7) and dracut.cmdline(7) manual pages
- How to install and boot custom kernels in Red Hat Enterprise Linux 8

4.2. WHAT GRUBBY IS

grubby is a utility for manipulating bootloader-specific configuration files.

You can use grubby also for changing the default boot entry, and for adding/removing arguments from a GRUB menu entry.

Additional resources

- the grubby(8) manual page

4.3. WHAT BOOT ENTRIES ARE

A boot entry is a collection of options which are stored in a configuration file and tied to a particular kernel version. In practice, you have at least as many boot entries as your system has installed kernels. The boot entry configuration file is located in the /boot/loader/entries/ directory and can look like this:
The file name above consists of a machine ID stored in the `/etc/machine-id` file, and a kernel version.

The boot entry configuration file contains information about the kernel version, the initial ramdisk image, and the kernel command-line parameters. The example contents of a boot entry config can be seen below:

```
title Red Hat Enterprise Linux (5.14.0-1.el9.x86_64) 9.0 (Plow)
version 5.14.0-1.el9.x86_64
linux /vmlinuz-5.14.0-1.el9.x86_64
initrd /initramfs-5.14.0-1.el9.x86_64.img
options root=/dev/mapper/rhel_kvm--02--guest08-root ro crashkernel=1G-4G:192M,4G-64G:256M,64G-512M resume=/dev/mapper/rhel_kvm--02--guest08-swap rd.lvm.lv=rhel_kvm-02-guest08/root rd.lvm.lv=rhel_kvm-02-guest08/swap console=ttyS0,115200
grub_users $grub_users
grub_arg --unrestricted
grub_class kernel
```

### Additional resources

- [How to install and boot custom kernels in Red Hat Enterprise Linux 8](#)

---

**4.4. CHANGING KERNEL COMMAND-LINE PARAMETERS FOR ALL BOOt ENTRIES**

This procedure describes how to change kernel command-line parameters for all boot entries on your system.

#### Prerequisites

- Verify that the `grubby` utility is installed on your system.
- Verify that the `zipl` utility is installed on your IBM Z system.

#### Procedure

- To add a parameter:

  ```
  # grubby --update-kernel=ALL --args="<NEW_PARAMETER>"
  ```

  For systems that use the GRUB bootloader and, on IBM Z that use the zIPL bootloader, the command adds a new kernel parameter to each `/boot/loader/entries/<ENTRY>.conf` file.

  On IBM Z that use the zIPL bootloader, the command adds a new kernel parameter to each `/boot/loader/entries/<ENTRY>.conf` file.

  - On IBM Z, execute the `zipl` command with no options to update the boot menu.

- To remove a parameter:

  ```
  # grubby --update-kernel=ALL --remove-args="<PARAMETER_TO_REMOVE>"
  ```
On IBM Z, execute the `zipl` command with no options to update the boot menu.

- After each update of your kernel package, propagate the configured kernel options to the new kernels:
  
  ```
  # grub2-mkconfig -o /etc/grub2.cfg
  ```

**IMPORTANT**

Newly installed kernels do not inherit the kernel command-line parameters from your previously configured kernels. You must run the `grub2-mkconfig` command on the newly installed kernel to propagate the needed parameters to your new kernel.

Additional resources

- Understanding kernel command-line parameters
- `grubby(8)` and `zipl(8)` manual pages
- `grubby tool`

**4.5. CHANGING KERNEL COMMAND-LINE PARAMETERS FOR A SINGLE BOOT ENTRY**

This procedure describes how to change kernel command-line parameters for a single boot entry on your system.

**Prerequisites**

- Verify that the `grubby` and `zipl` utilities are installed on your system.

**Procedure**

- To add a parameter:
  
  ```
  # grubby --update-kernel=/boot/vmlinuz-$\{uname -r\} --args="<NEW_PARAMETER>"
  ```

  - On IBM Z, execute the `zipl` command with no options to update the boot menu.

- To remove a parameter use the following:
  
  ```
  # grubby --update-kernel=/boot/vmlinuz-$\{uname -r\} --remove-args="<PARAMETER_TO_REMOVE>"
  ```

  - On IBM Z, execute the `zipl` command with no options to update the boot menu.

**IMPORTANT**

- `grubby` modifies and stores the kernel command-line parameters of an individual kernel boot entry in the `/boot/loader/entries/<ENTRY>.conf` file.

Additional resources
4.6. CHANGING KERNEL COMMAND-LINE PARAMETERS TEMPORARILY AT BOOT TIME

The following procedure allows you to make temporary changes to a Kernel Menu Entry by changing the kernel parameters only during a single boot process.

Procedure

1. Select the kernel you want to start when the GRUB 2 boot menu appears and press the `e` key to edit the kernel parameters.

2. Find the kernel command line by moving the cursor down. The kernel command line starts with `linux` on 64-Bit IBM Power Series and x86-64 BIOS-based systems, or `linuxefi` on UEFI systems.

3. Move the cursor to the end of the line.

   **NOTE**
   
   Press `Ctrl+a` to jump to the start of the line and `Ctrl+e` to jump to the end of the line. On some systems, `Home` and `End` keys might also work.

4. Edit the kernel parameters as required. For example, to run the system in emergency mode, add the `emergency` parameter at the end of the `linux` line:

   ```
   ```

   + To enable the system messages, remove the `rhgb` and `quiet` parameters.

   1. Press `Ctrl+x` to boot with the selected kernel and the modified command line parameters.

   **IMPORTANT**
   
   Press `Esc` key to leave command line editing and it will drop all the user made changes.

   **NOTE**
   
   This procedure applies only for a single boot and does not persistently make the changes.

4.7. CONFIGURING GRUB SETTINGS TO ENABLE SERIAL CONSOLE CONNECTION
The serial console is beneficial when you need to connect to a headless server or an embedded system and the network is down. Or when you need to avoid security rules and obtain login access on a different system.

You need to configure some default GRUB settings to use the serial console connection.

**Prerequisites**

- You have root permissions.

**Procedure**

1. Add the following two lines to the `/etc/default/grub` file:

   ```
   GRUB_TERMINAL="serial"
   GRUB_SERIAL_COMMAND="serial --speed=9600 --unit=0 --word=8 --parity=no --stop=1"
   ```

   The first line disables the graphical terminal. The `GRUB_TERMINAL` key overrides values of `GRUB_TERMINAL_INPUT` and `GRUB_TERMINAL_OUTPUT` keys.

   The second line adjusts the baud rate (`--speed`), parity and other values to fit your environment and hardware. Note that a much higher baud rate, for example 115200, is preferable for tasks such as following log files.

2. Update the GRUB configuration file.

   - On BIOS-based machines:

     ```
     # grub2-mkconfig -o /boot/grub2/grub.cfg
     ```

   - On UEFI-based machines:

     ```
     # grub2-mkconfig -o /boot/grub2/grub.cfg
     ```

3. Reboot the system for the changes to take effect.
CHAPTER 5. CONFIGURING KERNEL PARAMETERS AT RUNTIME

As a system administrator, you can modify many facets of the Red Hat Enterprise Linux kernel’s behavior at runtime. This section describes how to configure kernel parameters at runtime by using the `sysctl` command and by modifying the configuration files in the `/etc/sysctl.d/` and `/proc/sys/` directories.

5.1. WHAT ARE KERNEL PARAMETERS

Kernel parameters are tunable values which you can adjust while the system is running. There is no requirement to reboot or recompile the kernel for changes to take effect.

It is possible to address the kernel parameters through:

- The `sysctl` command
- The virtual file system mounted at the `/proc/sys/` directory
- The configuration files in the `/etc/sysctl.d/` directory

Tunables are divided into classes by the kernel subsystem. Red Hat Enterprise Linux has the following tunable classes:

**Table 5.1. Table of sysctl classes**

<table>
<thead>
<tr>
<th>Tunable class</th>
<th>Subsystem</th>
</tr>
</thead>
<tbody>
<tr>
<td>abi</td>
<td>Execution domains and personalities</td>
</tr>
<tr>
<td>crypto</td>
<td>Cryptographic interfaces</td>
</tr>
<tr>
<td>debug</td>
<td>Kernel debugging interfaces</td>
</tr>
<tr>
<td>dev</td>
<td>Device-specific information</td>
</tr>
<tr>
<td>fs</td>
<td>Global and specific file system tunables</td>
</tr>
<tr>
<td>kernel</td>
<td>Global kernel tunables</td>
</tr>
<tr>
<td>net</td>
<td>Network tunables</td>
</tr>
<tr>
<td>sunrpc</td>
<td>Sun Remote Procedure Call (NFS)</td>
</tr>
<tr>
<td>user</td>
<td>User Namespace limits</td>
</tr>
<tr>
<td>vm</td>
<td>Tuning and management of memory, buffers, and cache</td>
</tr>
</tbody>
</table>

**IMPORTANT**

Red Hat Enterprise Linux 9 Managing, monitoring, and updating the kernel
IMPORTANT

Configuring kernel parameters on a production system requires careful planning. Unplanned changes may render the kernel unstable, requiring a system reboot. Verify that you are using valid options before changing any kernel values.

Additional resources

- `sysctl(8)`, and `sysctl.d(5)` manual pages

5.2. CONFIGURING KERNEL PARAMETERS TEMPORARILY WITH SYSCTL

The following procedure describes how to use the `sysctl` command to temporarily set kernel parameters at runtime. The command is also useful for listing and filtering tunables.

Prerequisites

- Root permissions

Procedure

1. To list all parameters and their values, use the following:

   ```
   # sysctl -a
   ```

   **NOTE**
   
   The `sysctl -a` command displays kernel parameters, which can be adjusted at runtime and at boot time.

2. To configure a parameter temporarily, use the command as in the following example:

   ```
   # sysctl <TUNABLE_CLASS>.<PARAMETER>=<TARGET_VALUE>
   ```

   The sample command above changes the parameter value while the system is running. The changes take effect immediately, without a need for restart.

   **NOTE**
   
   The changes return back to default after your system reboots.

Additional resources

- `sysctl(8)` manual page
- Configuring kernel parameters permanently with `sysctl`
- Using configuration files in `/etc/sysctl.d/` to adjust kernel parameters

5.3. CONFIGURING KERNEL PARAMETERS PERMANENTLY WITH SYSCTL
The following procedure describes how to use the `sysctl` command to permanently set kernel parameters.

**Prerequisites**

- Root permissions

**Procedure**

1. To list all parameters, use the following:
   
   ```shell
   # sysctl -a
   ```
   
   The command displays all kernel parameters that can be configured at runtime.

2. To configure a parameter permanently:
   
   ```shell
   # sysctl -w <TUNABLE_CLASS>.<PARAMETER>=<TARGET_VALUE> >> /etc/sysctl.conf
   ```
   
   The sample command changes the tunable value and writes it to the `/etc/sysctl.conf` file, which overrides the default values of kernel parameters. The changes take effect immediately and persistently, without a need for restart.

   **NOTE**
   
   To permanently modify kernel parameters you can also make manual changes to the configuration files in the `/etc/sysctl.d/` directory.

**Additional resources**

- `sysctl(8)` and `sysctl.conf(5)` manual pages
- Using configuration files in `/etc/sysctl.d/` to adjust kernel parameters

### 5.4. USING CONFIGURATION FILES IN `/ETC/SYSCTL.D/` TO ADJUST KERNEL PARAMETERS

The following procedure describes how to manually modify configuration files in the `/etc/sysctl.d/` directory to permanently set kernel parameters.

**Prerequisites**

- Root permissions

**Procedure**

1. Create a new configuration file in `/etc/sysctl.d/`:
   
   ```shell
   # vim /etc/sysctl.d/<some_file.conf>
   ```

2. Include kernel parameters, one per line, as follows:
3. Save the configuration file.
4. Reboot the machine for the changes to take effect.
   - Alternatively, to apply changes without rebooting, execute:

```
# sysctl -p /etc/sysctl.d/<some_file.conf>
```

The command enables you to read values from the configuration file, which you created earlier.

Additional resources
- `sysctl(8)`, `sysctl.d(5)` manual pages

### 5.5. CONFIGURING KERNEL PARAMETERS TEMPORARILY THROUGH /PROC/SYS/

The following procedure describes how to set kernel parameters temporarily through the files in the virtual file system `/proc/sys/` directory.

Prerequisites
- Root permissions

Procedure
1. Identify a kernel parameter you want to configure:
   ```
   # ls -l /proc/sys/<TUNABLE_CLASS>/
   ```
   The writable files returned by the command can be used to configure the kernel. The files with read-only permissions provide feedback on the current settings.

2. Assign a target value to the kernel parameter:
   ```
   # echo <TARGET_VALUE> > /proc/sys/<TUNABLE_CLASS>/<PARAMETER>
   ```
   The command makes configuration changes that will disappear once the system is restarted.

3. Optionally, verify the value of the newly set kernel parameter:
   ```
   # cat /proc/sys/<TUNABLE_CLASS>/<PARAMETER>
   ```

Additional resources
- Configuring kernel parameters permanently with `sysctl`
- Using configuration files in `/etc/sysctl.d/` to adjust kernel parameters
CHAPTER 6. KEEPING KERNEL PANIC PARAMETERS DISABLED IN VIRTUALIZED ENVIRONMENTS

When configuring a virtualized environment in RHEL 9, you should not enable the `softlockup_panic` and `nmi_watchdog` kernel parameters, as the virtualized environment may trigger a spurious soft lockup that should not require a system panic.

The following sections explain the reasons behind this advice by summarizing:

- What causes a soft lockup.
- Describing the kernel parameters that control a system’s behavior on a soft lockup.
- Explaining how soft lockups may be triggered in a virtualized environment.

6.1. WHAT IS A SOFT LOCKUP

A soft lockup is a situation usually caused by a bug, when a task is executing in kernel space on a CPU without rescheduling. The task also does not allow any other task to execute on that particular CPU. As a result, a warning is displayed to a user through the system console. This problem is also referred to as the soft lockup firing.

Additional resources

- *What is a CPU soft lockup?*

6.2. PARAMETERS CONTROLLING KERNEL PANIC

The following kernel parameters can be set to control a system’s behavior when a soft lockup is detected.

`softlockup_panic`

Controls whether or not the kernel will panic when a soft lockup is detected.

<table>
<thead>
<tr>
<th>Type</th>
<th>Value</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>Integer</td>
<td>0</td>
<td>kernel does not panic on soft lockup</td>
</tr>
<tr>
<td>Integer</td>
<td>1</td>
<td>kernel panics on soft lockup</td>
</tr>
</tbody>
</table>

By default, on RHEL8 this value is 0.

In order to panic, the system needs to detect a hard lockup first. The detection is controlled by the `nmi_watchdog` parameter.

`nmi_watchdog`

Controls whether lockup detection mechanisms (watchdogs) are active or not. This parameter is of integer type.
<table>
<thead>
<tr>
<th>Value</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>disables lockup detector</td>
</tr>
<tr>
<td>1</td>
<td>enables lockup detector</td>
</tr>
</tbody>
</table>

The hard lockup detector monitors each CPU for its ability to respond to interrupts.

`watchdog_thresh`

Controls frequency of watchdog `hrtimer`, NMI events, and soft/hard lockup thresholds.

<table>
<thead>
<tr>
<th>Default threshold</th>
<th>Soft lockup threshold</th>
</tr>
</thead>
<tbody>
<tr>
<td>10 seconds</td>
<td>$2 \times \text{watchdog_thresh}$</td>
</tr>
</tbody>
</table>

Setting this parameter to zero disables lockup detection altogether.

**Additional resources**

- [Softlockup detector and hardlockup detector](#)
- [Kernel sysctl](#)

## 6.3. SPURIOUS SOFT LOCKUPS IN VIRTUALIZED ENVIRONMENTS

The soft lockup firing on physical hosts, as described in [What is a soft lockup](#), usually represents a kernel or hardware bug. The same phenomenon happening on guest operating systems in virtualized environments may represent a false warning.

Heavy work-load on a host or high contention over some specific resource such as memory, usually causes a spurious soft lockup firing. This is because the host may schedule out the guest CPU for a period longer than 20 seconds. Then when the guest CPU is again scheduled to run on the host, it experiences a **time jump** which triggers due timers. The timers include also watchdog `hrtimer`, which can consequently report a soft lockup on the guest CPU.

Because a soft lockup in a virtualized environment may be spurious, you should not enable the kernel parameters that would cause a system panic when a soft lockup is reported on a guest CPU.

### IMPORTANT

To understand soft lockups in guests, it is essential to know that the host schedules the guest as a task, and the guest then schedules its own tasks.

**Additional resources**

- [What is a soft lockup](#)
- [Virtual machine components and their interaction](#)
- Virtual machine reports a "BUG: soft lockup"
CHAPTER 7. ADJUSTING KERNEL PARAMETERS FOR DATABASE SERVERS

There are different sets of kernel parameters which can affect performance of specific database applications. The following sections explain what kernel parameters to configure to secure efficient operation of database servers and databases.

7.1. INTRODUCTION TO DATABASE SERVERS

A database server is a service that provides features of a database management system (DBMS). DBMS provides utilities for database administration and interacts with end users, applications, and databases.

Red Hat Enterprise Linux 9 provides the following database management systems:

- MariaDB 10.5
- MySQL 8.0
- PostgreSQL 13
- Redis 6

7.2. PARAMETERS AFFECTING PERFORMANCE OF DATABASE APPLICATIONS

The following kernel parameters affect performance of database applications.

**fs.aio-max-nr**

Defines the maximum number of asynchronous I/O operations the system can handle on the server.

**NOTE**

Raising the `fs.aio-max-nr` parameter produces no additional changes beyond increasing the aio limit.

**fs.file-max**

Defines the maximum number of file handles (temporary file names or IDs assigned to open files) the system supports at any instance.

The kernel dynamically allocates file handles whenever a file handle is requested by an application. The kernel however does not free these file handles when they are released by the application. The kernel recycles these file handles instead. This means that over time the total number of allocated file handles will increase even though the number of currently used file handles may be low.

**kernel.shmall**

Defines the total number of shared memory pages that can be used system-wide. To use the entire main memory, the value of the `kernel.shmall` parameter should be $\leq$ total main memory size.

**kernel.shmmmax**

Defines the maximum size in bytes of a single shared memory segment that a Linux process can allocate in its virtual address space.

**kernel.shmmni**
Defines the maximum number of shared memory segments the database server is able to handle.

**net.ipv4.ip_local_port_range**
Defines the port range the system can use for programs which want to connect to a database server without a specific port number.

**net.core.rmem_default**
Defines the default receive socket memory through Transmission Control Protocol (TCP).

**net.core.rmem_max**
Defines the maximum receive socket memory through Transmission Control Protocol (TCP).

**net.core.wmem_default**
Defines the default send socket memory through Transmission Control Protocol (TCP).

**net.core.wmem_max**
Defines the maximum send socket memory through Transmission Control Protocol (TCP).

**vm.dirty_bytes / vm.dirty_ratio**
Defines a threshold in bytes / in percentage of dirty-able memory at which a process generating dirty data is started in the `write()` function.

**NOTE**
Either `vm.dirty_bytes` or `vm.dirty_ratio` can be specified at a time.

**vm.dirty_background_bytes / vm.dirty_background_ratio**
Defines a threshold in bytes / in percentage of dirty-able memory at which the kernel tries to actively write dirty data to hard-disk.

**NOTE**
Either `vm.dirty_background_bytes` or `vm.dirty_background_ratio` can be specified at a time.

**vm.dirty_writeback_centisecs**
Defines a time interval between periodic wake-ups of the kernel threads responsible for writing dirty data to hard-disk.
This kernel parameters measures in 100th’s of a second.

**vm.dirty_expire_centisecs**
Defines the time after which dirty data is old enough to be written to hard-disk.
This kernel parameters measures in 100th’s of a second.

**Additional resources**
- *Dirty pagecache writeback and vm.dirty parameters*
CHAPTER 8. GETTING STARTED WITH KERNEL LOGGING

Log files are files that contain messages about the system, including the kernel, services, and applications running on it. The logging system in Red Hat Enterprise Linux is based on the built-in syslog protocol. Various utilities use this system to record events and organize them into log files. These files are useful when auditing the operating system or troubleshooting problems.

8.1. WHAT IS THE KERNEL RING BUFFER

During the boot process, the console provides a lot of important information about the initial phase of the system startup. To avoid loss of the early messages the kernel utilizes what is called a ring buffer. This buffer stores all messages, including boot messages, generated by the printk() function within the kernel code. The messages from the kernel ring buffer are then read and stored in log files on permanent storage, for example, by the syslog service.

The buffer mentioned above is a cyclic data structure which has a fixed size, and is hard-coded into the kernel. Users can display data stored in the kernel ring buffer through the dmesg command or the /var/log/boot.log file. When the ring buffer is full, the new data overwrites the old.

Additional resources

- syslog(2) and dmesg(1) manual page

8.2. ROLE OF PRINTK ON LOG-LEVELS AND KERNEL LOGGING

Each message the kernel reports has a log-level associated with it that defines the importance of the message. The kernel ring buffer, as described in What is the kernel ring buffer, collects kernel messages of all log-levels. It is the kernel.printk parameter that defines what messages from the buffer are printed to the console.

The log-level values break down in this order:

- 0 – Kernel emergency. The system is unusable.
- 1 – Kernel alert. Action must be taken immediately.
- 2 – Condition of the kernel is considered critical.
- 3 – General kernel error condition.
- 4 – General kernel warning condition.
- 5 – Kernel notice of a normal but significant condition.
- 6 – Kernel informational message.
- 7 – Kernel debug-level messages.

By default, kernel.printk in RHEL 9 contains the following four values:

```
# sysctl kernel.printk
kernel.printk = 7 4 1 7
```

The four values define the following:
1. value. Console log-level, defines the lowest priority of messages printed to the console.

2. value. Default log-level for messages without an explicit log-level attached to them.

3. value. Sets the lowest possible log-level configuration for the console log-level.

4. value. Sets default value for the console log-level at boot time.
   Each of these values above defines a different rule for handling error messages.

**IMPORTANT**

The default \texttt{7 4 1 7 printk} value allows for better debugging of kernel activity. However, when coupled with a serial console, this \texttt{printk} setting is able to cause intense I/O bursts that could lead to a RHEL system becoming temporarily unresponsive. To avoid these situations, setting a \texttt{printk} value of \texttt{4 4 1 7} typically works, but at the expense of losing the extra debugging information.

Also note that certain kernel command line parameters, such as \texttt{quiet} or \texttt{debug}, change the default \texttt{kernel.printk} values.

**Additional resources**

- \texttt{syslog(2)} manual page
CHAPTER 9. INSTALLING KDUMP

The **kdump** service is installed and activated by default on the new Red Hat Enterprise Linux installations. The following sections explain what **kdump** is and how to install **kdump** when it is not enabled by default.

### 9.1. WHAT IS KDUMP

**kdump** is a service which provides a crash dumping mechanism. The service enables you to save the contents of the system memory for analysis. **kdump** uses the **kexec** system call to boot into the second kernel (a *capture kernel*) without rebooting; and then captures the contents of the crashed kernel’s memory (a *crash dump* or a *vmcore*) and saves it into a file. The second kernel resides in a reserved part of the system memory.

**IMPORTANT**

A kernel crash dump can be the only information available in the event of a system failure (a critical bug). Therefore, operational **kdump** is important in mission-critical environments. Red Hat advise that system administrators regularly update and test **kexec-tools** in your normal kernel update cycle. This is especially important when new kernel features are implemented.

You can enable **kdump** for all installed kernels on a machine or only for specified kernels. This is useful when there are multiple kernels used on a machine, some of which are stable enough that there is no concern that they could crash.

When **kdump** is installed, a default /etc/kdump.conf file is created. The file includes the default minimum **kdump** configuration. You can edit this file to customize the **kdump** configuration, but it is not required.

### 9.2. INSTALLING KDUMP USING ANACONDA

The **Anaconda** installer provides a graphical interface screen for **kdump** configuration during an interactive installation. The installer screen is titled as **KDUMP** and is available from the main **Installation Summary** screen. You can enable **kdump** and reserve the required amount of memory.

**Procedure**

1. Go to the **Kdump** field.
2. Enable **kdump** if not already enabled.
3. Define how much memory should be reserved for `kdump`.

<table>
<thead>
<tr>
<th>KDUMP</th>
<th>RED HAT ENTERPRISE LINUX 9.0 INSTALLATION</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Kdump is a kernel crash dumping mechanism. In the event of a system crash, kdump will capture information from your system that can be invaluable in determining the cause of the crash. Note that kdump does require reserving a portion of system memory that will be unavailable for other uses.

- Enable kdump
- Kdump Memory Reservations: Automatic
- Memory To Be Reserved (MB): 160
- Total System Memory (MB): 2971
- Unused System Memory (MB): 2811

### 9.3. INSTALLING KDUMP ON THE COMMAND LINE

Some installation options, such as custom Kickstart installations, in some cases do not install or enable `kdump` by default. If this is your case, follow the procedure below.

#### Prerequisites
- An active RHEL subscription
- The `kexec-tools` package
- Fulfilled requirements for `kdump` configurations and targets. For details, see [Supported kdump configurations and targets](#).

#### Procedure

1. Check whether `kdump` is installed on your system:

   ```bash
   # rpm -q kexec-tools
   ```

   Output if the package is installed:

   ```bash
   # kexec-tools-2.0.22-13.el9.x86_64
   ```

   Output if the package is not installed:

   ```bash
   package kexec-tools is not installed
   ```

2. Install `kdump` and other necessary packages by:

   ```bash
   # dnf install kexec-tools
   ```
CHAPTER 10. CONFIGURING KDUMP ON THE COMMAND LINE

The memory for **kdump** is reserved during the system boot. The memory size is configured in the system’s Grand Unified Bootloader (GRUB) 2 configuration file. The memory size depends on the `crashkernel=` value specified in the configuration file and the size of the system’s physical memory.

10.1. CONFIGURING KDUMP MEMORY USAGE

The **kexec-tools** package maintains the default **crashkernel** memory reservation values. The **kdump** service uses the default value to reserve the **crashkernel** memory for each kernel.

The automatic memory allocation for **kdump** varies based on the system hardware architecture and available memory size. For example, on AMD and Intel 64-bit architectures, the **crashkernel** default parameters work only when the available memory is more than 1 GB. By default, **kexec-tools** configures the following memory reserves on AMD64 and Intel 64-bit architectures:

```
crashkernel=1G-4G:192M,4G-64G:256M,64G-:512M
```

The memory requirement of the crash kernel may vary depending on the hardware and machine specifications. If the default `crashkernel` value does not work on your system, you can run the **kdumpctl estimate** command and query a rough estimate value without triggering a crash. The estimated `crashkernel` value may not be accurate and can serve as a reference to set an appropriate `crashkernel` value.

**NOTE**

The `crashkernel=auto` option in the boot command line is no longer supported on RHEL 9 and later releases.

**Prerequisites**

- Root privileges
- Fulfilled **kdump** requirements for configurations and targets.
- On IBM Z systems, ensure the **zipl** utility is installed.

**Procedure**

1. Configure the default value for `crashkernel`:

   ```
   # kdumpctl reset-crashkernel --kernel=ALL
   ```

2. (Optional) To use a custom `crashkernel` value:
   a. Configure the required memory reserve:

   ```
   crashkernel=1G-4G:192M,4G-:256M
   ```

   The example reserves 192 MB of memory if the total amount of system memory is 1 GB or higher and lower than 4 GB. If the total amount of memory is more than 4 GB, 256 MB is reserved for **kdump** instead.
Some systems require to reserve memory with a certain fixed offset since crash kernel reservation is very early, and it wants to reserve some area for special usage. If the offset is set, the reserved memory begins there. To offset the reserved memory, use the following syntax:

```
crashkernel=192M@16M
```

The example above reserves 192 MB of memory starting at 16 MB (physical address 0x01000000). If the offset parameter is set to 0 or omitted entirely, `kdump` offsets the reserved memory automatically. You can also offset memory when setting a variable memory reservation by specifying the offset as the last value. For example,

```
crashkernel=1G-4G:192M,2G-64G:256M@16M
```

b. Update the bootloader configuration:

```
# grubby --update-kernel ALL --args "crashkernel=<CUSTOM-VALUE>"
```

On IBM Z systems that use the zIPL bootloader, the command adds a new kernel parameter to each `/boot/loader/entries/<ENTRY>.conf` file.

- On IBM Z systems, to update the boot menu, execute the `zipl` command with no options specified:

```
# zipl
```

3. Reboot for changes to take effect:

```
# reboot
```

**Verification**

1. Activate the `sysrq` key to boot into the `kdump` kernel:

```
# echo 1 > /proc/sys/kernel/sysrq
# echo c > /proc/sysrq-trigger
```

This forces the Linux kernel to crash and copy the `address-YYYY-MM-DD-HH:MM:SS/vmcore` file to the target location specified in the configuration file.

2. Verify that the `vmcore` file is dumped in the target as specified in the `/etc/kdump.conf` file:

```
$ ls /var/crash/127.0.0.1-2022-01-18-0
/var/crash/127.0.0.1-2022-01-18-05:23:10:
  kexec-dmesg.log  vmcore  vmcore-dmesg.txt
```

In this example, the kernel saves the `vmcore` in the default target directory, `/var/crash/`.

**10.2. CONFIGURING THE KDUMP TARGET**

When a kernel crash is captured, the core dump can be either stored as a file in a local file system, written directly to a device, or sent over a network using the `NFS` (Network File System) or `SSH` (Secure
Shell) protocol. Only one of these options can be set at a time, and the default behavior is to store the vmcore file in the /var/crash/ directory of the local file system.

Prerequisites

- Fulfilled **kdump** requirements for configurations and targets.

Procedure

- To store the **vmcore** file in /var/crash/ directory of the local file system, edit the /etc/kdump.conf file and specify the path:

  ```
  # path /var/crash
  ```

  The option **path /var/crash** represents the path to the file system in which **kdump** saves the **vmcore** file. When you specify a dump target in the /etc/kdump.conf file, then the **path** is relative to the specified dump target.

  If you do not specify a dump target in the /etc/kdump.conf file, then the **path** represents the absolute path from the root directory. Depending on what is mounted in the current system, the dump target and the adjusted dump path are taken automatically.

  **WARNING**

  **kdump** saves the **vmcore** file in /var/crash/var/crash directory, when the dump target is mounted at /var/crash and the option **path** is also set as /var/crash in the /etc/kdump.conf file. For example, in the following instance, the **ext4** file system is already mounted at /var/crash and the **path** are set as /var/crash:

  ```
  # grep -v ^# /etc/kdump.conf | grep -v ^$
  ext4 /dev/mapper/vg00-varcrashvol
  path /var/crash
  core_collector makedumpfile -c --message-level 1 -d 31
  ```

  This results in the /var/crash/var/crash path. To solve this problem, use the option **path /** instead of **path /var/crash**

- To change the local directory in which the core dump is to be saved, as **root**, edit the /etc/kdump.conf configuration file as described below.
  1. Remove the hash sign (“#”) from the beginning of the **#path /var/crash** line.
  2. Replace the value with the intended directory path. For example:

     ```
     path /usr/local/cores
     ```

In RHEL 9, the directory defined as the kdump target using the `path` directive must exist when the `kdump` systemd service is started - otherwise the service fails. This behavior is different from earlier releases of RHEL, where the directory was being created automatically if it did not exist when starting the service.

- To write the file to a different partition, as `root`, edit the `/etc/kdump.conf` configuration file as described below.

  1. Remove the hash sign ("#") from the beginning of the `#ext4` line, depending on your choice.
     - device name (the `#ext4 /dev/vg/lv_kdump` line)
     - file system label (the `#ext4 LABEL=/boot` line)
     - UUID (the `#ext4 UUID=03138356-5e61-4ab3-b58e-27507ac41937` line)
  2. Change the file system type as well as the device name, label or UUID to the desired values. For example:

     ```
     ext4 UUID=03138356-5e61-4ab3-b58e-27507ac41937
     ```

- It is recommended to specify storage devices using a `LABEL=` or `UUID=`. Disk device names such as `/dev/sda3` are not guaranteed to be consistent across reboot.

- When dumping to Direct Access Storage Device (DASD) on IBM Z hardware, it is essential that the dump devices are correctly specified in `/etc/dasd.conf` before proceeding.

- To write the dump directly to a device:

  1. Remove the hash sign ("#") from the beginning of the `#raw /dev/vg/lv_kdump` line.
  2. Replace the value with the intended device name. For example:

     ```
     raw /dev/sdb1
     ```

- To store the dump to a remote machine using the `NFS` protocol:

  1. Remove the hash sign ("#") from the beginning of the `#nfs my.server.com:/export/tmp` line.
  2. Replace the value with a valid hostname and directory path. For example:

     ```
     nfs penguin.example.com:/export/cores
     ```

- To store the dump to a remote machine using the `SSH` protocol:
1. Remove the hash sign ("#") from the beginning of the `#ssh user@my.server.com` line.

2. Replace the value with a valid username and hostname.

3. Include your SSH key in the configuration.
   - Remove the hash sign from the beginning of the `#sshkey /root/.ssh/kdump_id_rsa` line.
   - Change the value to the location of a key valid on the server you are trying to dump to.
     For example:
     ```
     ssh john@penguin.example.com
     sshkey /root/.ssh/mykey
     ```

10.3. CONFIGURING THE CORE COLLECTOR

The `kdump` service uses the `core_collector` program to capture the `vmcore` image. In RHEL, the `makedumpfile` utility is the default core collector.

`makedumpfile` is a dump program that helps to copy only the required pages using various dump levels and compress the size of a dump file.

Using `makedumpfile`, you can create a small size dump file either by compressing dump data or by excluding pages or both. It needs the first kernel debug information to distinguish the not necessary pages by analyzing how the first kernel uses the memory.

**Syntax**

```bash
core_collector makedumpfile -z -d 31 --message-level 1
```

**Options**

- `-c`, `-l`, `-z`, or `-p`: specifies the compress dump file format by each page when you use one of these options: `-c` for zlib, `-l` for lzo, `-z` for zstd, or `-p` for snappy.

- `-d (dump_level)`: excludes pages so that they do not get copied to the dump file.

- `--message-level`: specifies the message types.
  Using `--message-level`, you can restrict the outputs to print. For example, specifying 7 as the message level prints common messages and error messages. The maximum value for `--message-level` is 31.

**Prerequisites**

- Fulfilled `kdump` requirements for configurations and targets.

**Procedure**

1. As `root` user, edit the `/etc/kdump.conf` configuration file to remove the hash sign ("#") from the beginning of the following command:
core_collector makedumpfile -z -d 31 --message-level 1

2. To enable dump file compression, specify one of the makedumpfile options:

   core_collector makedumpfile -z -d 31 --message-level 1

where,

- `-z` specifies the dump compressed file format.
- `-d` specifies dump level as 31.
- `--message-level` specifies message level as 1.

Also, consider the following example that uses `-l`:

- To compress a dump file using `-l`:
  
  core_collector makedumpfile -l -d 31 --message-level 1

Additional resources

- makedumpfile(8) manual page

10.4. CONFIGURING THE KDUMP DEFAULT FAILURE RESPONSES

By default, when kdump fails to create a crash dump file at the configured target location, the system reboots and the dump is lost in the process. To change this behavior, follow the procedure below.

Prerequisites

- Root permissions.
- Fulfilled requirements for kdump configurations and targets.

Procedure

1. As root, remove the hash sign ("#") from the beginning of the `#failure_action` line in the `/etc/kdump.conf` configuration file.

2. Replace the value with a desired action.

   failure_action poweroff

10.5. CONFIGURATION FILE FOR KDUMP

The configuration file for kdump kernel is `/etc/sysconfig/kdump`. This file controls the kdump kernel command line parameters.

For most configurations, use the default options. However, in some scenarios you might need to modify certain parameters to control the kdump kernel behavior. For example, modifying to append the kdump kernel command-line to obtain a detailed debugging output.
This section covers information on modifying the KDUMP_COMMANDLINE_REMOVE and KDUMP_COMMANDLINE_APPEND options for kdump. For information on additional configuration options refer to Documentation/admin-guide/kernel-parameters.txt or the /etc/sysconfig/kdump file.

- **KDUMP_COMMANDLINE_REMOVE**
  This option removes arguments from the current kdump command line. It removes parameters that may cause kdump errors or kdump kernel boot failures. These parameters may have been parsed from the previous KDUMP_COMMANDLINE process or inherited from the /proc/cmdline file. When this variable is not configured, it inherits all values from the /proc/cmdline file. Configuring this option also provides information that is helpful in debugging an issue.

  **Example**

  To remove certain arguments, add them to KDUMP_COMMANDLINE_REMOVE as follows:

  ```
  # KDUMP_COMMANDLINE_REMOVE="hugepages hugepagesize slub_debug quiet log_buf_len swiotlb"
  ```

- **KDUMP_COMMANDLINE_APPEND**
  This option appends arguments to the current command line. These arguments may have been parsed by the previous KDUMP_COMMANDLINE_REMOVE variable.

  For the kdump kernel, disabling certain modules such as mce, cgroup, numa, hest_disable can help prevent kernel errors. These modules may consume a significant portion of the kernel memory reserved for kdump or cause kdump kernel boot failures.

  **Example**

  To disable memory cgroups on the kdump kernel command line, run the command as follows:

  ```
  # KDUMP_COMMANDLINE_APPEND="cgroup_disable=memory"
  ```

**Additional resources**

- Documentation/admin-guide/kernel-parameters.txt file
- /etc/sysconfig/kdump file

**10.6. ENABLING AND DISABLING THE KDUMP SERVICE**

To start the kdump service at boot time, follow the procedure below.

**Prerequisites**

- Fulfilled kdump requirements for configurations and targets.
- All configurations for installing kdump are set up according to your needs.

**Procedure**

1. To enable the kdump service, use the following command:
# systemctl enable kdump.service
This enables the service for `multi-user.target`.

2. To start the service in the current session, use the following command:

```bash
# systemctl start kdump.service
```

3. To stop the `kdump` service, type the following command:

```bash
# systemctl stop kdump.service
```

4. To disable the `kdump` service, execute the following command:

```bash
# systemctl disable kdump.service
```

---

**WARNING**

It is recommended to set `kptr_restrict=1` as default. When `kptr_restrict` is set to (1) as default, the `kdumpctl` service loads the crash kernel even if Kernel Address Space Layout (KASLR) is enabled or not enabled.

---

**Troubleshooting step**

When `kptr_restrict` is not set to (1), and if KASLR is enabled, the contents of `/proc/kore` file are generated as all zeros. Consequently, the `kdumpctl` service fails to access the `/proc/kcore` and load the crash kernel.

To work around this problem, the `kexec-kdump-howto.txt` file displays a warning message, which specifies to keep the recommended setting as `kptr_restrict=1`.

To ensure that `kdumpctl` service loads the crash kernel, verify that:

- Kernel `kptr_restrict=1` in the `sysctl.conf` file.

---

### 10.7. TESTING THE KDUMP CONFIGURATION

You can test that the crash dump process works and is valid before the machine enters production.

---

**WARNING**

The commands below cause the kernel to crash. Use caution when following these steps, and never carelessly use them on active production system.
Procedure

1. Reboot the system with **kdump** enabled.

2. Make sure that **kdump** is running:

   ```bash
   # systemctl is-active kdump
   active
   ```

3. Force the Linux kernel to crash:

   ```bash
   echo 1 > /proc/sys/kernel/sysrq
   echo c > /proc/sysrq-triger
   ```

   **WARNING**

   The command above crashes the kernel, and a reboot is required.

   Once booted again, the *address*-YYYY-MM-DD-HH:MM:SS/vmcore file is created at the location you have specified in the `/etc/kdump.conf` file (by default to `/var/crash`).

   **NOTE**

   This action confirms the validity of the configuration. Also it is possible to use this action to record how long it takes for a crash dump to complete with a representative work-load.

Additional resources

- Configuring the kdump target

### 10.8. PREVENTING KERNEL DRIVERS FROM LOADING FOR KDUMP

This section explains how to prevent the capture kernel from loading certain kernel drivers using the `/etc/sysconfig/kdump` configuration file. You can prevent the **kdump** initramfs from loading the specified kernel module. To achieve this, you need to put the `KDUMP_COMMANDLINE_APPEND=` variable in the `/etc/sysconfig/kdump` file. This helps to prevent the out-of-memory (oom) killer or other crash kernel failures.

You can append the `KDUMP_COMMANDLINE_APPEND=` variable using one of the following configuration options:

- `rd.driver.blacklist=<modules>`
- `modprobe.blacklist=<modules>

**Procedure**

1. Select a kernel module that you intend to block from loading.
The `lsmod` command displays a list of modules that are loaded to the currently running kernel.

2. Update the `KDUMP_COMMANDLINE_APPEND=` variable in the `/etc/sysconfig/kdump` file.

   ```
   #
   KDUMP_COMMANDLINE_APPEND="rd.driver.blacklist=hw_vmbus,hv_storvsc,hv_utils,
   hv_netvsc,hid-hyperv"
   ```

   Also, consider the following example using the `modprobe.blacklist=<modules>` configuration option.

   ```
   # KDUMP_COMMANDLINE_APPEND="modprobe.blacklist=emcp
   modprobe.blacklist=bnx2fc modprobe.blacklist=libfcoe modprobe.blacklist=fcoe"
   ```

3. Restart the `kdump` service.

   ```
   # systemctl restart kdump
   ```

**Additional resources**
- [dracut.cmdline](#) manual page

### 10.9. RUNNING KDUMP ON SYSTEMS WITH ENCRYPTED DISK

When you run a Linux Unified Key Setup (LUKS) encrypted partition, the system requires certain amount of available memory. If the system has less than the required amount of available memory, the `systemd-cryptsetup` service fails to mount the partition. As a result, capturing the `vmcore` file to an encrypted target location fails in the second kernel (capture kernel).

The `kdumpctl estimate` command helps you estimate the amount of memory you need for `kdump`. It prints the recommended `crashkernel` value, which is the most suitable memory size required for `kdump`.

The recommended `crashkernel` value is calculated based on the current kernel size, kernel modules, `initramfs`, and the LUKS encrypted target memory requirement.

In case you use the custom `crashkernel` option, `kdumpctl estimate` prints the LUKS required size value. The value is the memory size required for LUKS encrypted target.

**Procedure**

1. Print the estimate `crashkernel` value:

   ```
   # kdumpctl estimate
   ```
Encrypted kdump target requires extra memory, assuming using the keyslot with minimum memory requirement
- Reserved crashkernel: 256M
- Recommended crashkernel: 652M
- Kernel image size: 47M
- Kernel modules size: 8M
- Initramfs size: 20M
- Runtime reservation: 64M
- LUKS required size: 512M
- Large modules: none

WARNING: Current crashkernel size is lower than recommended size 652M.

2. Configure the amount of required memory by increasing `crashkernel` to the desired value.

   ```
   # grubby -args="crashkernel=652M" --update-kernel=ALL
   ```

3. Reboot for changes to take effect.

   ```
   # reboot
   ```

NOTE

If the `kdump` service still fails to save the dump file to the encrypted target, increase the `crashkernel` value gradually to configure an appropriate amount of memory.
CHAPTER 11. SUPPORTED KDUMP CONFIGURATIONS AND TARGETS

11.1. MEMORY REQUIREMENTS FOR KDUMP

In order for kdump to be able to capture a kernel crash dump and save it for further analysis, a part of the system memory has to be permanently reserved for the capture kernel. When reserved, this part of the system memory is not available to the main kernel.

The memory requirements vary based on certain system parameters. One of the major factors is the system’s hardware architecture. To find out the exact machine architecture (such as Intel 64 and AMD64, also known as x86_64) and print it to standard output, use the following command:

```bash
$ uname -m
```

The following table lists the minimum memory requirements to automatically reserve a memory size for kdump on the latest available versions. The size changes according to the system’s architecture and total available physical memory.

Table 11.1. Minimum amount of reserved memory required for kdump

<table>
<thead>
<tr>
<th>Architecture</th>
<th>Available Memory</th>
<th>Minimum Reserved Memory</th>
</tr>
</thead>
<tbody>
<tr>
<td>AMD64 and Intel 64 (x86_64)</td>
<td>1 GB to 4 GB</td>
<td>160 MB of RAM.</td>
</tr>
<tr>
<td></td>
<td>4 GB to 64 GB</td>
<td>192 MB of RAM.</td>
</tr>
<tr>
<td></td>
<td>64 GB to 1 TB</td>
<td>256 MB of RAM.</td>
</tr>
<tr>
<td></td>
<td>1 TB and more</td>
<td>512 MB of RAM.</td>
</tr>
<tr>
<td>64-bit ARM architecture (arm64)</td>
<td>2 GB and more</td>
<td>448 MB of RAM.</td>
</tr>
<tr>
<td>IBM Power Systems (ppc64le)</td>
<td>2 GB to 4 GB</td>
<td>384 MB of RAM.</td>
</tr>
<tr>
<td></td>
<td>4 GB to 16 GB</td>
<td>512 MB of RAM.</td>
</tr>
<tr>
<td></td>
<td>16 GB to 64 GB</td>
<td>1 GB of RAM.</td>
</tr>
<tr>
<td></td>
<td>64 GB to 128 GB</td>
<td>2 GB of RAM.</td>
</tr>
<tr>
<td></td>
<td>128 GB and more</td>
<td>4 GB of RAM.</td>
</tr>
<tr>
<td>IBM Z (s390x)</td>
<td>1 GB to 4 GB</td>
<td>160 MB of RAM.</td>
</tr>
<tr>
<td></td>
<td>4 GB to 64 GB</td>
<td>192 MB of RAM.</td>
</tr>
<tr>
<td></td>
<td>64 GB to 1 TB</td>
<td>256 MB of RAM.</td>
</tr>
</tbody>
</table>
On many systems, kdump is able to estimate the amount of required memory and reserve it automatically. This behavior is enabled by default, but only works on systems that have more than a certain amount of total available memory, which varies based on the system architecture.

**IMPORTANT**

The configuration of reserved memory based on the total amount of memory in the system is a best effort estimation. The actual required memory may vary due to other factors such as I/O devices. Using not enough of memory might cause that a debug kernel is not able to boot as a capture kernel in case of a kernel panic. To avoid this problem, sufficiently increase the crash kernel memory.

Additional resources

- How has the crashkernel parameter changed between RHEL8 minor releases?
- Technology capabilities and limits tables

### 11.2. MINIMUM THRESHOLD FOR MEMORY RESERVATION

The **kexec-tools** utility, by default, configures the crashkernel command line parameter and reserves a certain amount of memory for kdump. For the default memory reservation to work, a certain amount of total memory must be available in the system. The amount of memory required differs based on the system’s architecture.

The following table lists the minimum threshold values for memory allocation. If the system has memory less than the specified threshold value, you must configure the memory manually.

**Table 11.2. Minimum amount of memory required for memory reservation**

<table>
<thead>
<tr>
<th>Architecture</th>
<th>Required Memory</th>
</tr>
</thead>
<tbody>
<tr>
<td>AMD64 and Intel 64 (<strong>x86_64</strong>)</td>
<td>1 GB</td>
</tr>
<tr>
<td>IBM Power Systems (<strong>ppc64le</strong>)</td>
<td>2 GB</td>
</tr>
<tr>
<td>IBM Z (<strong>s390x</strong>)</td>
<td>1 GB</td>
</tr>
<tr>
<td>ARM (aarch64)</td>
<td>2 GB</td>
</tr>
</tbody>
</table>

### 11.3. SUPPORTED KDUMP TARGETS

When a kernel crash is captured, the vmcore dump file can be either written directly to a device, stored as a file on a local file system, or sent over a network. The table below contains a complete list of dump targets that are currently supported or explicitly unsupported by kdump.
<table>
<thead>
<tr>
<th>Type</th>
<th>Supported Targets</th>
<th>Unsupported Targets</th>
</tr>
</thead>
<tbody>
<tr>
<td>Raw device</td>
<td>All locally attached raw disks and partitions.</td>
<td>Any local file system not explicitly listed as supported in this table, including the auto type (automatic file system detection).</td>
</tr>
<tr>
<td>Local file system</td>
<td>ext2, ext3, ext4, and xfs file systems on directly attached disk drives, hardware RAID logical drives, LVM devices, and mdraid arrays.</td>
<td></td>
</tr>
<tr>
<td>Remote directories accessed using the iSCSI protocol over both hardware and software initiators.</td>
<td>Remote directories accessed using the iSCSI protocol on be2iscsi hardware.</td>
<td>Multipath-based storages.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Remote directories accessed using the SMB or CIFS protocol.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Remote directories accessed using the FCoE (Fibre Channel over Ethernet) protocol.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Remote directories accessed using wireless network interfaces.</td>
</tr>
</tbody>
</table>

**IMPORTANT**

Utilizing firmware assisted dump (fadump) to capture a vmcore and store it to a remote machine using SSH or NFS protocol causes renaming of the network interface to `kdump-<interface-name>`. The renaming happens if the `<interface-name>` is generic, for example `*eth#, net#`, and so on. This problem occurs because the vmcore capture scripts in the initial RAM disk (`initrd`) add the `kdump-` prefix to the network interface name to secure persistent naming. Since the same `initrd` is used also for a regular boot, the interface name is changed for the production kernel too.

**11.4. SUPPORTED KDUMP FILTERING LEVELS**

To reduce the size of the dump file, `kdump` uses the `makedumpfile` core collector to compress the data and optionally to omit unwanted information. The table below contains a complete list of filtering levels that are currently supported by the `makedumpfile` utility.
### Option | Description
--- | ---
1 | Zero pages
2 | Cache pages
4 | Cache private
8 | User pages
16 | Free pages

**NOTE**

The `makedumpfile` command supports removal of transparent huge pages and hugetlbfs pages. Consider both these types of hugepages User Pages and remove them using the `-8` level.

### 11.5. SUPPORTED DEFAULT FAILURE RESPONSES

By default, when `kdump` fails to create a core dump, the operating system reboots. You can, however, configure `kdump` to perform a different operation in case it fails to save the core dump to the primary target. The table below lists all default actions that are currently supported.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>dump_to_rootfs</code></td>
<td>Attempt to save the core dump to the root file system. This option is especially useful in combination with a network target: if the network target is unreachable, this option configures <code>kdump</code> to save the core dump locally. The system is rebooted afterwards.</td>
</tr>
<tr>
<td><code>reboot</code></td>
<td>Reboot the system, losing the core dump in the process.</td>
</tr>
<tr>
<td><code>halt</code></td>
<td>Halt the system, losing the core dump in the process.</td>
</tr>
<tr>
<td><code>poweroff</code></td>
<td>Power off the system, losing the core dump in the process.</td>
</tr>
<tr>
<td><code>shell</code></td>
<td>Run a shell session from within the initramfs, allowing the user to record the core dump manually.</td>
</tr>
<tr>
<td><code>final_action</code></td>
<td>Enable additional operations such as <code>reboot</code>, <code>halt</code>, and <code>poweroff</code> actions after a successful <code>kdump</code> or when <code>shell</code> or <code>dump_to_rootfs</code> failure action completes. The default <code>final_action</code> option is <code>reboot</code>.</td>
</tr>
</tbody>
</table>
11.6. USING FINAL_ACTION PARAMETER

The `final_action` parameter enables you to use certain additional operations such as `reboot`, `halt`, and `poweroff` actions after a successful `kdump` or when an invoked `failure_action` mechanism using `shell` or `dump_to_rootfs` completes. If the `final_action` option is not specified, it defaults to `reboot`.

**Procedure**

1. To configure `final_action`, edit the `/etc/kdump.conf` file and add one of the following options:

   ```
   # final_action <reboot | halt | poweroff>
   ```

2. Restart the `kdump` service for the changes to take effect:

   ```
   # kdumpctl restart
   ```

11.7. USING FAILURE_ACTION PARAMETER

The `failure_action` parameter specifies the action to perform when a dump fails in the event of a kernel crash. The default action for `failure_action` is `reboot`, which reboots the system.

`failure_action` specifies one of the following actions to take:

- `reboot`: reboots the system after a dump failure.
- `dump_to_rootfs`: saves the dump file on a root file system when a non-root dump target is configured.
- `halt`: halts the system.
- `poweroff`: stops the running operations on the system.
- `shell`: starts a shell session inside `initramfs`, from which you can manually perform additional recovery actions.

**Procedure**

1. To configure an action to take if the dump fails, edit the `/etc/kdump.conf` file and specify one of the `failure_action` options:

   ```
   # failure_action <reboot | halt | poweroff | shell | dump_to_rootfs>
   ```

2. Restart the `kdump` service for the changes to take effect:

   ```
   # kdumpctl restart
   ```
Firmware assisted dump (fadump) is a dump capturing mechanism, provided as an alternative to the kdump mechanism on IBM POWER systems. The kexec and kdump mechanisms are useful for capturing core dumps on AMD64 and Intel 64 systems. However, some hardware such as mini systems and mainframe computers, leverage the onboard firmware to isolate regions of memory and prevent any accidental overwriting of data that is important to the crash analysis. This section covers fadump mechanisms and how they integrate with RHEL. The fadump utility is optimized for these expanded dumping features on IBM POWER systems.

12.1. FIRMWARE ASSISTED DUMP ON IBM POWERPC HARDWARE

The fadump utility captures the vmcore file from a fully-reset system with PCI and I/O devices. This mechanism uses firmware to preserve memory regions during a crash and then reuses the kdump userspace scripts to save the vmcore file. The memory regions consist of all system memory contents, except the boot memory, system registers, and hardware Page Table Entries (PTEs).

The fadump mechanism offers improved reliability over the traditional dump type, by rebooting the partition and using a new kernel to dump the data from the previous kernel crash. The fadump requires an IBM POWER6 processor-based or later version hardware platform.

For further details about the fadump mechanism, including PowerPC specific methods of resetting hardware, see the /usr/share/doc/kexec-tools/fadump-howto.txt file.

**NOTE**

The area of memory that is not preserved, known as boot memory, is the amount of RAM required to successfully boot the kernel after a crash event. By default, the boot memory size is 256MB or 5% of total system RAM, whichever is larger.

Unlike kexec-initiated event, the fadump mechanism uses the production kernel to recover a crash dump. When booting after a crash, PowerPC hardware makes the device node /proc/device-tree/rtas/ibm.kernel-dump available to the proc filesystem (procfs). The fadump-aware kdump scripts, check for the stored vmcore, and then complete the system reboot cleanly.

12.2. ENABLING FIRMWARE ASSISTED DUMP MECHANISM

You can enhance the crash dumping capabilities of IBM POWER systems by enabling the firmware assisted dump (fadump) mechanism

**Prerequisites**

- Root access

**Procedure**

1. Install the kexec-tools package.

2. Configure the default value for crashkernel.

   ```
   # kdumpctl reset-crashkernel --fadump=on --kernel=ALL
   ```

3. (Optional) Reserve boot memory instead of the default value.
# grubby --update-kernel ALL --args="fadump=on crashkernel=xxM"

where, xx is the required memory size in megabytes.

**NOTE**

When specifying boot configuration options, test the configuration by rebooting the kernel with `kdump` enabled. If the `kdump` kernel fails to boot, increase the `crashkernel` value gradually to set an appropriate value.

4. Reboot for changes to take effect.

```bash
# reboot
```

### 12.3. FIRMWARE ASSISTED DUMP MECHANISMS ON IBM Z HARDWARE

IBM Z systems supports two firmware assisted dump mechanisms: Stand-alone dump (`sadump`) and `VMDUMP` dump file.

The `kdump` infrastructure is supported and utilized on IBM Z systems. However, using one of the firmware assisted dump (`fadump`) methods for IBM Z can provide various benefits:

- The `sadump` mechanism is initiated and controlled from the system console, and is stored on an IPL bootable device.
- The `VMDUMP` mechanism is similar to `sadump`. This tool is initiated from the system console, but retrieves the resulting dump from hardware and copies it to the system for analysis.
- These methods, similar to other hardware-based dump mechanisms, have the ability to capture the state of a machine in the early boot phase, before the `kdump` service starts.
- Although `VMDUMP` contains a mechanism to receive the dump file into a Red Hat Enterprise Linux system, the configuration and control of `VMDUMP` is managed from the IBM Z Hardware console.

**Additional resources**

- Using the Dump Tools on Red Hat Enterprise Linux 8.5
- Stand-alone dump
- Creating dumps on z/VM with VMDUMP

### 12.4. USING SADUMP ON FUJITSU PRIMEQUEST SYSTEMS

The Fujitsu `sadump` mechanism is designed to provide a fallback dump capture in an event when `kdump` is unable to complete successfully. The `sadump` mechanism is invoked manually from the system Management Board (MMB) interface. Using MMB, configure `kdump` like for an Intel 64 or AMD 64 server and then perform the following additional steps to enable `sadump`.

**Procedure**
1. Add or edit the following lines in the `/etc/sysctl.conf` file to ensure that `kdump` starts as expected for `sadump`:

   ```
   kernel.panic=0
   kernel.unknown_nmi_panic=1
   ```

   **WARNING**

   In particular, ensure that after `kdump`, the system does not reboot. If the system reboots after `kdump` has fails to save the `vmcore` file, then it is not possible to invoke the `sadump`.

2. Set the `failure_action` parameter in `/etc/kdump.conf` appropriately as `halt` or `shell`.

   ```
   failure_action shell
   ```

Additional resources

CHAPTER 13. ANALYZING A CORE DUMP

To determine the cause of the system crash, you can use the `crash` utility, which provides an interactive prompt very similar to the GNU Debugger (GDB). This utility allows you to interactively analyze a core dump created by `kdump`, `netdump`, `diskdump` or `xendump` as well as a running Linux system. Alternatively, you have the option to use Kernel Oops Analyzer or the Kdump Helper tool.

13.1. INSTALLING THE CRASH UTILITY

The following procedure describes how to install the `crash` analyzing tool.

Procedure

1. Enable the relevant repositories:

   ```
   # subscription-manager repos --enable baseos repository
   # subscription-manager repos --enable appstream repository
   # subscription-manager repos --enable rhel-8-for-x86_64-baseos-debug-rpms
   ```

2. Install the `crash` package:

   ```
   # dnf install crash
   ```

3. Install the `kernel-debuginfo` package:

   ```
   # dnf install kernel-debuginfo
   ```

   The package corresponds to your running kernel and provides the data necessary for the dump analysis.

Additional resources

- Configuring basic system settings

13.2. RUNNING AND EXITING THE CRASH UTILITY

The following procedure describes how to start the crash utility for analyzing the cause of the system crash.

Prerequisites

- Identify the currently running kernel (for example 5.14.0-1.el9.x86_64).

Procedure

1. To start the `crash` utility, two necessary parameters need to be passed to the command:

   - The debug-info (a decompressed vmlinuz image), for example `/usr/lib/debug/lib/modules/5.14.0-1.el9.x86_64/vmlinux` provided through a specific `kernel-debuginfo` package.
• The actual vmcore file, for example /var/crash/127.0.0.1-2021-09-13-14:05:33/vmcore
  The resulting crash command then looks like this:

  # crash /usr/lib/debug/lib/modules/5.14.0-1.el9.x86_64/vmlinux /var/crash/127.0.0.1-
  2021-09-13-14:05:33/vmcore

  Use the same <kernel> version that was captured by kdump.

**Example 13.1. Running the crash utility**

The following example shows analyzing a core dump created on September 13 2021 at
14:05 PM, using the 5.14.0-1.el9.x86_64 kernel.

```
WARNING: kernel relocated [202MB]: patching 90160 gdb minimal_symbol values

  KERNEL: /usr/lib/debug/lib/modules/5.14.0-1.el9.x86_64/vmlinux
  DUMPFILE: /var/crash/127.0.0.1-2021-09-13-14:05:33/vmcore [PARTIAL DUMP]
  CPUS: 2
  DATE: Mon Sep 13 14:05:16 2021
  UPTIME: 01:03:57
  LOAD AVERAGE: 0.00, 0.00, 0.00
  TASKS: 586
  NODENAME: localhost.localdomain
  RELEASE: 5.14.0-1.el9.x86_64
  VERSION: #1 SMP Wed Aug 29 11:51:55 UTC 2018
  MACHINE: x86_64 (2904 Mhz)
  MEMORY: 2.9 GB
  PANIC: "sysrq: SysRq : Trigger a crash"
  PID: 10635
  COMMAND: "bash"
  TASK: ffff8d6c84271800 [THREAD_INFO: ffff8d6c84271800]
  CPU: 1
  STATE: TASK_RUNNING (SYSRQ)

  crash>
```

2. To exit the interactive prompt and terminate crash, type exit or q.

**Example 13.2. Exiting the crash utility**

```
  crash> exit
  ~>#
```

**NOTE**

The crash command can also be used as a powerful tool for debugging a live system.
However use it with caution so as not to break your system.

**Additional resources**

- [A Guide to Unexpected System Restarts](#)
13.3. DISPLAYING VARIOUS INDICATORS IN THE CRASH UTILITY

The following procedures describe how to use the crash utility and display various indicators, such as a kernel message buffer, a backtrace, a process status, virtual memory information and open files.

Displaying the message buffer

- To display the kernel message buffer, type the `log` command at the interactive prompt as displayed in the example below:

```
crash> log
... several lines omitted ...
EIP: 0060:[<c068124f>] EFLAGS: 00010096 CPU: 2
EIP is at sysrq_handle_crash+0xf/0x20
EAX: 00000063 EBX: 00000063 ECX: c09e1c8c EDX: 00000000
ESI: c0a9ca0 EDI: 00000286 EBP: 00000000 ESP: ef4dbf24
DS: 007b ES: 007b FS: 00d8 GS: 00e0 SS: 0068
Process bash (pid: 5591, ti=ef4da000 task=f196d560 task.ti=ef4da000)
Stack:
c068146b c0960891 c0968653 00000003 00000000 00000002 efade5c0 c06814d0
<0> ffffffff c068150f b7776000 f2600c40 c0569ec4 ef4dbf9c 00000002 b7776000
<0> efade5c0 00000002 b7776000 c0569e60 c051de50 ef4dbf9c f196d560 ef4dbfb4
Call Trace:
[<c068146b>] ? __handle_sysrq+0xfb/0x160
[<c06814d0>] ? write_sysrq_trigger+0x0/0x50
[<c068150f>] ? write_sysrq_trigger+0x3f/0x50
[<c0569ec4>] ? proc_reg_write+0x64/0xa0
[<c0569e60>] ? proc_reg_write+0x0/0xa0
[<c051de50>] ? vfs_write+0xa0/0x190
[<c051e8d1>] ? sys_write+0x41/0x70
[<c0409adc>] ? syscall_call+0x7/0xb
Code: a0 c0 01 0f b6 41 03 19 d2 f7 d2 83 e2 03 83 e0 cf c1 e2 04 09 d0 88 41 03 f3 c3 90 c7 05
c8 1b 9e c0 01 00 00 00 0f ae f8 89 f6 <c6> 05 00 00 00 00 01 c3 89 f6 8d bc 27 00 00 00 00 8d 50
d0 83
EIP: [<c068124f>] sysrq_handle_crash+0xf/0x20 SS:ESP 0068:ef4dbf24
CR2: 0000000000000000
```

Type `help log` for more information on the command usage.

NOTE

The kernel message buffer includes the most essential information about the system crash and, as such, it is always dumped first in to the `vmcore-dmesg.txt` file. This is useful when an attempt to get the full `vmcore` file failed, for example because of lack of space on the target location. By default, `vmcore-dmesg.txt` is located in the `/var/crash/` directory.

Displaying a backtrace

- To display the kernel stack trace, use the `bt` command.

```
crash> bt
PID: 5591   TASK: f196d560 CPU: 2   COMMAND: "bash"
#0 [ef4dbdce] crash_kexec at c0499222
#1 [ef4dbdce] crash_kexec at c0499222
```
Displaying a process status

- To display the status of processes in the system, use the `ps` command.

```
crash> ps
    PID    PPID  CPU   TASK    ST  %MEM     VSZ    RSS  COMM
>     0      0   0  c09dc560  RU   0.0       0      0  [swapper]
>     0      0   1  f7072030  RU   0.0       0      0  [swapper]
  0      0   2  f70a3a90  RU   0.0       0      0  [swapper]
>     0      0   3  f70ac560  RU   0.0       0      0  [swapper]
  1      0   1  f705ba90  IN   0.0    2828   1424  init
... several lines omitted ...
  5566   1   1  f2592560  IN   0.0   12876   784 auditd
  5567   1   2  ef427560  IN   0.0   12876   784 auditd
  5587  5132   0  f196d030  IN   0.0  11064  3184 sshd
>  5591  5587   2  f196d560  RU   0.0   5084  1648 bash
```

Use `ps <pid>` to display the status of a single specific process. Use `help ps` for more information on `ps` usage.

Displaying virtual memory information

- To display basic virtual memory information, type the `vm` command at the interactive prompt.

```
crash> vm
     PID: 5591   TASK: f196d560   CPU: 2   COMMAND: "bash"
     MM   PGD   RSS   TOTAL_VM
f19b5900  ef9c6000  1648k    5084k
VMA   START   END   FLAGS   FILE
f1bb0310  242000  260000 8000875  /lib/ld-2.12.so
f26af0b8  260000  261000 8100871  /lib/ld-2.12.so
efbc275c  261000  262000 8100873  /lib/ld-2.12.so
```
Use `vm <pid>` to display information on a single specific process, or use `help vm` for more information on `vm` usage.

Displaying open files

- To display information about open files, use the `files` command.

```
crash> files
PID: 5591 TASK: f196d560 CPU: 2 COMMAND: "bash"
ROOT: / CWD: /root
FD FILE DENTRY INODE TYPE PATH
0 f734f640 eedc2c6c eecd6048 CHR /pts/0
1 efade5c0 eee14090 f00431d4 REG /proc/sysrq-trigger
2 f734f640 eedc2c6c eecd6048 CHR /pts/0
10 f734f640 eedc2c6c eecd6048 CHR /pts/0
255 f734f640 eedc2c6c eecd6048 CHR /pts/0
```

Use `files <pid>` to display files opened by only one selected process, or use `help files` for more information on `files` usage.

### 13.4. USING KERNEL OOPS ANALYZER

The Kernel Ooops Analyzer tool analyzes the crash dump by comparing the oops messages with known issues in the knowledge base.

**Prerequisites**

- Secure an oops message to feed the Kernel Ooops Analyzer.

**Procedure**

1. Access the Kernel Ooops Analyzer tool.

2. To diagnose a kernel crash issue, upload a kernel oops log generated in `vmcore`.  

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• Alternatively you can also diagnose a kernel crash issue by providing a text message or a `vmcore-dmesg.txt` as an input.

3. Click DETECT to compare the oops message based on information from the `makedumpfile` against known solutions.

Additional resources

• The Kernel Oops Analyzer article

• A Guide to Unexpected System Restarts

13.5. THE KDUMP HELPER TOOL

The Kdump Helper tool helps to set up the kdump using the provided information. Kdump Helper generates a configuration script based on your preferences. Initiating and running the script on your server sets up the kdump service.

Additional resources

• Kdump Helper
CHAPTER 14. APPLYING PATCHES WITH KERNEL LIVE PATCHING

You can use the Red Hat Enterprise Linux kernel live patching solution to patch a running kernel without rebooting or restarting any processes.

With this solution, system administrators:

- Can immediately apply critical security patches to the kernel.
- Do not have to wait for long-running tasks to complete, for users to log off, or for scheduled downtime.
- Control the system's uptime more and do not sacrifice security or stability.

Note that not every critical or important CVE will be resolved using the kernel live patching solution. Our goal is to reduce the required reboots for security-related patches, not to eliminate them entirely. For more details about the scope of live patching, see the Customer Portal Solutions article.

WARNING

Some incompatibilities exist between kernel live patching and other kernel subcomponents. Read the Limitations of kpatch carefully before using kernel live patching.

14.1. LIMITATIONS OF KPATCH

- The kpatch feature is not a general-purpose kernel upgrade mechanism. It is used for applying simple security and bug fix updates when rebooting the system is not immediately possible.

- Do not use the SystemTap or kprobe tools during or after loading a patch. The patch could fail to take effect until after such probes have been removed.

14.2. SUPPORT FOR THIRD-PARTY LIVE PATCHING

The kpatch utility is the only kernel live patching utility supported by Red Hat with the RPM modules provided by Red Hat repositories. Red Hat will not support any live patches which were not provided by Red Hat itself.

If you require support for an issue that arises with a third-party live patch, Red Hat recommends that you open a case with the live patching vendor at the outset of any investigation in which a root cause determination is necessary. This allows the source code to be supplied if the vendor allows, and for their support organization to provide assistance in root cause determination prior to escalating the investigation to Red Hat Support.

For any system running with third-party live patches, Red Hat reserves the right to ask for reproduction with Red Hat shipped and supported software. In the event that this is not possible, we require a similar system and workload be deployed on your test environment without live patches applied, to confirm if the same behavior is observed.
For more information about third-party software support policies, see How does Red Hat Global Support Services handle third-party software, drivers, and/or uncertified hardware/hypervisors or guest operating systems?

14.3. ACCESS TO KERNEL LIVE PATCHES

Kernel live patching capability is implemented as a kernel module (kmod) that is delivered as an RPM package.

All customers have access to kernel live patches, which are delivered through the usual channels. However, customers who do not subscribe to an extended support offering will lose access to new patches for the current minor release once the next minor release becomes available. For example, customers with standard subscriptions will only be able to live patch RHEL 9.1 kernel until the RHEL 9.2 kernel is released.

14.4. COMPONENTS OF KERNEL LIVE PATCHING

The components of kernel live patching are as follows:

Kernel patch module
- The delivery mechanism for kernel live patches.
- A kernel module which is built specifically for the kernel being patched.
- The patch module contains the code of the desired fixes for the kernel.
- The patch modules register with the livepatch kernel subsystem and provide information about original functions to be replaced, with corresponding pointers to the replacement functions. Kernel patch modules are delivered as RPMs.
- The naming convention is kpatch_<kernel version>_<kpatch version>_<kpatch release>. The "kernel version" part of the name has dots replaced with underscores.

The kpatch utility
A command-line utility for managing patch modules.

The kpatch service
A systemd service required by multiuser.target. This target loads the kernel patch module at boot time.

The kpatch-dnf package
A DNF plugin delivered in the form of an RPM package. This plugin manages automatic subscription to kernel live patches.

14.5. HOW KERNEL LIVE PATCHING WORKS

The kpatch kernel patching solution uses the livepatch kernel subsystem to redirect old functions to new ones. When a live kernel patch is applied to a system, the following things happen:

1. The kernel patch module is copied to the /var/lib/kpatch/ directory and registered for re-application to the kernel by systemd on next boot.

2. The kpatch module is loaded into the running kernel and the new functions are registered to the ftrace mechanism with a pointer to the location in memory of the new code.
3. When the kernel accesses the patched function, it is redirected by the `ftrace` mechanism which bypasses the original functions and redirects the kernel to patched version of the function.

**Figure 14.1. How kernel live patching works**

A kernel patch module is delivered in an RPM package, specific to the version of the kernel being patched. Each RPM package will be cumulatively updated over time.

The following procedure explains how to subscribe to all future cumulative live patching updates for a given kernel. Because live patches are cumulative, you cannot select which individual patches are deployed for a given kernel.

**WARNING**

Red Hat does not support any third party live patches applied to a Red Hat supported system.

**Prerequisites**

- Root permissions

**Procedure**

1. Optionally, check your kernel version:

```
# uname -r
5.14.0-1.el9.x86_64
```
2. Search for a live patching package that corresponds to the version of your kernel:

```
# dnf search $(uname -r)
```

3. Install the live patching package:

```
# dnf install "kpatch-patch = $(uname -r)"
```

The command above installs and applies the latest cumulative live patches for that specific kernel only.

If the version of a live patching package is 1-1 or higher, the package will contain a patch module. In that case the kernel will be automatically patched during the installation of the live patching package.

The kernel patch module is also installed into the `/var/lib/kpatch/` directory to be loaded by the `systemd` system and service manager during the future reboots.

**NOTE**

An empty live patching package will be installed when there are no live patches available for a given kernel. An empty live patching package will have a `kpatch_version-kpatch_release` of 0-0, for example `kpatch-patch-5_14_0-1-0-0.x86_64.rpm`. The installation of the empty RPM subscribes the system to all future live patches for the given kernel.

4. Optionally, verify that the kernel is patched:

```
# kpatch list
Loaded patch modules:
kpatch_5_14_0_1_0_1 [enabled]

Installed patch modules:
kpatch_5_14_0_1_0_1 (5.14.0-1.el9.x86_64)
...
```

The output shows that the kernel patch module has been loaded into the kernel, which is now patched with the latest fixes from the `kpatch-patch-5_14_0-1-0-1.el9.x86_64.rpm` package.

**Additional resources**

- `kpatch(1)` manual page
- `Configuring basic system settings` in RHEL

**14.7. AUTOMATICALLY SUBSCRIBING ANY FUTURE KERNEL TO THE LIVE PATCHING STREAM**

You can use the `kpatch-dnf` DNF plugin to subscribe your system to fixes delivered by the kernel patch module, also known as kernel live patches. The plugin enables automatic subscription for any kernel the system currently uses, and also for kernels to-be-installed in the future.

**Prerequisites**
You have root permissions.

**Procedure**

1. Optionally, check all installed kernels and the kernel you are currently running:

   ```
   # dnf list installed | grep kernel
   Updating Subscription Management repositories.
   Installed Packages
   ...
   kernel-core.x86_64            5.14.0-1.el9              @beaker-BaseOS
   kernel-core.x86_64            5.14.0-2.el9              @@commandline
   ...
   # uname -r
   5.14.0-2.el9.x86_64
   ```

2. Install the **kpatch-dnf** plugin:

   ```
   # dnf install kpatch-dnf
   ```

3. Enable automatic subscription to kernel live patches:

   ```
   # dnf kpatch auto
   Updating Subscription Management repositories.
   Last metadata expiration check: 1:38:21 ago on Fri 17 Sep 2021 07:29:53 AM EDT.
   Dependencies resolved.
   Package Architecture
   Installing:
   kpatch-patch-5_14_0-1               x86_64
   kpatch-patch-5_14_0-2               x86_64
   Transaction Summary
   Install  2 Packages
   ```

This command subscribes all currently installed kernels to receiving kernel live patches. The command also installs and applies the latest cumulative live patches, if any, for all installed kernels.

In the future, when you update the kernel, live patches will automatically be installed during the new kernel installation process.

The kernel patch module is also installed into the `/var/lib/kpatch/` directory to be loaded by the `systemd` system and service manager during future reboots.
NOTE

An empty live patching package will be installed when there are no live patches available for a given kernel. An empty live patching package will have a `kpatch_version-kpatch_release` of 0-0, for example `kpatch-patch-5_14_0-1-0-0.el9.x86_64.rpm`. The installation of the empty RPM subscribes the system to all future live patches for the given kernel.

Verification step

- Verify that all installed kernels have been patched:

```
# kpatch list
Loaded patch modules:
  kpatch_5_14_0_2_0_1 [enabled]

Installed patch modules:
  kpatch_5_14_0_1_0_1 (5.14.0-1.el9.x86_64)
  kpatch_5_14_0_2_0_1 (5.14.0-2.el9.x86_64)
```

The output shows that both the kernel you are running, and the other installed kernel have been patched with fixes from `kpatch-patch-5_14_0-1-0-1.el9.x86_64.rpm` and `kpatch-patch-5_14_0-2-0-1.el9.x86_64.rpm` packages respectively.

Additional resources

- `kpatch(1)` and `dnf-kpatch(8)` manual pages
- Configuring basic system settings in RHEL

14.8. DISABLING AUTOMATIC SUBSCRIPTION TO THE LIVE PATCHING STREAM

When you subscribe your system to fixes delivered by the kernel patch module, your subscription is automatic. You can disable this feature, and thus disable automatic installation of `kpatch-patch` packages.

Prerequisites

- You have root permissions.

Procedure

1. Optionally, check all installed kernels and the kernel you are currently running:

```
# dnf list installed | grep kernel
Updating Subscription Management repositories.
Installed Packages
  ...
  kernel-core.x86_64  5.14.0-1.el9   @beaker-BaseOS
  kernel-core.x86_64  5.14.0-2.el9   @@commandline
  ...
```
# unittest

5.14.0-2.el9.x86_64

2. Disable automatic subscription to kernel live patches:

```
# dnf kpatch manual
Updating Subscription Management repositories.
```

**Verification step**

- You can check for the successful outcome:

```
# yum kpatch status
```

**Additional resources**

- `kpatch(1)` and `dnf-kpatch(8)` manual pages

### 14.9. UPDATING KERNEL PATCH MODULES

Since kernel patch modules are delivered and applied through RPM packages, updating a cumulative kernel patch module is like updating any other RPM package.

**Prerequisites**

- The system is subscribed to the live patching stream, as described in [Subscribing the currently installed kernels to the live patching stream](#).

**Procedure**

- Update to a new cumulative version for the current kernel:

  ```
  # dnf update "kpatch-patch = $(uname -r)"
  ```

  The command above automatically installs and applies any updates that are available for the currently running kernel. Including any future released cumulative live patches.

- Alternatively, update all installed kernel patch modules:

  ```
  # dnf update "kpatch-patch"
  ```

**NOTE**

When the system reboots into the same kernel, the kernel is automatically live patched again by the `kpatch.service` systemd service.

**Additional resources**
14.10. REMOVING THE LIVE PATCHING PACKAGE

The following procedure describes how to disable the Red Hat Enterprise Linux kernel live patching solution by removing the live patching package.

Prerequisites

- Root permissions
- The live patching package is installed.

Procedure

1. Select the live patching package:

   ```
   # dnf list installed | grep kpatch-patch
   kpatch-patch-5_14_0-1.x86_64 0-1.el9 @@commandline
   ...
   ```

   The example output above lists live patching packages that you installed.

2. Remove the live patching package:

   ```
   # dnf remove kpatch-patch-5_14_0-1.x86_64
   ```

   When a live patching package is removed, the kernel remains patched until the next reboot, but the kernel patch module is removed from disk. On future reboot, the corresponding kernel will no longer be patched.

3. Reboot your system.

4. Verify that the live patching package has been removed:

   ```
   # dnf list installed | grep kpatch-patch
   ```

   The command displays no output if the package has been successfully removed.

5. Optionally, verify that the kernel live patching solution is disabled:

   ```
   # kpatch list
   Loaded patch modules:
   ```

   The example output shows that the kernel is not patched and the live patching solution is not active because there are no patch modules that are currently loaded.

**IMPORTANT**

Currently, Red Hat does not support reverting live patches without rebooting your system. In case of any issues, contact our support team.

Additional resources
Prerequisites

- Root permissions
- A live patching package is installed.
- A kernel patch module is installed and loaded.

Procedure

1. Select a kernel patch module:

   ```
   # kpatch list
   Loaded patch modules:
   kpatch_5_14_0_1_0_1 [enabled]
   
   Installed patch modules:
   kpatch_5_14_0_1_0_1 (5.14.0-1.el9.x86_64)
   ... 
   ```

2. Uninstall the selected kernel patch module:

   ```
   # kpatch uninstall kpatch_5_14_0_1_0_1
   uninstalling kpatch_5_14_0_1_0_1 (5.14.0-1.el9.x86_64)
   ```

   - Note that the uninstalled kernel patch module is still loaded:

   ```
   # kpatch list
   Loaded patch modules:
   kpatch_5_14_0_1_0_1 [enabled]
   
   Installed patch modules:
   <NO_RESULT>
   ```

   When the selected module is uninstalled, the kernel remains patched until the next reboot, but the kernel patch module is removed from disk.

3. Reboot your system.

4. Optionally, verify that the kernel patch module has been uninstalled:

   ```
   # kpatch list
   Loaded patch modules:
   ...
   ```
The example output above shows no loaded or installed kernel patch modules, therefore the kernel is not patched and the kernel live patching solution is not active.

**IMPORTANT**
Currently, Red Hat does not support reverting live patches without rebooting your system. In case of any issues, contact our support team.

Additional resources
- kpatch(1) manual page

### 14.12. DISABLING KPATCH.SERVICE

The following procedure describes how to prevent the Red Hat Enterprise Linux kernel live patching solution from applying all kernel patch modules globally on subsequent boots.

**Prerequisites**
- Root permissions
- A live patching package is installed.
- A kernel patch module is installed and loaded.

**Procedure**

1. Verify **kpatch.service** is enabled:

   ```bash
   # systemctl is-enabled kpatch.service
   enabled
   ```

2. Disable **kpatch.service**:

   ```bash
   # systemctl disable kpatch.service
   Removed /etc/systemd/system/multi-user.target.wants/kpatch.service.
   ```

   - Note that the applied kernel patch module is still loaded:

     ```bash
     # kpatch list
     Loaded patch modules:
     kpatch_5_14_0_1_0_1 [enabled]
     Installed patch modules:
     kpatch_5_14_0_1_0_1 (5.14.0-1.el9.x86_64)
     ```

3. Reboot your system.

4. Optionally, verify the status of **kpatch.service**:

   ```bash
   # systemctl status kpatch.service
   kpatch.service - "Apply kpatch kernel patches"
   Loaded: loaded (/usr/lib/systemd/system/kpatch.service; disabled; vendor preset: disabled)
   Active: inactive (dead)
   ```
The example output testifies that `kpatch.service` has been disabled and is not running. Thereby, the kernel live patching solution is not active.

5. Verify that the kernel patch module has been unloaded:

```bash
# kpatch list
Loaded patch modules:

Installed patch modules:
   kpatch_5_14_0_1_0_1 (5.14.0-1.el9.x86_64)
```

The example output above shows that a kernel patch module is still installed but the kernel is not patched.

**IMPORTANT**

Currently, Red Hat does not support reverting live patches without rebooting your system. In case of any issues, contact our support team.

Additional resources

- `kpatch(1)` manual page
- *Configuring basic system settings* in RHEL
CHAPTER 15. USING SYSTEMD TO MANAGE RESOURCES USED BY APPLICATIONS

RHEL 9 moves the resource management settings from the process level to the application level by binding the system of cgroup hierarchies with the systemd unit tree. Therefore, you can manage the system resources with the systemctl command, or by modifying the systemd unit files.

To achieve this, systemd takes various configuration options from the unit files or directly via the systemctl command. Then systemd applies those options to specific process groups by utilizing the Linux kernel system calls and features like cgroups and namespaces.

NOTE

You can review the full set of configuration options for systemd in the following manual pages:

- systemd.resource-control(5)
- systemd.exec(5)

15.1. ALLOCATING SYSTEM RESOURCES USING SYSTEMD

To modify the distribution of system resources, you can apply one or more of the following distribution models:

Weights

You can distribute the resource by adding up the weights of all sub-groups and giving each sub-group the fraction matching its ratio against the sum.

For example, if you have 10 cgroups, each with weight of value 100, the sum is 1000. Each cgroup receives one tenth of the resource.

Weight is usually used to distribute stateless resources. For example the CPUWeight= option is an implementation of this resource distribution model.

Limits

A cgroup can consume up to the configured amount of the resource. The sum of sub-group limits can exceed the limit of the parent cgroup. Therefore it is possible to overcommit resources in this model.

For example the MemoryMax= option is an implementation of this resource distribution model.

Protections

You can set up a protected amount of a resource for a cgroup. If the resource usage is below the protection boundary, the kernel will try not to penalize this cgroup in favor of other cgroups that compete for the same resource. An overcommit is also possible.

For example the MemoryLow= option is an implementation of this resource distribution model.

Allocations

Exclusive allocations of an absolute amount of a finite resource. An overcommit is not possible. An example of this resource type in Linux is the real-time budget.

unit file option

A setting for resource control configuration.
For example, you can configure CPU resource with options like `CPUAccounting=` or `CPUQuota=`. Similarly, you can configure memory or I/O resources with options like `AllowedMemoryNodes=` and `IOAccounting=`.

**Procedure**

To change the required value of the unit file option of your service, you can adjust the value in the unit file, or use `systemctl` command:

1. Check the assigned values for the service of your choice.
   
   ```
   # systemctl show --property <unit file option> <service name>
   ```

2. Set the required value of the CPU time allocation policy option:

   ```
   # systemctl set-property <service name> <unit file option>=<value>
   ```

**Verification steps**

- Check the newly assigned values for the service of your choice.

  ```
  # systemctl show --property <unit file option> <service name>
  ```

**Additional resources**

- `systemd.resource-control(5)`, `systemd.exec(5)` manual pages

### 15.2. ROLE OF SYSTEMD IN RESOURCE MANAGEMENT

The core function of `systemd` is service management and supervision. The `systemd` system and service manager ensures that managed services start at the right time and in the correct order during the boot process. The services have to run smoothly to use the underlying hardware platform optimally. Therefore, `systemd` also provides capabilities to define resource management policies, and to tune various options, which can improve the performance of the service.

**IMPORTANT**

In general, Red Hat recommends you use `systemd` for controlling the usage of system resources. You should manually configure the `cgroups` virtual file system only in special cases. For example, when you need to use `cgroup-v1` controllers that have no equivalents in `cgroup-v2` hierarchy.

### 15.3. OVERVIEW OF SYSTEMD HIERARCHY FOR CGROUPS

On the backend, the `systemd` system and service manager makes use of the `slice`, the `scope` and the `service` units to organize and structure processes in the control groups. You can further modify this hierarchy by creating custom unit files or using the `systemctl` command. Also, `systemd` automatically mounts hierarchies for important kernel resource controllers at the `/sys/fs/cgroup/` directory.

Three `systemd` unit types are used for resource control:
• **Service** - A process or a group of processes, which **systemd** started according to a unit configuration file. Services encapsulate the specified processes so that they can be started and stopped as one set. Services are named in the following way:

```
<name>.service
```

• **Scope** - A group of externally created processes. Scopes encapsulate processes that are started and stopped by the arbitrary processes through the **fork()** function and then registered by **systemd** at runtime. For example, user sessions, containers, and virtual machines are treated as scopes. Scopes are named as follows:

```
<name>.scope
```

• **Slice** - A group of hierarchically organized units. Slices organize a hierarchy in which scopes and services are placed. The actual processes are contained in scopes or in services. Every name of a slice unit corresponds to the path to a location in the hierarchy. The dash ("-"), character acts as a separator of the path components to a slice from the **-.slice** root slice. In the following example:

```
<parent-name>.slice
```

**parent-name.slice** is a sub-slice of **parent.slice**, which is a sub-slice of the **-.slice** root slice. **parent-name.slice** can have its own sub-slice named **parent-name-name2.slice**, and so on.

The service, the scope, and the slice units directly map to objects in the control group hierarchy. When these units are activated, they map directly to control group paths built from the unit names.

The following is an abbreviated example of a control group hierarchy:

```
Control group /:
├── -.slice
│   ├── user.slice
│   │   ├── user-42.slice
│   │       ├── session-c1.scope
│   │           ├── 967 gdm-session-worker [pam/gdm-launch-environment]
│   │           ├── 1035 /usr/libexec/gdm-x-session gnome-session --autostart
│   │           /usr/share/gdm/greeter/autostart
│   │               ├── 1054 /usr/libexec/Xorg vt1 -displayfd 3 -auth /run/user/42/gdm/Xauthority -background none
│   │               │   -noreset -keeptty -verbose 3
│   │               │       ├── 1212 /usr/libexec/gnome-session-binary --autostart /usr/share/gdm/greeter/autostart
│   │               │           ├── 1369 /usr/bin/gnome-shell
│   │               │           ├── 1732 ibus-daemon --xim --panel disable
│   │               │           ├── 1752 /usr/libexec/ibus-dconf
│   │               │           ├── 1762 /usr/libexec/ibus-x11 --kill-daemon
│   │               │           ├── 1912 /usr/libexec/gsd-xsettings
│   │               │           ├── 1917 /usr/libexec/gsd-a11y-settings
│   │               │           └── 1920 /usr/libexec/gsd-clipboard
│   │               …
│   │   └── init.scope
│   │       ├── 1 /usr/lib/systemd/systemd --switched-root --system --deserialize 18
│   │       └── systemd.slice
│           └── mgd.service
│               └── 800 /sbin/mgd -f
│   └── systemd-udevd.service
│       └── 659 /usr/lib/systemd/systemd-udevd
```

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The example above shows that services and scopes contain processes and are placed in slices that do not contain processes of their own.

**Additional resources**

- *Configuring basic system settings* in Red Hat Enterprise Linux
- What are kernel resource controllers
- *systemd.resource-control(5), systemd.exec(5), cgroups(7), fork(), fork(2)* manual pages
- Understanding cgroups

**15.4. LISTING SYSTEMD UNITS**

The following procedure describes how to use the *systemd* system and service manager to list its units.

**Procedure**

- To list all active units on the system, execute the `# systemctl` command and the terminal will return an output similar to the following example:

```
# systemctl
UNIT                                                LOAD   ACTIVE SUB       DESCRIPTION
...                                                ...
init.scope                                          loaded active running System and Service Manager
session-2.scope                                     loaded active running Session 2 of user jdoe
abrt-ccpp.service                                   loaded active exited Install ABRT coredump hook
abrt-oops.service                                   loaded active running ABRT kernel log watcher
abrt-vmcore.service                                 loaded active exited Harvest vmcores for ABRT
abrt-xorg.service                                   loaded active running ABRT Xorg log watcher
...                                                ...
-.slice                                             loaded active active  Root Slice
machine.slice                                       loaded active active  Virtual Machine and Container
Slice system-getty.slice                            loaded active active  Virtual Machine and Container
system-getty.slice                                  loaded active active  Virtual Machine and Container
system-lvm2\x2dpvsan.slice                          loaded active active  system-
lvm2\x2dpvsan.slice                                  loaded active active  system-
system-sshd\x2dkeygen.slice                         loaded active active  system-
sshd\x2dkeygen.slice                                 loaded active active  system-
system-systemd\x2dhibernate\x2dresume.slice          loaded active active  system-
systemd\x2dhibernate\x2dresume>                     loaded active active  system-
system-user\x2druntime\x2ddir.slice                 loaded active active  system-
```
UNIT - a name of a unit that also reflects the unit position in a control group hierarchy. The units relevant for resource control are a slice, a scope, and a service.

LOAD - indicates whether the unit configuration file was properly loaded. If the unit file failed to load, the field contains the state error instead of loaded. Other unit load states are: stub, merged, and masked.

ACTIVE - the high-level unit activation state, which is a generalization of SUB.

SUB - the low-level unit activation state. The range of possible values depends on the unit type.

DESCRIPTION - the description of the unit content and functionality.

To list inactive units, execute:

```
# systemctl --all
```

To limit the amount of information in the output, execute:

```
# systemctl --type service,masked
```

The --type option requires a comma-separated list of unit types such as a service and a slice, or unit load states such as loaded and masked.

Additional resources

- Configuring basic system settings in RHEL
- systemd.resource-control(5), systemd.exec(5) manual pages

15.5. VIEWING SYSTEMD CONTROL GROUP HIERARCHY

The following procedure describes how to display control groups (cgroups) hierarchy and processes running in specific cgroups.

Procedure

- To display the whole cgroups hierarchy on your system, execute # systemctl-cgls:

```
# systemctl-cgls
Control group /:
- .slice
  └─ user.slice
    └─ user-42.slice
      └─ session-c1.scope
        └─ 965 gdm-session-worker [pam/gdm-launch-environment]
        └─ 1040 /usr/libexec/gdm-x-session gnome-session --autostart
```
The example output returns the entire **cgroups** hierarchy, where the highest level is formed by **slices**.

- To display the **cgroups** hierarchy filtered by a resource controller, execute **# systemd-cgls**

```
# systemd-cgls memory
Controller memory; Control group /:
├─ 1 /usr/lib/systemd/systemd --switched-root --system --deserialize 18
├─ user.slice
│  ├─ user-42.slice
│  │  ├─ session-c1.scope
│  │  │  └─ 965 gdm-session-worker [pam/gdm-launch-environment]
│  └─ system.slice
│      ...
├─ chronyd.service
│  └─ 844 /usr/sbin/chronyd
└─ example.service
    ├─ 8914 /bin/bash /home/jdoe/example.sh
    └─ 8916 sleep 1
```

The example output of the above command lists the services that interact with the selected controller.

- To display detailed information about a certain unit and its part of the **cgroups** hierarchy, execute **# systemctl status**

```
# systemctl status example.service
● example.service - My example service
 Loaded: loaded (/usr/lib/systemd/system/example.service; enabled; vendor preset: disabled)
 Active: active (running) since Tue 2019-04-16 12:12:39 CEST; 3s ago
 Main PID: 17737 (bash)
 Tasks: 2 (limit: 11522)
 Memory: 496.0K (limit: 1.5M)
 CGroup: /system.slice/example.service
    └─ 17737 /bin/bash /home/jdoe/example.sh
    └─ 17743 sleep 1
```
15.6. VIEWING CGROUPS OF PROCESSES

The following procedure describes how to learn which control group (cgroup) a process belongs to. Then you can check the cgroup to learn which controllers and controller-specific configurations it uses.

Procedure

1. To view which cgroup a process belongs to, run the `cat /proc/<PID>/cgroup` command:

```
# cat /proc/2467/cgroup
0::/system.slice/example.service
```

The example output relates to a process of interest. In this case, it is a process identified by PID 2467, which belongs to the example.service unit. You can determine whether the process was placed in a correct control group as defined by the systemd unit file specifications.

2. To display what controllers the cgroup utilizes and the respective configuration files, check the cgroup directory:

```
# cat /sys/fs/cgroup/system.slice/example.service/cgroup.controllers
memory pids

# ls /sys/fs/cgroup/system.slice/example.service/
cgroup.controllers
cgroup.events
...
cpu.pressure
cpu.stat
io.pressure
memory.current
memory.events
...
pids.current
pids.events
pids.max
```

**NOTE**

The version 1 hierarchy of cgroups uses a per-controller model. Therefore the output from the `/proc/PID/cgroup` file shows, which cgroups under each controller the PID belongs to. You can find the respective cgroups under the controller directories at `/sys/fs/cgroup/<controller_name>/`. 
15.7. MONITORING RESOURCE CONSUMPTION

The following procedure describes how to view a list of currently running control groups (cgroups) and their resource consumption in real-time.

Procedure

1. To see a dynamic account of currently running cgroups, execute the `# systemd-cgtop` command:

```
# systemd-cgtop
Control Group     Tasks %CPU Memory Input/s Output/s
/                 607   29.8  1.5G   -   -
/system.slice     125   -428.7M  -   -   -
/system.slice/ModemManager.service 3     -8.6M  -   -   -
/system.slice/NetworkManager.service 3     -12.8M  -   -   -
/system.slice/accounts-daemon.service 3     -1.8M  -   -   -
/system.slice/boot.mount           -   -48.0K  -   -   -
/system.slice/chronyd.service     1     -2.0M  -   -   -
/system.slice/cockpit.socket      -   -1.3M  -   -   -
/system.slice/colord.service      3     -3.5M  -   -   -
/system.slice/crond.service       1     -1.8M  -   -   -
/system.slice/cups.service        1     -3.1M  -   -   -
/system.slice/dev-hugepages.mount -   -244.0K -   -   -
/system.slice/dev-mapper-rhel\x2dswap.swap -   -912.0K -   -   -
/system.slice/dev-mqueue.mount    -   -48.0K  -   -   -
/system.slice/example.service     2     -2.0M  -   -   -
/system.slice/firewalld.service   2     -28.8M -   -   -
```

The example output displays currently running cgroups ordered by their resource usage (CPU, memory, disk I/O load). The list refreshes every 1 second by default. Therefore, it offers a dynamic insight into the actual resource usage of each control group.

Additional resources

- `systemd-cgtop(1)` manual page

15.8. USING SYSTEMD UNIT FILES TO SET LIMITS FOR APPLICATIONS

Each existing or running unit is supervised by the systemd, which also creates control groups for them. The units have configuration files in the `/usr/lib/systemd/system/` directory. You can manually modify the unit files to set limits, prioritize, or control access to hardware resources for groups of processes.

Prerequisites
You have the root privileges.

Procedure

1. Modify the /usr/lib/systemd/system/example.service file to limit the memory usage of a service:

   ...
   [Service]
   MemoryMax=1500K
   ...

   The configuration above places a maximum memory limit, which the processes in a control group cannot exceed. The example.service service is part of such a control group which has imposed limitations. You can use suffixes K, M, G, or T to identify Kilobyte, Megabyte, Gigabyte, or Terabyte as a unit of measurement.

2. Reload all unit configuration files:

   # systemctl daemon-reload

3. Restart the service:

   # systemctl restart example.service

   NOTE

   You can review the full set of configuration options for systemd in the following manual pages:

   • systemd.resource-control(5)
   • systemd.exec(5)

Verification

1. Check that the changes took effect:

   # cat /sys/fs/cgroup/system.slice/example.service/memory.max
   1536000

   The example output shows that the memory consumption was limited at around 1,500 KB.

Additional resources

• Understanding cgroups

• Configuring basic system settings in Red Hat Enterprise Linux

• systemd.resource-control(5), systemd.exec(5), cgroups(7) manual pages

15.9. USING SYSTEMCTL COMMAND TO SET LIMITS TO APPLICATIONS
CPU affinity settings help you restrict the access of a particular process to some CPUs. Effectively, the CPU scheduler never schedules the process to run on the CPU that is not in the affinity mask of the process.

The default CPU affinity mask applies to all services managed by systemd.

To configure CPU affinity mask for a particular systemd service, systemd provides CPUAffinity= both as a unit file option and a manager configuration option in the /etc/systemd/system.conf file.

The CPUAffinity= unit file option sets a list of CPUs or CPU ranges that are merged and used as the affinity mask.

After configuring CPU affinity mask for a particular systemd service, you must restart the service to apply the changes.

Procedure

To set CPU affinity mask for a particular systemd service using the CPUAffinity unit file option:

1. Check the values of the CPUAffinity unit file option in the service of your choice:

   $ systemctl show --property <CPU affinity configuration option> <service name>

2. As a root, set the required value of the CPUAffinity unit file option for the CPU ranges used as the affinity mask:

   # systemctl set-property <service name> CPUAffinity=<value>

3. Restart the service to apply the changes.

   # systemctl restart <service name>

   **NOTE**

   You can review the full set of configuration options for systemd in the following manual pages:

   - systemd.resource-control(5)
   - systemd.exec(5)

15.10. SETTING GLOBAL DEFAULT CPU AFFINITY THROUGH MANAGER CONFIGURATION

The CPUAffinity option in the /etc/systemd/system.conf file defines an affinity mask for the process identification number (PID) 1 and all processes forked off of PID1. You can then override the CPUAffinity on a per-service basis.

To set default CPU affinity mask for all systemd services using the manager configuration option:

1. Set the CPU numbers for the CPUAffinity= option in the /etc/systemd/system.conf file.

2. Save the edited file and reload the systemd service:
# systemctl daemon-reload

3. Reboot the server to apply the changes.

NOTE
You can review the full set of configuration options for `systemd` in the following manual pages:

- `systemd.resource-control(5)`
- `systemd.exec(5)`

15.11. CONFIGURING NUMA POLICIES USING SYSTEMD

Non-uniform memory access (NUMA) is a computer memory subsystem design, in which the memory access time depends on the physical memory location relative to the processor.

Memory close to the CPU has lower latency (local memory) than memory that is local for a different CPU (foreign memory) or is shared between a set of CPUs.

In terms of the Linux kernel, NUMA policy governs where (for example, on which NUMA nodes) the kernel allocates physical memory pages for the process.

`systemd` provides unit file options `NUMAPolicy` and `NUMAMask` to control memory allocation policies for services.

Procedure
To set the NUMA memory policy through the `NUMAPolicy` unit file option:

1. Check the values of the `NUMAPolicy` unit file option in the service of your choice:

   ```
   $ systemctl show --property <NUMA policy configuration option> <service name>
   ```

2. As a root, set the required policy type of the `NUMAPolicy` unit file option:

   ```
   # systemctl set-property <service name> NUMAPolicy=<value>
   ```

3. Restart the service to apply the changes.

   ```
   # systemctl restart <service name>
   ```

To set a global `NUMAPolicy` setting through the `manager configuration` option:

1. Search in the `/etc/systemd/system.conf` file for the `NUMAPolicy` option.

2. Edit the policy type and save the file.

3. Reload the `systemd` configuration:

   ```
   # systemctl daemon-reload
   ```

4. Reboot the server.
IMPORTANT

When you configure a strict NUMA policy, for example `bind`, make sure that you also appropriately set the `CPUAffinity` unit file option.

Additional resources

- Using `systemctl` command to set limits to applications
- `systemd.resource-control(5)`, `systemd.exec(5)`, `set_mempolicy(2)` manual pages.

15.12. NUMA POLICY CONFIGURATION OPTIONS FOR SYSTEMD

Systemd provides the following options to configure the NUMA policy:

**NUMAPolicy**

Controls the NUMA memory policy of the executed processes. The following policy types are possible:

- default
- preferred
- bind
- interleave
- local

**NUMAMask**

Controls the NUMA node list which is associated with the selected NUMA policy.

Note that the `NUMAMask` option is not required to be specified for the following policies:

- default
- local

For the preferred policy, the list specifies only a single NUMA node.

Additional resources

- `systemd.resource-control(5)`, `systemd.exec(5)`, and `set_mempolicy(2)` manual pages

15.13. CREATING TRANSIENT CGROUPS USING SYSTEMD-RUN COMMAND

The transient `cgroups` set limits on resources consumed by a unit (service or scope) during its runtime.

**Procedure**

- To create a transient control group, use the `systemd-run` command in the following format:

  ```shell
  # systemd-run --unit=<name> --slice=<name>.slice <command>
  ```
This command creates and starts a transient service or a scope unit and runs a custom command in such a unit.

- The \texttt{--unit=<name>} option gives a name to the unit. If \texttt{--unit} is not specified, the name is generated automatically.

- The \texttt{--slice=<name>.slice} option makes your service or scope unit a member of a specified slice. Replace \texttt{<name>.slice} with the name of an existing slice (as shown in the output of \texttt{systemctl -t slice}), or create a new slice by passing a unique name. By default, services and scopes are created as members of the \texttt{system.slice}.

- Replace \texttt{<command>} with the command you wish to execute in the service or the scope unit.

The following message is displayed to confirm that you created and started the service or the scope successfully:

\begin{verbatim}
# Running as unit <name>.service
\end{verbatim}

- Optionally, keep the unit running after its processes finished to collect run-time information:

\begin{verbatim}
# systemd-run --unit=<name> --slice=<name>.slice --remain-after-exit <command>
\end{verbatim}

The command creates and starts a transient service unit and runs a custom command in such a unit. The \texttt{--remain-after-exit} option ensures that the service keeps running after its processes have finished.

Additional resources

- Understanding control groups
- \texttt{Configuring basic system settings} in RHEL
- the \texttt{systemd-run(1)} manual page

15.14. REMOVING TRANSIENT CONTROL GROUPS

You can use the \texttt{systemd} system and service manager to remove transient control groups (\texttt{cgroups}) if you no longer need to limit, prioritize, or control access to hardware resources for groups of processes.

Transient \texttt{cgroups} are automatically released once all the processes that a service or a scope unit contains, finish.

Procedure

- To stop the service unit with all its processes, execute:

\begin{verbatim}
# systemctl stop name.service
\end{verbatim}

- To terminate one or more of the unit processes, execute:

\begin{verbatim}
# systemctl kill name.service --kill-who=PID,... --signal=<signal>
\end{verbatim}

The command above uses the \texttt{--kill-who} option to select process(es) from the control group.
you wish to terminate. To kill multiple processes at the same time, pass a comma-separated list of PIDs. The **--signal** option determines the type of POSIX signal to be sent to the specified processes. The default signal is *SIGTERM*.

**Additional resources**

- Understanding control groups
- What are kernel resource controllers
- `systemd.resource-control(5), cgroups(7)` manual pages
- *Configuring basic system settings* in RHEL
CHAPTER 16. UNDERSTANDING CGROUPS

You can use the control groups (cgroups) kernel functionality to set limits, prioritize or isolate the hardware resources of processes. This allows you to granularly control resource usage of applications to utilize them more efficiently.

16.1. UNDERSTANDING CONTROL GROUPS

Control groups is a Linux kernel feature that enables you to organize processes into hierarchically ordered groups - cgroups. The hierarchy (control groups tree) is defined by providing structure to cgroups virtual file system, mounted by default on the /sys/fs/cgroup/ directory. The systemd system and service manager utilizes cgroups to organize all units and services that it governs. Alternatively, you can manage cgroups hierarchies manually by creating and removing sub-directories in the /sys/fs/cgroup/ directory.

The resource controllers (a kernel component) then modify the behavior of processes in cgroups by limiting, prioritizing or allocating system resources, (such as CPU time, memory, network bandwidth, or various combinations) of those processes.

The added value of cgroups is process aggregation which enables division of hardware resources among applications and users. Thereby an increase in overall efficiency, stability and security of users’ environment can be achieved.

Control groups version 1

Control groups version 1 (cgroups-v1) provide a per-resource controller hierarchy. It means that each resource, such as CPU, memory, I/O, and so on, has its own control group hierarchy. It is possible to combine different control group hierarchies in a way that one controller can coordinate with another one in managing their respective resources. However, the two controllers may belong to different process hierarchies, which does not permit their proper coordination.

The cgroups-v1 controllers were developed across a large time span and as a result, the behavior and naming of their control files is not uniform.

Control groups version 2

The problems with controller coordination, which stemmed from hierarchy flexibility, led to the development of control groups version 2.

Control groups version 2 (cgroups-v2) provides a single control group hierarchy against which all resource controllers are mounted.

The control file behavior and naming is consistent among different controllers.

IMPORTANT

RHEL 9, by default, mounts and utilizes cgroups-v2.

This sub-section was based on a Devconf.cz 2019 presentation.[1]

Additional resources

- What are kernel resource controllers
- cgroups(7) manual page
16.2. WHAT ARE KERNEL RESOURCE CONTROLLERS

The functionality of control groups is enabled by kernel resource controllers. RHEL 9 supports various controllers for control groups version 1 (cgroups-v1) and control groups version 2 (cgroups-v2).

A resource controller, also called a control group subsystem, is a kernel subsystem that represents a single resource, such as CPU time, memory, network bandwidth or disk I/O. The Linux kernel provides a range of resource controllers that are mounted automatically by the systemd system and service manager. Find a list of currently mounted resource controllers in the /proc/cgroups file.

The following controllers are available for cgroups-v1:

- **blkio** - can set limits on input/output access to and from block devices.
- **cpu** - can adjust the parameters of the Completely Fair Scheduler (CFS) scheduler for control group’s tasks. It is mounted together with the cpuacct controller on the same mount.
- **cpuacct** - creates automatic reports on CPU resources used by tasks in a control group. It is mounted together with the cpu controller on the same mount.
- **cpuset** - can be used to restrict control group tasks to run only on a specified subset of CPUs and to direct the tasks to use memory only on specified memory nodes.
- **devices** - can control access to devices for tasks in a control group.
- **freezer** - can be used to suspend or resume tasks in a control group.
- **memory** - can be used to set limits on memory use by tasks in a control group and generates automatic reports on memory resources used by those tasks.
- **net_cls** - tags network packets with a class identifier (classid) that enables the Linux traffic controller (the tc command) to identify packets that originate from a particular control group task. A subsystem of net_cls, the net_filter (iptables), can also use this tag to perform actions on such packets. The net_filter tags network sockets with a firewall identifier (fwid) that allows the Linux firewall (through iptables command) to identify packets originating from a particular control group task.
- **net_prio** - sets the priority of network traffic.
- **pids** - can set limits for a number of processes and their children in a control group.
- **perf_event** - can group tasks for monitoring by the perf performance monitoring and reporting utility.
- **rdma** - can set limits on Remote Direct Memory Access/InfiniBand specific resources in a control group.
- **hugetlb** - can be used to limit the usage of large size virtual memory pages by tasks in a control group.

The following controllers are available for cgroups-v2:

- **io** - A follow-up to blkio of cgroups-v1.
• **memory** - A follow-up to **memory** of **cgroups-v1**.

• **pids** - Same as **pids** in **cgroups-v1**.

• **rdma** - Same as **rdma** in **cgroups-v1**.

• **cpu** - A follow-up to **cpu** and **cpuacct** of **cgroups-v1**.

• **cpuset** - Supports only the core functionality (**cpus{,.effective}, mems{,.effective}**) with a new partition feature.

• **perf_event** - Support is inherent, no explicit control file. You can specify a **v2 cgroup** as a parameter to the **perf** command that will profile all the tasks within that **cgroup**.

**IMPORTANT**

A resource controller can be used either in a **cgroups-v1** hierarchy or a **cgroups-v2** hierarchy, not simultaneously in both.

**Additional resources**

• **cgroups(7)** manual page


16.3. WHAT ARE NAMESPACES

Namespaces are one of the most important methods for organizing and identifying software objects.

A namespace wraps a global system resource (for example a mount point, a network device, or a hostname) in an abstraction that makes it appear to processes within the namespace that they have their own isolated instance of the global resource. One of the most common technologies that utilize namespaces are containers.

Changes to a particular global resource are visible only to processes in that namespace and do not affect the rest of the system or other namespaces.

To inspect which namespaces a process is a member of, you can check the symbolic links in the `/proc/<PID>/ns/` directory.

The following table shows supported namespaces and resources which they isolate:

<table>
<thead>
<tr>
<th>Namespace</th>
<th>Isolates</th>
</tr>
</thead>
<tbody>
<tr>
<td>Mount</td>
<td>Mount points</td>
</tr>
<tr>
<td>UTS</td>
<td>Hostname and NIS domain name</td>
</tr>
<tr>
<td>IPC</td>
<td>System V IPC, POSIX message queues</td>
</tr>
<tr>
<td>PID</td>
<td>Process IDs</td>
</tr>
<tr>
<td>Namespace</td>
<td>Isolates</td>
</tr>
<tr>
<td>---------------------</td>
<td>-----------------------------------------------</td>
</tr>
<tr>
<td>Network</td>
<td>Network devices, stacks, ports, etc</td>
</tr>
<tr>
<td>User</td>
<td>User and group IDs</td>
</tr>
<tr>
<td>Control groups</td>
<td>Control group root directory</td>
</tr>
</tbody>
</table>

Additional resources

- namespaces(7) and cgroup_namespaces(7) manual pages
- Understanding control groups

[1] Linux Control Group v2 – An Introduction, Devconf.cz 2019 presentation by Waiman Long
You can improve system performance by enabling the \texttt{zswap} kernel feature.

### 17.1. WHAT IS ZSWAP

This section explains what \texttt{zswap} is and how it can lead to system performance improvement.

\texttt{zswap} is a kernel feature that provides a compressed RAM cache for swap pages. The mechanism works as follows: \texttt{zswap} takes pages that are in the process of being swapped out and attempts to compress them into a dynamically allocated RAM-based memory pool. When the pool becomes full or the RAM becomes exhausted, \texttt{zswap} evicts pages from compressed cache on an LRU basis (least recently used) to the backing swap device. After the page has been decompressed into the swap cache, \texttt{zswap} frees the compressed version in the pool.

The benefits of \texttt{zswap}:

- significant I/O reduction
- significant improvement of workload performance

In Red Hat Enterprise Linux 9, \texttt{zswap} is enabled by default.

**Additional resources**

- What is Zswap?

### 17.2. ENABLING ZSWAP AT RUNTIME

You can enable the \texttt{zswap} feature at system runtime using the \texttt{sysfs} interface.

**Prerequisites**

- You have root permissions.

**Procedure**

- Enable \texttt{zswap}:

  ```
  # echo 1 > /sys/module/zswap/parameters/enabled
  ```

**Verification step**

- Verify that \texttt{zswap} is enabled:

  ```
  # grep -r . /sys/kernel/debug/zswap
  duplicate_entry:0
  pool_limit_hit:13422200
  pool_total_size:6184960 (pool size in total in pages)
  reject_alloc_fail:5
  ```
17.3. ENABLING ZSWAP PERMANENTLY

You can enable the `zswap` feature permanently by providing the `zswap.enabled=1` kernel command-line parameter.

Prerequisites

- You have root permissions.
- The `grubby` or `zipl` utility is installed on your system.

Procedure

1. Enable `zswap` permanently:

   ```
   # grubby --update-kernel=/boot/vmlinuz-$(uname -r) --args="zswap.enabled=1"
   ```

2. Reboot the system for the changes to take effect.

Verification steps

- Verify that `zswap` is enabled:

   ```
   # cat /proc/cmdline
   ```

   ```
   BOOT_IMAGE=(hd0,msdos1)/vmlinuz-5.14.0-70.5.1.el9_0.x86_64
   root=/dev/mapper/rhel-root ro crashkernel=1G-4G:192M,4G-64G:256M,64G-:512M
   resume=/dev/mapper/rhel-swap rd.lvm.lv=rhel/root
   rd.lvm.lv=rhel/swap rhgb quiet
   zswap.enabled=1
   ```

Additional resources

- How to enable Zswap feature?
- Configuring kernel command-line parameters
CHAPTER 18. USING CGROUPFS TO MANUALLY MANAGE CGROUPS

You can manage cgroup hierarchies on your system by creating directories on the cgroupfs virtual file system. The file system is mounted by default on the /sys/fs/cgroup/ directory and you can specify desired configurations in dedicated control files.

IMPORTANT

In general, Red Hat recommends you use systemd for controlling the usage of system resources. You should manually configure the cgroups virtual file system only in special cases. For example, when you need to use cgroup-v1 controllers that have no equivalents in cgroup-v2 hierarchy.

18.1. CREATING CGROUPS AND ENABLING CONTROLLERS IN CGROUPS-V2 FILE SYSTEM

You can manage the control groups (cgroups) by creating or removing directories and by writing to files in the cgroups virtual file system. The file system is by default mounted on the /sys/fs/cgroup/ directory. To use settings from the cgroups controllers, you also need to enable the desired controllers for child cgroups. The root cgroup has, by default, enabled the memory and pids controllers for its child cgroups. Therefore, Red Hat recommends to create at least two levels of child cgroups inside the /sys/fs/cgroup/ root cgroup. This way you optionally remove the memory and pids controllers from the child cgroups and maintain better organizational clarity of cgroup files.

Prerequisites

- You have root permissions.

Procedure

1. Create the /sys/fs/cgroup/Example/ directory:

   # mkdir /sys/fs/cgroup/Example/

   The /sys/fs/cgroup/Example/ directory defines a child group. When you create the /sys/fs/cgroup/Example/ directory, some cgroups-v2 interface files are automatically created in the directory. The /sys/fs/cgroup/Example/ directory contains also controller-specific files for the memory and pids controllers.

2. Optionally, inspect the newly created child control group:

   # ll /sys/fs/cgroup/Example/
   -r--r--r-- 1 root root 0 Jun  1 10:33 cgroup.controllers
   -r--r--r-- 1 root root 0 Jun  1 10:33 cgroup.events
   -rw-r--r-- 1 root root 0 Jun  1 10:33 cgroup.freeze
   -rw-r--r-- 1 root root 0 Jun  1 10:33 cgroup.procs
   ...
   -rw-r--r-- 1 root root 0 Jun  1 10:33 cgroup.subtree_control
   -r--r--r-- 1 root root 0 Jun  1 10:33 memory.events.local
   -rw-r--r-- 1 root root 0 Jun  1 10:33 memory.high
   -rw-r--r-- 1 root root 0 Jun  1 10:33 memory.low
   ...

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The example output shows general `cgroup` control interface files such as `cgroup.procs` or `cgroup.controllers`. These files are common to all control groups, regardless of enabled controllers.

The files such as `memory.high` and `pids.max` relate to the `memory` and `pids` controllers, which are in the root control group (`/sys/fs/cgroup/`), and are enabled by default by `systemd`.

By default, the newly created child group inherits all settings from the parent `cgroup`. In this case, there are no limits from the root `cgroup`.

3. Verify that the desired controllers are available in the `/sys/fs/cgroup/cgroup.controllers` file:

```
# cat /sys/fs/cgroup/cgroup.controllers
cpuset cpu io memory hugetlb pids rdma
```

4. Enable the desired controllers. In this example it is `cpu` and `cpuset` controllers:

```
# echo "+cpu" >> /sys/fs/cgroup/cgroup.subtree_control
# echo "+cpuset" >> /sys/fs/cgroup/cgroup.subtree_control
```

These commands enable the `cpu` and `cpuset` controllers for the immediate child groups of the `/sys/fs/cgroup/` root control group. Including the newly created `Example` control group. A `child group` is where you can specify processes and apply control checks to each of the processes based on your criteria.

Users can read the contents of the `cgroup.subtree_control` file at any level to get an idea of what controllers are going to be available for enablement in the immediate child group.

**NOTE**

By default, the `/sys/fs/cgroup/cgroup.subtree_control` file in the root control group contains `memory` and `pids` controllers.

5. Enable the desired controllers for child `cgroups` of the `Example` control group:

```
# echo "+cpu +cpuset" >> /sys/fs/cgroup/Example/cgroup.subtree_control
```

This command ensures that the immediate child control group will only have controllers relevant to regulate the CPU time distribution - not to `memory` or `pids` controllers.

6. Create the `/sys/fs/cgroup/Example/tasks/` directory:

```
# mkdir /sys/fs/cgroup/Example/tasks/
```

The `/sys/fs/cgroup/Example/tasks/` directory defines a child group with files that relate purely to `cpu` and `cpuset` controllers. You can now assign processes to this control group and utilize `cpu` and `cpuset` controller options for your processes.

7. Optionally, inspect the child control group:
# ll /sys/fs/cgroup/Example/tasks
-r--r--r-- 1 root root 0 Jun 11:45 cgroup.controllers
-r--r--r-- 1 root root 0 Jun 11:45 cgroup.events
-rw-r--r-- 1 root root 0 Jun 11:45 cgroup.freeze
-rw-r--r-- 1 root root 0 Jun 11:45 cgroup.max.depth
-rw-r--r-- 1 root root 0 Jun 11:45 cgroup.max.descendants
-rw-r--r-- 1 root root 0 Jun 11:45 cgroup.procs
-r--r--r-- 1 root root 0 Jun 11:45 cgroup.stat
-rw-r--r-- 1 root root 0 Jun 11:45 cgroup.subtree_control
-rw-r--r-- 1 root root 0 Jun 11:45 cgroup.type
-rw-r--r-- 1 root root 0 Jun 11:45 cpu.max
-rw-r--r-- 1 root root 0 Jun 11:45 cpu.pressure
-rw-r--r-- 1 root root 0 Jun 11:45 cpuset.cpus
-r--r--r-- 1 root root 0 Jun 11:45 cpuset.cpus.effective
-rw-r--r-- 1 root root 0 Jun 11:45 cpuset.cpus.partition
-rw-r--r-- 1 root root 0 Jun 11:45 cpuset.mems
-r--r--r-- 1 root root 0 Jun 11:45 cpuset.mems.effective
-r--r--r-- 1 root root 0 Jun 11:45 cpu.stat
-rw-r--r-- 1 root root 0 Jun 11:45 cpu.weight
-rw-r--r-- 1 root root 0 Jun 11:45 cpu.weight.nice
-rw-r--r-- 1 root root 0 Jun 11:45 io.pressure
-rw-r--r-- 1 root root 0 Jun 11:45 memory.pressure

**IMPORTANT**

The **cpu** controller is only activated if the relevant child control group has at least 2 processes which compete for time on a single CPU.

**Verification steps**

- Optional: confirm that you have created a new **cgroup** with only the desired controllers active:

  ```
  # cat /sys/fs/cgroup/Example/tasks/cgroup.controllers
  cpuset cpu
  ```

**Additional resources**

- **Understanding control groups**
- **What are kernel resource controllers**
- **Mounting cgroups-v1**
- **cgroups(7), sysfs(5) manual pages**

### 18.2. CONTROLLING DISTRIBUTION OF CPU TIME FOR APPLICATIONS BY ADJUSTING CPU WEIGHT

You need to assign values to the relevant files of the **cpu** controller to regulate distribution of the CPU time to applications under the specific cgroup tree.

**Prerequisites**
- You have root permissions.
- You have applications for which you want to control distribution of CPU time.
- You created a two level hierarchy of child control groups inside the /sys/fs/cgroup/ root control group as in the following example:

```
...  
|    Example
|    ├── g1
|    │    ├── g2
|    │    │    └── g3

```

- You enabled the cpu controller in the parent control group and in child control groups similarly as described in Creating cgroups and enabling controllers in cgroups-v2 file system.

**Procedure**

1. Configure desired CPU weights to achieve resource restrictions within the control groups:

```
# echo "150" > /sys/fs/cgroup/Example/g1/cpu.weight
# echo "100" > /sys/fs/cgroup/Example/g2/cpu.weight
# echo "50" > /sys/fs/cgroup/Example/g3/cpu.weight
```

2. Add the applications' PIDs to the g1, g2, and g3 child groups:

```
# echo "33373" > /sys/fs/cgroup/Example/g1/cgroup.procs
# echo "33374" > /sys/fs/cgroup/Example/g2/cgroup.procs
# echo "33377" > /sys/fs/cgroup/Example/g3/cgroup.procs
```

The example commands ensure that desired applications become members of the Example/g*/ child cgroups and will get their CPU time distributed as per the configuration of those cgroups.

The weights of the children cgroups (g1, g2, g3) that have running processes are summed up at the level of the parent cgroup (Example). The CPU resource is then distributed proportionally based on the respective weights.

As a result, when all processes run at the same time, the kernel allocates to each of them the proportionate CPU time based on their respective cgroup’s cpu.weight file:

<table>
<thead>
<tr>
<th>Child cgroup</th>
<th>cpu.weight file</th>
<th>CPU time allocation</th>
</tr>
</thead>
<tbody>
<tr>
<td>g1</td>
<td>150</td>
<td>~50% (150/300)</td>
</tr>
<tr>
<td>g2</td>
<td>100</td>
<td>~33% (100/300)</td>
</tr>
<tr>
<td>g3</td>
<td>50</td>
<td>~16% (50/300)</td>
</tr>
</tbody>
</table>

The value of the cpu.weight controller file is not a percentage.
If one process stopped running, leaving cgroup g2 with no running processes, the calculation would omit the cgroup g2 and only account weights of cgroups g1 and g3:

<table>
<thead>
<tr>
<th>Child cgroup</th>
<th>cpu.weight file</th>
<th>CPU time allocation</th>
</tr>
</thead>
<tbody>
<tr>
<td>g1</td>
<td>150</td>
<td>~75% (150/200)</td>
</tr>
<tr>
<td>g3</td>
<td>50</td>
<td>~25% (50/200)</td>
</tr>
</tbody>
</table>

**IMPORTANT**

If a child cgroup had multiple running processes, the CPU time allocated to the respective cgroup would be distributed equally to the member processes of that cgroup.

**Verification**

1. Verify that the applications run in the specified control groups:

   ```
   # cat /proc/33373/cgroup /proc/33374/cgroup /proc/33377/cgroup
   0::/Example/g1
   0::/Example/g2
   0::/Example/g3
   ```

   The command output shows the processes of the specified applications that run in the Example/g*/ child cgroups.

2. Inspect the current CPU consumption of the throttled applications:

   ```
   # top
   top - 05:17:18 up 1 day, 18:25, 1 user, load average: 3.03, 3.03, 3.00
   Tasks: 95 total, 4 running, 91 sleeping, 0 stopped, 0 zombie
   %Cpu(s): 18.1 us, 81.6 sy, 0.0 ni, 0.0 id, 0.0 wa, 0.3 hi, 0.0 si, 0.0 st
   MiB Mem: 3737.0 total, 3233.7 free, 132.8 used, 370.5 buff/cache
   MiB Swap: 4060.0 total, 4060.0 free, 0.0 used. 3373.1 avail Mem
   
   PID USER PR NI VIRT RES SHR %CPU %MEM TIME+ COMMAND
   33373 root 20 0 18720 1748 1460 R 49.5 0.0 415:05.87 sha1sum
   33374 root 20 0 18720 1756 1464 R 32.9 0.0 412:58.33 sha1sum
   33377 root 20 0 18720 1860 1568 R 16.3 0.0 411:03.12 sha1sum
   760 root 20 0 416620 28540 15296 S 0.3 0.7 0:10:23 tuned
   1 root 20 0 186328 14108 9484 S 0.0 0.4 0:02.00 systemd
   2 root 20 0 0 0 0 S 0.0 0.0 0:00.01 kthread
   ...  
   ```

**NOTE**

We forced all the example processes to run on a single CPU for clearer illustration. The CPU weight applies the same principles also when used on multiple CPUs.
Notice that the CPU resource for the **PID 33373**, **PID 33374**, and **PID 33377** was allocated based on the weights, 150, 100, 50, you assigned to the respective child cgroups. The weights correspond to around 50%, 33%, and 16% allocation of CPU time for each application.

### Additional resources

- Understanding control groups
- What are kernel resource controllers
- Creating cgroups and enabling controllers in cgroups-v2 file system
- Resource Distribution Models
- cgroups(7), sysfs(5) manual pages

### 18.3. MOUNTING CGROUPS-V1

During the boot process, RHEL 9 mounts the cgroup-v2 virtual filesystem by default. To utilize cgroup-v1 functionality in limiting resources for your applications, manually configure the system.

**NOTE**

Both cgroup-v1 and cgroup-v2 are fully enabled in the kernel. There is no default control group version from the kernel point of view, and is decided by systemd to mount at startup.

#### Prerequisites

- You have root permissions.

#### Procedure

1. Configure the system to mount cgroups-v1 by default during system boot by the systemd system and service manager:
   ```bash
   # grubby --update-kernel=/boot/vmlinuz-$(uname -r) --
   args="systemd.unified_cgroup_hierarchy=0
   systemd.legacy_systemd_cgroup_controller"
   ```
   This adds the necessary kernel command-line parameters to the current boot entry.

   To add the same parameters to all kernel boot entries:
   ```bash
   # grubby --update-kernel=ALL --args="systemd.unified_cgroup_hierarchy=0
   systemd.legacy_systemd_cgroup_controller"
   ```

2. Reboot the system for the changes to take effect.

#### Verification

1. Optionally, verify that the cgroups-v1 filesystem was mounted:
   ```bash
   # mount -l | grep cgroup
   ```
The cgroups-v1 filesystems that correspond to various cgroup-v1 controllers, were successfully mounted on the /sys/fs/cgroup/ directory.

2. Optionally, inspect the contents of the /sys/fs/cgroup/ directory:

```
# ll /sys/fs/cgroup/
```

dr-xr-xr-x. 10 root root 0 Mar 16 09:34 blkio
lrwxrwxrwx. 10 root root 0 Mar 16 09:34 cpu → cpu,cpuacct
lrwxrwxrwx. 10 root root 0 Mar 16 09:34 cpuacct → cpu,cpuacct
dr-xr-xr-x. 10 root root 0 Mar 16 09:34 cpu,cpuacct
dr-xr-xr-x. 2 root root 0 Mar 16 09:34 cpuset
dr-xr-xr-x. 10 root root 0 Mar 16 09:34 devices
dr-xr-xr-x. 2 root root 0 Mar 16 09:34 freezer
dr-xr-xr-x. 2 root root 0 Mar 16 09:34 hugetlb
dr-xr-xr-x. 10 root root 0 Mar 16 09:34 memory
dr-xr-xr-x. 2 root root 0 Mar 16 09:34 misc
lrwxrwxrwx. 10 root root 0 Mar 16 09:34 net_cls → net_cls,net_prio
dr-xr-xr-x. 2 root root 0 Mar 16 09:34 net_cls,net_prio
lrwxrwxrwx. 10 root root 0 Mar 16 09:34 net_cls,net_prio
lrwxrwxrwx. 10 root root 0 Mar 16 09:34 perf_event
dr-xr-xr-x. 10 root root 0 Mar 16 09:34 pids
dr-xr-xr-x. 2 root root 0 Mar 16 09:34 rdma
dr-xr-xr-x. 11 root root 0 Mar 16 09:34 systemd
```

The /sys/fs/cgroup/ directory, also called the root control group, by default, contains controller-specific directories such as cpuset. In addition, there are some directories related to systemd.

Additional resources

- Understanding control groups
What are kernel resource controllers

cgroups(7), sysfs(5) manual pages

cgroup-v2 enabled by default in RHEL 9

18.4. SETTING CPU LIMITS TO APPLICATIONS USING CGROUPS-V1

Sometimes an application consumes a lot of CPU time, which may negatively impact the overall health of your environment. Use the /sys/fs/ virtual file system to configure CPU limits to an application using control groups version 1 (cgroups-v1).

Prerequisites

- You have root permissions.
- You have an application whose CPU consumption you want to restrict.
- You configured the system to mount cgroups-v1 by default during system boot by the systemctl system and service manager:

```bash
# grubby --update-kernel=/boot/vmlinuz-$uname -r --
args="systemd.unified_cgroup_hierarchy=0
systemd.legacy_systemd_cgroup_controller"
```

This adds the necessary kernel command-line parameters to the current boot entry.

Procedure

1. Identify the process ID (PID) of the application you want to restrict in CPU consumption:

```bash
# top
```

Tasks: 267 total, 3 running, 264 sleeping, 0 stopped, 0 zombie
%Cpu(s): 49.0 us, 3.3 sy, 0.0 ni, 47.5 id, 0.0 wa, 0.2 hi, 0.0 si, 0.0 st
MiB Mem : 1826.8 total, 303.4 free, 1046.8 used, 476.5 buff/cache
MiB Swap: 1536.0 total, 1396.0 free, 140.0 used. 616.4 avail Mem

<table>
<thead>
<tr>
<th>PID</th>
<th>USER</th>
<th>PR</th>
<th>NI</th>
<th>VIRT</th>
<th>RES</th>
<th>SHR</th>
<th>S</th>
<th>%CPU</th>
<th>%MEM</th>
<th>TIME+</th>
<th>COMMAND</th>
</tr>
</thead>
<tbody>
<tr>
<td>6955</td>
<td>root</td>
<td>20</td>
<td>0</td>
<td>228440</td>
<td>1752</td>
<td>1472</td>
<td>R</td>
<td>99.3</td>
<td>0.1</td>
<td>0:32.71</td>
<td>sha1sum</td>
</tr>
<tr>
<td>5760</td>
<td>jdoe</td>
<td>20</td>
<td>0</td>
<td>3603868</td>
<td>205188</td>
<td>64196</td>
<td>S</td>
<td>3.7</td>
<td>11.0</td>
<td>0:17:19</td>
<td>gnome-shell</td>
</tr>
<tr>
<td>6448</td>
<td>jdoe</td>
<td>20</td>
<td>0</td>
<td>743648</td>
<td>30640</td>
<td>19488</td>
<td>S</td>
<td>0.7</td>
<td>1.6</td>
<td>0:02.73</td>
<td>gnome-terminal-1</td>
</tr>
<tr>
<td>1</td>
<td>root</td>
<td>20</td>
<td>0</td>
<td>245300</td>
<td>6568</td>
<td>4116</td>
<td>S</td>
<td>0.3</td>
<td>0.4</td>
<td>0:01.87</td>
<td>systemd</td>
</tr>
<tr>
<td>505</td>
<td>root</td>
<td>20</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>l</td>
<td>0.3</td>
<td>0.0</td>
<td>0:00.75</td>
<td>kworker/u4:4-events_unbound</td>
</tr>
</tbody>
</table>

The example output of the `top` program reveals that PID 6955 (illustrative application sha1sum) consumes a lot of CPU resources.

2. Create a sub-directory in the cpu resource controller directory:

```bash
# mkdir /sys/fs/cgroup/cpu/Example/
```
The directory above represents a control group, where you can place specific processes and apply certain CPU limits to the processes. At the same time, some cgroups-v1 interface files and cpu controller-specific files will be created in the directory.

3. Optionally, inspect the newly created control group:

```bash
# ll /sys/fs/cgroup/cpu/Example/
-rw-r--r-- 1 root root 0 Mar 11 11:42 cgroup.clone_children
-rw-r--r-- 1 root root 0 Mar 11 11:42 cgroup.procs
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.stat
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.usage
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.usage_all
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.usage_percpu
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.usage_percpu_sys
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.usage_percpu_user
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.usage_sys
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpuacct.usage_user
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpu.cfs_period_us
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpu.cfs_quota_us
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpu.rt_period_us
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpu.rt_runtime_us
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpu.shares
-rw-r--r-- 1 root root 0 Mar 11 11:42 cpu.stat
-rw-r--r-- 1 root root 0 Mar 11 11:42 notify_on_release
-rw-r--r-- 1 root root 0 Mar 11 11:42 tasks
```

The example output shows files, such as `cpuacct.usage`, `cpu.cfs_period_us`, that represent specific configurations and/or limits, which can be set for processes in the Example control group. Notice that the respective file names are prefixed with the name of the control group controller to which they belong.

By default, the newly created control group inherits access to the system’s entire CPU resources without a limit.

4. Configure CPU limits for the control group:

```bash
# echo "1000000" > /sys/fs/cgroup/cpu/Example/cpu.cfs_period_us
# echo "200000" > /sys/fs/cgroup/cpu/Example/cpu.cfs_quota_us
```

The `cpu.cfs_period_us` file represents a period of time in microseconds (µs, represented here as "us") for how frequently a control group’s access to CPU resources should be reallocated. The upper limit is 1 second and the lower limit is 1000 microseconds.

The `cpu.cfs_quota_us` file represents the total amount of time in microseconds for which all processes collectively in a control group can run during one period (as defined by `cpu.cfs_period_us`). As soon as processes in a control group, during a single period, use up all the time specified by the quota, they are throttled for the remainder of the period and not allowed to run until the next period. The lower limit is 1000 microseconds.

The example commands above set the CPU time limits so that all processes collectively in the Example control group will be able to run only for 0.2 seconds (defined by `cpu.cfs_quota_us`) out of every 1 second (defined by `cpu.cfs_period_us`).

5. Optionally, verify the limits:

```bash
# cat /sys/fs/cgroup/cpu/Example/cpu.cfs_period_us
```
6. Add the application’s PID to the Example control group:

```bash
# echo "6955" > /sys/fs/cgroup/cpu/Example/cgroup.procs
```

or

```bash
# echo "6955" > /sys/fs/cgroup/cpu/Example/tasks
```

The previous command ensures that a desired application becomes a member of the Example control group and hence does not exceed the CPU limits configured for the Example control group. The PID should represent an existing process in the system. The PID 6955 here was assigned to process `sha1sum /dev/zero &`, used to illustrate the use-case of the `cpu` controller.

7. Verify that the application runs in the specified control group:

```bash
# cat /proc/6955/cgroup
12:cpuset:/
11:hugetlb:/
10:net_cls,net_prio:/
 9:memory:/user.slice/user-1000.slice/user@1000.service
 8:devices:/user.slice
 7:blkio:/
 6:freezer:/
 5:rdma:/
 4:pids:/user.slice/user-1000.slice/user@1000.service
 3:perf_event:/
 2:cpu,cpuacct:/Example
 1:name=systemd:/user.slice/user-1000.slice/user@1000.service/gnome-terminal-server.service
```

The example output above shows that the process of the desired application runs in the Example control group, which applies CPU limits to the application’s process.

8. Identify the current CPU consumption of your throttled application:

```bash
# top
```

Notice that the CPU consumption of the PID 6955 has decreased from 99% to 20%.
IMPORTANT

The `cgroups-v2` counterpart for `cpu.cfs_period_us` and `cpu.cfs_quota_us` is the `cpu.max` file. The `cpu.max` file is available through the `cpu` controller.

Additional resources

- Understanding control groups
- What kernel resource controllers are
- `cgroups(7), sysfs(5)` manual pages
CHAPTER 19. ANALYZING SYSTEM PERFORMANCE WITH BPF COMPILER COLLECTION

As a system administrator, you can use the BPF Compiler Collection (BCC) library to create tools for analyzing the performance of your Linux operating system and gathering information, which could be difficult to obtain through other interfaces.

19.1. AN INTRODUCTION TO BCC

BPF Compiler Collection (BCC) is a library, which facilitates the creation of the extended Berkeley Packet Filter (eBPF) programs. The main utility of eBPF programs is analyzing OS performance and network performance without experiencing overhead or security issues.

BCC removes the need for users to know deep technical details of eBPF, and provides many out-of-the-box starting points, such as the bcc-tools package with pre-created eBPF programs.

NOTE

The eBPF programs are triggered on events, such as disk I/O, TCP connections, and process creations. It is unlikely that the programs should cause the kernel to crash, loop or become unresponsive because they run in a safe virtual machine in the kernel.

19.2. INSTALLING THE BCC-TOOLS PACKAGE

This section describes how to install the bcc-tools package, which also installs the BPF Compiler Collection (BCC) library as a dependency.

Procedure

1. Install bcc-tools:

   # dnf install bcc-tools

   The BCC tools are installed in the /usr/share/bcc/tools/ directory.

2. Optionally, inspect the tools:

   # ll /usr/share/bcc/tools/

   ...  
   -rwxr-xr-x. 1 root root  4198 Dec 14 17:53 dcsnoop
   -rwxr-xr-x. 1 root root  3931 Dec 14 17:53 dcstat
   -rwxr-xr-x. 1 root root 20040 Dec 14 17:53 deadlock_detector
   -rw-r--r--. 1 root root  7105 Dec 14 17:53 deadlock_detector.c
   drwxr-xr-x. 3 root root  8192 Mar 11 10:28 doc
   -rwxr-xr-x. 1 root root  7588 Dec 14 17:53 execsnoop
   -rwxr-xr-x. 1 root root  6373 Dec 14 17:53 ext4dist
   -rwxr-xr-x. 1 root root 10401 Dec 14 17:53 ext4slower
   ...

   The doc directory in the listing above contains documentation for each tool.

19.3. USING SELECTED BCC-TOOLS FOR PERFORMANCE ANALYSES
This section describes how to use certain pre-created programs from the BPF Compiler Collection (BCC) library to efficiently and securely analyze the system performance on the per-event basis. The set of pre-created programs in the BCC library can serve as examples for creation of additional programs.

**Prerequisites**

- Installed bcc-tools package
- Root permissions

**Using execsnoop to examine the system processes**

1. Execute the `execsnoop` program in one terminal:

   ```
   # /usr/share/bcc/tools/execsnoop
   ```

2. In another terminal execute for example:

   ```
   $ ls /usr/share/bcc/tools/doc/
   ```

   The above creates a short-lived process of the `ls` command.

3. The terminal running `execsnoop` shows the output similar to the following:

   ```
   PCOMM PID    PPID   RET  ARGS
   ls    8382   8287     0 /usr/bin/ls --color=auto /usr/share/bcc/tools/doc/ ...
   ```

   The `execsnoop` program prints a line of output for each new process, which consumes system resources. It even detects processes of programs that run very shortly, such as `ls`, and most monitoring tools would not register them.

   The `execsnoop` output displays the following fields:

   - **PCOMM** - The parent process name. (`ls`)
   - **PID** - The process ID. (`8382`)
   - **PPID** - The parent process ID. (`8287`)
   - **RET** - The return value of the `exec()` system call (`0`), which loads program code into new processes.
   - **ARGS** - The location of the started program with arguments.

   To see more details, examples, and options for `execsnoop`, refer to the `/usr/share/bcc/tools/doc/execsnoop_example.txt` file.

   For more information about `exec()`, see `exec(3)` manual pages.

**Using opensnoop to track what files a command opens**

1. Execute the `opensnoop` program in one terminal:

   ```
   # /usr/share/bcc/tools/opensnoop -n uname
   ```
The above prints output for files, which are opened only by the process of the `uname` command.

2. In another terminal execute:

   ```
   $ uname
   ```

   The command above opens certain files, which are captured in the next step.

3. The terminal running `opensnoop` shows the output similar to the following:

   ```
   PID COMM FD ERR PATH
   8596  uname  3  0  /etc/ld.so.cache
   8596  uname  3  0  /lib64/libc.so.6
   8596  uname  3  0  /usr/lib/locale/locale-archive
   ... 
   ```

   The `opensnoop` program watches the `open()` system call across the whole system, and prints a line of output for each file that `uname` tried to open along the way.

   The `opensnoop` output displays the following fields:

   - **PID** - The process ID. (8596)
   - **COMM** - The process name. (uname)
   - **FD** - The file descriptor - a value that `open()` returns to refer to the open file. (3)
   - **ERR** - Any errors.
   - **PATH** - The location of files that `open()` tried to open.

   If a command tries to read a non-existent file, then the FD column returns -1 and the ERR column prints a value corresponding to the relevant error. As a result, `opensnoop` can help you identify an application that does not behave properly.

   To see more details, examples, and options for `opensnoop`, refer to the `/usr/share/bcc/tools/doc/opensnoop_example.txt` file.

   For more information about `open()`, see `open(2)` manual pages.

**Using biotop to examine the I/O operations on the disk**

1. Execute the `biotop` program in one terminal:

   ```
   # /usr/share/bcc/tools/biotop 30
   ```

   The command enables you to monitor the top processes, which perform I/O operations on the disk. The argument ensures that the command will produce a 30 second summary.

   **NOTE**

   When no argument provided, the output screen by default refreshes every 1 second.

2. In another terminal execute for example:
The command above reads the content from the local hard disk device and writes the output to the /dev/zero file. This step generates certain I/O traffic to illustrate biotop.

3. The terminal running biotop shows the output similar to the following:

<table>
<thead>
<tr>
<th>PID</th>
<th>COMM</th>
<th>DISK</th>
<th>I/O</th>
<th>Kbytes</th>
<th>AVGms</th>
</tr>
</thead>
<tbody>
<tr>
<td>9568</td>
<td>dd</td>
<td>vda</td>
<td>16294</td>
<td>14440636.0</td>
<td>3.69</td>
</tr>
<tr>
<td>48</td>
<td>kswapd0</td>
<td>vda</td>
<td>1763</td>
<td>120696.0</td>
<td>1.65</td>
</tr>
<tr>
<td>7571</td>
<td>gnome-shell</td>
<td>vda</td>
<td>834</td>
<td>83612.0</td>
<td>0.33</td>
</tr>
<tr>
<td>1891</td>
<td>gnome-shell</td>
<td>vda</td>
<td>1379</td>
<td>19792.0</td>
<td>0.15</td>
</tr>
<tr>
<td>7515</td>
<td>Xorg</td>
<td>vda</td>
<td>280</td>
<td>9940.0</td>
<td>0.28</td>
</tr>
<tr>
<td>7579</td>
<td>lvmpipe-1</td>
<td>vda</td>
<td>228</td>
<td>6928.0</td>
<td>0.19</td>
</tr>
<tr>
<td>9515</td>
<td>gnome-control-c</td>
<td>vda</td>
<td>62</td>
<td>6444.0</td>
<td>0.43</td>
</tr>
<tr>
<td>8112</td>
<td>gnome-terminal-</td>
<td>vda</td>
<td>67</td>
<td>2572.0</td>
<td>1.54</td>
</tr>
<tr>
<td>7807</td>
<td>gnome-software</td>
<td>vda</td>
<td>31</td>
<td>2336.0</td>
<td>0.73</td>
</tr>
<tr>
<td>9578</td>
<td>awk</td>
<td>vda</td>
<td>17</td>
<td>2228.0</td>
<td>0.66</td>
</tr>
<tr>
<td>7578</td>
<td>lvmpipe-0</td>
<td>vda</td>
<td>156</td>
<td>2204.0</td>
<td>0.07</td>
</tr>
<tr>
<td>9581</td>
<td>pgrep</td>
<td>vda</td>
<td>58</td>
<td>1748.0</td>
<td>0.42</td>
</tr>
<tr>
<td>7531</td>
<td>InputThread</td>
<td>vda</td>
<td>30</td>
<td>1200.0</td>
<td>0.48</td>
</tr>
<tr>
<td>7504</td>
<td>gdbus</td>
<td>vda</td>
<td>3</td>
<td>1164.0</td>
<td>0.30</td>
</tr>
<tr>
<td>1983</td>
<td>lvmpipe-1</td>
<td>vda</td>
<td>39</td>
<td>724.0</td>
<td>0.08</td>
</tr>
<tr>
<td>1982</td>
<td>lvmpipe-0</td>
<td>vda</td>
<td>36</td>
<td>652.0</td>
<td>0.06</td>
</tr>
</tbody>
</table>

The biotop output displays the following fields:

- **PID** - The process ID. (9568)
- **COMM** - The process name. (dd)
- **DISK** - The disk performing the read operations. (vda)
- **I/O** - The number of read operations performed. (16294)
- **Kbytes** - The amount of Kbytes reached by the read operations. (14,440,636)
- **AVGms** - The average I/O time of read operations. (3.69)

To see more details, examples, and options for biotop, refer to the /usr/share/bcc/tools/doc/biotop_example.txt file.

For more information about dd, see dd(1) manual pages.

**Using xfsslower to expose unexpectedly slow file system operations**

1. Execute the xfsslower program in one terminal:

```
# /usr/share/bcc/tools/xfsslower 1
```

The command above measures the time the XFS file system spends in performing read, write, open or sync (fsync) operations. The 1 argument ensures that the program shows only the operations that are slower than 1 ms.
NOTE

When no arguments provided, `xfsslower` by default displays operations slower than 10 ms.

2. In another terminal execute, for example, the following:

```bash
$ vim text
```

The command above creates a text file in the `vim` editor to initiate certain interaction with the XFS file system.

3. The terminal running `xfsslower` shows something similar upon saving the file from the previous step:

```
TIME     COMM           PID    T BYTES   OFF_KB   LAT(ms) FILENAME
13:07:14 b'bash'        4754   R 256      0   7.11 b'vim'
13:07:14 b'vim'         4754   R 832      0   4.03 b'libgpm.so.2.1.0'
13:07:14 b'vim'         4754   R 32       20   1.04 b'libgpm.so.2.1.0'
13:07:14 b'vim'         4754   R 1982     0   2.30 b'vimrc'
13:07:14 b'vim'         4754   R 1393     0   2.52 b'getscriptPlugin.vim'
13:07:45 b'vim'         4754   S 0        0   6.71 b'text'
13:07:45 b'pool'        2588   R 16       0   5.58 b'text'
... 
```

Each line above represents an operation in the file system, which took more time than a certain threshold. `xfsslower` is good at exposing possible file system problems, which can take form of unexpectedly slow operations.

The `xfsslower` output displays the following fields:

- **COMM** - The process name. (b'bash')
- **T** - The operation type. (R)
  - Read
  - Write
  - Sync
- **OFF_KB** - The file offset in KB. (0)
- **FILENAME** - The file being read, written, or synced.

To see more details, examples, and options for `xfsslower`, refer to the `/usr/share/bcc/tools/doc/xfsslower_example.txt` file.

For more information about `fsync`, see `fsync(2)` manual pages.
You can increase the protection of your system by utilizing components of the kernel integrity subsystem. The following sections introduce the relevant components and provide guidance on their configuration.

**NOTE**
The features involving cryptographic signatures are available only for the intra Red Hat usage. This is because only the certificates for Red Hat signature keys are present in the kernel keyring. The user can use any hash feature, however, be complete tamper-proof mechanisms will not finish successfully.

## 20.1. THE KERNEL INTEGRITY SUBSYSTEM

The integrity subsystem is a part of the kernel that is responsible for maintaining the overall system data integrity. This subsystem helps to keep the state of a certain system the same from the time it was built and thus prevents undesired modification on specific system files.

The kernel integrity subsystem consists of two major components:

**Integrity Measurement Architecture (IMA)**
- Measures file content whenever it is executed or opened by cryptographically hashing or signing with cryptographic keys. The keys are stored in the kernel keyring subsystem.
- Places the measured values within the kernel’s memory space thereby it prevents any modification from the users of the system.
- Allows local and remote parties to verify the measured values.
- Provides local validation of the current content of files against the values previously stored in the measurement list within the kernel memory. This extension forbids to conduct any operation on a specific file in case the current and the previous measures do not match.

**Extended Verification Module (EVM)**
- Protects extended attributes of files (also known as xattr) that are related to the system security, like IMA measurements and SELinux attributes. The component cryptographically hashes their corresponding values or signs them with cryptographic keys. The keys are stored in the kernel keyring subsystem.

The kernel integrity subsystem can harness the Trusted Platform Module (TPM) to harden the system security even more. TPM is a specification by the Trusted Computing Group (TCG) for important cryptographic functions. TPMs are usually built as dedicated hardware that is attached to the platform’s motherboard and prevents software-based attacks by providing cryptographic functions from a protected and tamper-proof area of the hardware chip. Some of the TPM features are:
  - Random-number generator
  - Generator and secure storage for cryptographic keys
  - Hashing generator
Remote attestation

Additional resources
- *Integrity Measurement Architecture (IMA)*
- *Trusted Computing Group resources*
- Security hardening
- Basic and advanced configuration of Security-Enhanced Linux (SELinux)

### 20.2. INTEGRITY MEASUREMENT ARCHITECTURE

Integrity Measurement Architecture (IMA) is a component of the kernel integrity subsystem. IMA aims to maintain the contents of local files. Specifically, IMA measures, stores, and appraises file hashes before they are accessed, which prevents the reading and execution of unreliable data. Thereby, usage of IMA with digital signatures enhances the security of the system and prevents from measurement data forgery. The scenario assumes the attacker does not possess the original signature key.

Additional resources
- *Integrity Measurement Architecture (IMA)*
- Security hardening
- Basic and advanced configuration of Security-Enhanced Linux (SELinux)

### 20.3. EXTENDED VERIFICATION MODULE

Extended Verification Module (EVM) is a component of the kernel integrity subsystem that monitors changes in file extended attributes (`xattr`). Many security-oriented technologies store sensitive file information, such as content hashes and signatures, in extended attributes of files. EVM goes one step further and hashes or signs those extended attributes. The resulting data is validated every time the extended attribute is used. For example, when IMA appraises the file.

Additional resources
- *Integrity Measurement Architecture (IMA)*
- Security hardening
- Basic and advanced configuration of Security-Enhanced Linux (SELinux)

### 20.4. TRUSTED AND ENCRYPTED KEYS

The following section introduces trusted and encrypted keys as an important part of enhancing system security.

*Trusted* and *encrypted keys* are variable-length symmetric keys generated by the kernel that utilize the kernel keyring service. The fact that this type of keys never appear in the user space in an unencrypted form means that their integrity can be verified, which in turn means that they can be used, for example, by the extended verification module (EVM) to verify and confirm the integrity of a running system. User-level programs can only access the keys in the form of encrypted *blobs*. 
Trusted keys need a hardware component: the Trusted Platform Module (TPM) chip, which is used to both create and encrypt (seal) the keys. The TPM seals the keys using a 2048-bit RSA key called the storage root key (SRK).

**NOTE**

To use a TPM 1.2 specification, enable and activate it through a setting in the machine firmware or by using the `tpm_setactive` command from the `tpm-tools` package of utilities. Also, the TrouSers software stack needs to be installed and the `tcsd` daemon needs to be running to communicate with the TPM (dedicated hardware). The `tcsd` daemon is part of the TrouSers suite, which is available through the `trousers` package. The more recent and backward incompatible TPM 2.0 uses a different software stack, where the `tpm2-tools` or `ibm-tss` utilities provide access to the dedicated hardware.

In addition to that, the user can seal the trusted keys with a specific set of the TPM's platform configuration register (PCR) values. PCR contains a set of integrity-management values that reflect the firmware, boot loader, and operating system. This means that PCR-sealed keys can only be decrypted by the TPM on the same system on which they were encrypted. However, once a PCR-sealed trusted key is loaded (added to a keyring), and thus its associated PCR values are verified, it can be updated with new (or future) PCR values, so that a new kernel, for example, can be booted. A single key can also be saved as multiple blobs, each with different PCR values.

Encrypted keys do not require a TPM, as they use the kernel Advanced Encryption Standard (AES), which makes them faster than trusted keys. Encrypted keys are created using kernel-generated random numbers and encrypted by a master key when they are exported into user-space blobs. The master key is either a trusted key or a user key. If the master key is not trusted, the encrypted key is only as secure as the user key used to encrypt it.

### 20.5. WORKING WITH TRUSTED KEYS

The following section describes how to create, export, load or update trusted keys with the `keyctl` utility to improve the system security.

**Prerequisites**

- For the 64-bit ARM architecture and IBM Z, the `trusted` kernel module needs to be loaded. For more information on how to load kernel modules, see Managing kernel modules.

- Trusted Platform Module (TPM) needs to be enabled and active. For more information about TPM see, The kernel integrity subsystem and Trusted and encrypted keys.

**Procedure**

1. To create a trusted key using a TPM, execute:

   ```
   # keyctl add trusted <name> "new <key_length> [options]" <key_ring>
   ```

   - Based on the syntax, construct an example command as follows:

     ```
     # keyctl add trusted kmk "new 32" @u 642500861
     ```
The command creates a trusted key called `kmk` with the length of 32 bytes (256 bits) and places it in the user keyring (@u). The keys may have a length of 32 to 128 bytes (256 to 1024 bits).

2. To list the current structure of the kernel keyrings:
   ```bash
   # keyctl show
   Session Keyring
   -3 --alswrv 500 500 keyring: ses 97833714 --alswrv 500 -1 \ keyring: uid.1000
   642500861 --alswrv 500 500 \ trusted: kmk
   ```

3. To export the key to a user-space blob, execute:
   ```bash
   # keyctl pipe 642500861 > kmk.blob
   ```
   The command uses the `pipe` subcommand and the serial number of `kmk`.

4. To load the trusted key from the user-space blob, use the `add` subcommand with the blob as an argument:
   ```bash
   # keyctl add trusted kmk "load `cat kmk.blob`" @u
   ``
   ```bash
   268728824
   ```

5. Create secure encrypted keys based on the TPM-sealed trusted key:
   ```bash
   # keyctl add encrypted <pass:quotes[name]> "new [format] <pass:quotes[key_type]>:
   <pass:quotes[primary_key_name]> <pass:quotes[keyleNGTH]>
   <pass:quotes[key_ring]>
   ```
   - Based on the syntax, generate an encrypted key using the already created trusted key:
     ```bash
     # keyctl add encrypted encr-key "new trusted:kmk 32" @u
     ``
     ```bash
     159771175
     ```
     The command uses the TPM-sealed trusted key (`kmk`), produced in the previous step, as a `primary key` for generating encrypted keys.

Additional resources
- keyctl(1) manual page
- Trusted and encrypted keys
- Kernel Key Retention Service
- The kernel integrity subsystem

20.6. WORKING WITH ENCRYPTED KEYS

The following section describes managing encrypted keys to improve the system security on systems where a Trusted Platform Module (TPM) is not available.

Prerequisites
For the 64-bit ARM architecture and IBM Z, the encrypted-keys kernel module needs to be loaded. For more information on how to load kernel modules, see Managing kernel modules.

Procedure

1. Use a random sequence of numbers to generate a user key:

   ```shell
   # keyctl add user kmk-user "$(dd if=/dev/urandom bs=1 count=32 2>/dev/null)" @u 427069434
   ```

   The command generates a user key called kmk-user which acts as a primary key and is used to seal the actual encrypted keys.

2. Generate an encrypted key using the primary key from the previous step:

   ```shell
   # keyctl add encrypted encr-key "new user:kmk-user 32" @u 1012412758
   ```

3. Optionally, list all keys in the specified user keyring:

   ```shell
   # keyctl list @u
   2 keys in keyring:
   427069434: --alswrv 1000 1000 user: kmk-user
   1012412758: --alswrv 1000 1000 encrypted: encr-key
   ```

   **IMPORTANT**

   Keep in mind that encrypted keys that are not sealed by a trusted primary key are only as secure as the user primary key (random-number key) that was used to encrypt them. Therefore, the primary user key should be loaded as securely as possible and preferably early during the boot process.

Additional resources

- [keyctl(1) manual page](#)
- [Kernel Key Retention Service](#)

## 20.7. ENABLING INTEGRITY MEASUREMENT ARCHITECTURE AND EXTENDED VERIFICATION MODULE

Integrity measurement architecture (IMA) and extended verification module (EVM) belong to the kernel integrity subsystem and enhance the system security in various ways. The following section describes how to enable and configure IMA and EVM to improve the security of the operating system.

**IMPORTANT**

This procedure works only when using certificates and keys issued by Red Hat.

Prerequisites

- Verify that the `securityfs` filesystem is mounted on the `/sys/kernel/security/` directory and the `/sys/kernel/security/integrity/ima/` directory exists.
mount ...
securityfs on /sys/kernel/security type securityfs (rw,nosuid,nodev,noexec,relatime)
...

- Verify that the **systemd** service manager is already patched to support IMA and EVM on boot time:

```
# dmesg | grep -i -e EVM -e IMA
[    0.000000] kvm-clock: cpu 0, msr 23601001, primary cpu clock
[    0.000000] Using crashkernel=1G-4G:192M,4G-64G:256M,64G-:512M, the size chosen is a best effort estimation.
[    0.911527] ima: No TPM chip found, activating TPM-bypass!
[    0.911538] ima: Allocated hash algorithm: sha1
[    0.911580] evm: Initialising EVM extended attributes:
[    0.911581] evm: security.selinux
[    0.911581] evm: security.ima
[    0.911582] evm: security.capability
[    0.911582] evm: HMAC attrs: 0x1
[    1.715151] systemd[1]: systemd 239 running in system mode. (+PAM +AUDIT +SELINUX +IMA -APPARMOR +SMACK +SYSSVINIT +UPTYPE +LIBCRYPTSETUP +GCRYPT +GNUTLS +ACL +XZ +LZ4 +SECCOMP +BLKID +ELFUTILS +KMOD +IDN2 -IDN +PCRE2 default-hierarchy=legacy)
[    3.824198] fbcon: qxldrmfb (fb0) is primary device
[    4.673457] PM: Image not found (code -22)
[    6.549966] systemd[1]: systemd 239 running in system mode. (+PAM +AUDIT +SELINUX +IMA -APPARMOR +SMACK +SYSSVINIT +UPTYPE +LIBCRYPTSETUP +GCRYPT +GNUTLS +ACL +XZ +LZ4 +SECCOMP +BLKID +ELFUTILS +KMOD +IDN2 -IDN +PCRE2 default-hierarchy=legacy)
```

### Procedure

1. Add the following kernel command line parameters:

```
# grubby --update-kernel=/boot/vmlinuz-$\{uname -r\} --args="ima_policy=appraise_tcb ima_appraise=fix evm=fix"
```

The command enables IMA and EVM in the *fix* mode for the current boot entry and allows users to gather and update the IMA measurements.

The `ima_policy=appraise_tcb` kernel command line parameter ensures that the kernel uses the default Trusted Computing Base (TCB) measurement policy and the appraisal step. The appraisal part forbids access to files, whose prior and current measures do not match.

2. Reboot to make the changes come into effect.

3. Optionally, verify that the parameters have been added to the kernel command line:

```
# cat /proc/cmdline
```
4. Create a kernel master key to protect the EVM key:

```
# keyctl add user kmk "$(dd if=/dev/urandom bs=1 count=32 2> /dev/null)" @u
748544121
```

The kernel master key (kmk) is kept entirely in the kernel space memory. The 32-byte long value of the kernel master key kmk is generated from random bytes from the /dev/urandom file and placed in the user (@u) keyring. The key serial number is on the second line of the previous output.

5. Create an encrypted EVM key based on the kmk key:

```
# keyctl add encrypted evm-key "new user:kmk 64" @u
641780271
```

The command uses kmk to generate and encrypt a 64-byte long user key (named evm-key) and places it in the user (@u) keyring. The key serial number is on the second line of the previous output.

**IMPORTANT**

It is necessary to name the user key as evm-key because that is the name the EVM subsystem is expecting and is working with.

6. Create a directory for exported keys.

```
# mkdir -p /etc/keys/
```

7. Search for the kmk key and export its value into a file:

```
# keyctl pipe $(keyctl search @u user kmk) > /etc/keys/kmk
```

The command places the unencrypted value of the kernel master key (kmk) into a file of previously defined location (/etc/keys/).

8. Search for the evm-key user key and export its value into a file:

```
# keyctl pipe $(keyctl search @u encrypted evm-key) > /etc/keys/evm-key
```

The command places the encrypted value of the user evm-key key into a file of arbitrary location. The evm-key has been encrypted by the kernel master key earlier.

9. Optionally, view the newly created keys:

```
# keyctl show
Session Keyring
974575405 --alswrv 0 0 keyring: ses 299489774 --alswrv 0 65534 \ keyring: uid.0
748544121 --alswrv 0 0 \ user: kmk
641780271 --alswrv 0 0 \ encrypted: evm-key
```
You should be able to see a similar output.

10. Activate EVM:

```
# echo 1 > /sys/kernel/security/evm
```

Verification step

- You can verify that EVM has been initialized by entering:

```
# dmesg | tail -1
[...] evm: key initialized
```

**NOTE**

If the system is rebooted, the keys are removed from the keyring. In such a case, you can import the already exported `kmk` and `evm-key` keys.

Procedure

1. Add the user `kmk` key (already exported to the `/etc/keys/kmk` file in step 7).

```
# keyctl add user kmk "$(cat /etc/keys/kmk)" @u
451342217
```

```
# keyctl show
Session Keyring
695566911 --alswrv 0 0 keyring: ses 58982213 --alswrv 0 65534
keyring: uid.0 451342217 --alswrv 0 0 user: kmk
```

2. Import the user `evm-key` key (already exported to the `/etc/keys/evm-key` file in step 8).

```
# keyctl add encrypted evm-key "load $(cat /etc/keys/evm-key)" @u
924537557
```

```
# keyctl show
Session Keyring
695566911 --alswrv 0 0 keyring: ses 58982213 --alswrv 0 65534
keyring: uid.0 451342217 --alswrv 0 0 user: kmk
924537557 --alswrv 0 0 _ encrypted: evm-key
```

Additional resources

- Extended verification module
- Integrity measurement architecture
20.8. COLLECTING FILE HASHES WITH INTEGRITY MEASUREMENT ARCHITECTURE

The first level of operation of integrity measurement architecture (IMA) is the measurement phase that allows to create file hashes and store them as extended attributes (xattrs) of those files. The following section describes how to create and inspect the files' hashes.

Prerequisites

- Enable integrity measurement architecture (IMA) and extended verification module (EVM) as described in Enabling integrity measurement architecture and extended verification module.
- Verify that the ima-evm-utils, attr, and keyutils packages are already installed:

```bash
# dnf install ima-evm-utils attr keyutils
Updating Subscription Management repositories.
This system is registered to Red Hat Subscription Management, but is not receiving updates.
You can use subscription-manager to assign subscriptions.
Last metadata expiration check: 0:58:22 ago on Fri 14 Feb 2020 09:58:23 AM CET.
Package ima-evm-utils-1.1-5.el8.x86_64 is already installed.
Package attr-2.4.48-3.el8.x86_64 is already installed.
Package keyutils-1.5.10-7.el8.x86_64 is already installed.
Dependencies resolved.
Nothing to do.
Complete!
```

Procedure

1. Create a test file:

```bash
# echo <Test_text> > test_file
```

IMA and EVM ensure that the test_file example file has assigned hash values that are stored as its extended attributes.

2. Inspect extended attributes of the file:

```bash
# getfattr -m . -d test_file
# file: test_file
security.vhm=0sAnDly4VPA0HArpPO/EqiutnNyBql
security.ima=0sAQOEDeuUnWzwwKYk+n66h/vby3eD
security.selinux="unconfined_u:object_r:admin_home_t:s0"
```

The example output shows extended attributes related to SELinux and the IMA and EVM hash values. EVM adds a security.vhm extended attribute related to the other attributes. You can use the evmctl utility on security.vhm to generate either an RSA based digital signature or a Hash-based Message Authentication Code (HMAC-SHA1). To successfully conduct the previous operation you need a valid trusted or encrypted key, which is stored in the kernel keyring. This configuration prevents offline tampering attacks on the extended attributes.
Additional resources

- Security hardening
- Extended verification module
- Integrity measurement architecture
You can use the `kernel_settings` role to configure kernel parameters on multiple clients at once. This solution:

- Provides a friendly interface with efficient input setting.
- Keeps all intended kernel parameters in one place.

After you run the `kernel_settings` role from the control machine, the kernel parameters are applied to the managed systems immediately and persist across reboots.

**IMPORTANT**

Note that RHEL System Role delivered over RHEL channels are available to RHEL customers as an RPM package in the default AppStream repository. RHEL System Role are also available as a collection to customers with Ansible subscriptions over Ansible Automation Hub.

### 21.1. INTRODUCTION TO THE KERNEL SETTINGS ROLE

RHEL System Roles is a set of roles that provide a consistent configuration interface to remotely manage multiple systems.

RHEL System Roles were introduced for automated configurations of the kernel using the `kernel_settings` System Role. The `rhel-system-roles` package contains this system role, and also the reference documentation.

To apply the kernel parameters on one or more systems in an automated fashion, use the `kernel_settings` role with one or more of its role variables of your choice in a playbook. A playbook is a list of one or more plays that are human-readable, and are written in the YAML format.

With the `kernel_settings` role you can configure:

- The kernel parameters using the `kernel_settings_sysctl` role variable
- Various kernel subsystems, hardware devices, and device drivers using the `kernel_settings_sysfs` role variable
- The CPU affinity for the `systemd` service manager and processes it forks using the `kernel_settings_systemd_cpu_affinity` role variable
- The kernel memory subsystem transparent hugepages using the `kernel_settings_transparent_hugepages` and `kernel_settings_transparent_hugepages_defrag` role variables

**Additional resources**

- `README.md` and `README.html` files in the `/usr/share/doc/rhel-system-roles/kernel_settings/` directory
- Working with playbooks
- How to build your inventory
21.2. APPLYING SELECTED KERNEL PARAMETERS USING THE KERNEL_SETTINGS ROLE

Follow these steps to prepare and apply an Ansible playbook to remotely configure kernel parameters with persisting effect on multiple managed operating systems.

Prerequisites

- You have root permissions.
- Entitled by your RHEL subscription, you installed the ansible-core and rhel-system-roles packages on the control machine.
- An inventory of managed hosts is present on the control machine and Ansible is able to connect to them.

IMPORTANT

RHEL 8.0 - 8.5 provided access to a separate Ansible repository that contains Ansible Engine 2.9 for automation based on Ansible. Ansible Engine contains command-line utilities such as ansible, ansible-playbook; connectors such as docker and podman; and the entire world of plugins and modules. For information on how to obtain and install Ansible Engine, refer to How do I Download and Install Red Hat Ansible Engine?.

RHEL 8.6 and 9.0 has introduced Ansible Core (provided as ansible-core RPM), which contains the Ansible command-line utilities, commands, and a small set of built-in Ansible plugins. The AppStream repository provides ansible-core, which has a limited scope of support. You can learn more by reviewing Scope of support for the ansible-core package included in the RHEL 9 AppStream.

Procedure

1. Optionally, review the inventory file for illustration purposes:

```
# cat /home/jdoe/<ansible_project_name>/inventory
[testingservers]
pdoe@192.168.122.98
fdoe@192.168.122.226

[db-servers]
db1.example.com
db2.example.com

[webservers]
web1.example.com
web2.example.com
192.0.2.42
```

The file defines the [testing.servers] group and other groups. It allows you to run Ansible more effectively against a specific set of systems.

2. Create a configuration file to set defaults and privilege escalation for Ansible operations.
   a. Create a new YAML file and open it in a text editor, for example:
b. Insert the following content into the file:

```yaml
[defaults]
inventory = ./inventory

[privilege_escalation]
become = true
become_method = sudo
become_user = root
become_ask_pass = true
```

The `[defaults]` section specifies a path to the inventory file of managed hosts. The `[privilege_escalation]` section defines that user privileges be shifted to `root` on the specified managed hosts. This is necessary for successful configuration of kernel parameters. When Ansible playbook is run, you will be prompted for user password. The user automatically switches to `root` by means of `sudo` after connecting to a managed host.

3. Create an Ansible playbook that uses the `kernel_settings` role.

   a. Create a new YAML file and open it in a text editor, for example:

   ```yaml
   # vi /home/jdoe/<ansible_project_name>/kernel-roles.yml
   ```

   This file represents a playbook and usually contains an ordered list of tasks, also called `plays`, that are run against specific managed hosts selected from your `inventory` file.

   b. Insert the following content into the file:

   ```yaml
   ---
   -
     hosts: testingservers
     name: "Configure kernel settings"
     roles:
     - rhel-system-roles.kernel_settings
     vars:
       kernel_settings_sysctl:
         - name: fs.file-max
           value: 400000
         - name: kernel.threads-max
           value: 65536
       kernel_settings_sysfs:
         - name: /sys/class/net/lo/mtu
           value: 65000
       kernel_settings_transparent_hugepages: madvise
   ```

   The `name` key is optional. It associates an arbitrary string with the play as a label and identifies what the play is for. The `hosts` key in the play specifies the hosts against which the play is run. The value or values for this key can be provided as individual names of managed hosts or as groups of hosts as defined in the `inventory` file.

   The `vars` section represents a list of variables containing selected kernel parameter names and values to which they have to be set.
The **roles** key specifies what system role is going to configure the parameters and values mentioned in the **vars** section.

**NOTE**

You can modify the kernel parameters and their values in the playbook to fit your needs.

4. Optionally, verify that the syntax in your play is correct.

```bash
# ansible-playbook --syntax-check kernel-roles.yml
playbook: kernel-roles.yml
```

This example shows the successful verification of a playbook.

5. Execute your playbook.

```bash
# ansible-playbook kernel-roles.yml
...
BECOME password:
PLAY [Configure kernel settings]
**********************************************************************************
PLAY RECAP
********************************************************************************************************
fdoe@192.168.122.226  : ok=10  changed=4  unreachable=0  failed=0  skipped=6
rescued=0  ignored=0
PDOE@192.168.122.98  : ok=10  changed=4  unreachable=0  failed=0  skipped=6
rescued=0  ignored=0
```

Before Ansible runs your playbook, you are going to be prompted for your password and so that a user on managed hosts can be switched to **root**, which is necessary for configuring kernel parameters.

The recap section shows that the play finished successfully (**failed=0**) for all managed hosts, and that 4 kernel parameters have been applied (**changed=4**).

6. Restart your managed hosts and check the affected kernel parameters to verify that the changes have been applied and persist across reboots.

**Additional resources**

- Preparing a control node and managed nodes to use RHEL System Roles
- README.html and README.md files in the `/usr/share/doc/rhel-system-roles/kernel_settings/` directory
- Build Your Inventory
● Configuring Ansible
● Working With Playbooks
● Using Variables
● Roles